



# **STF- Energy Efficient Scheduling Algorithm for WSN**

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**ABSTRACT:** As it is well known that sensors are deployed to perform an important task of gathering sensory information from its surrounding which may be very crucial for monitoring and analysis purpose, thus it should be taken care that sensors remain alive for a long time to sense data. All the sensor nodes have a limited supply of power which keeps them alive. Therefore, it is very necessary to reduce the energy consumed by sensors. This paper aims at increasing the lifetime of sensor nodes by adopting an efficient scheduling algorithm named Shortest Time First (STF). Merging this concept with the existing “intelligent sleeping mechanism” of Low Energy Adaptive Clustering Hierarchy Centralised Sleeping (LEACH CS) [3] protocol helps to increase the network lifetime as it allows sensor nodes to sleep for some more time that has been saved using our new scheduling mechanism. According to STF, all the cluster heads will be scheduled to send their aggregated data on the basis of time taken for their signal to reach base station, which would help in reducing the waiting time of other cluster heads in a particular round and ultimately conserve energy. Simulation Results show that proposed STF schedule in LEACH CS increases the energy efficiency by approximately 20% more than LEACH CS [3] algorithm.

**KEYWORDS:** STF (Shortest Time First), Network Lifetime, QoS (Quality Of Service), Average waiting time, Sleeping time, Threshold value.

## **I. INTRODUCTION**

With advances in micro-electronics and communication technology, it is now possible to manufacture low cost sensors. These sensors are used by the sensor nodes to capture data from their surroundings. A Wireless Sensor Network is comprised of a large number of sensor nodes which may vary from a few to several hundreds. Each of these sensor nodes have limited energy as power is supplied through a battery which discharge after a certain period of time. Also, these sensors are usually placed in remote areas, where battery cannot be replaced regularly. Thus, our main aim should be to conserve the energy of sensors which would extend the lifetime of sensors and the network as well. It is necessary to keep these sensor nodes alive for a longer duration so that they could sense values and transmit it to the base station. [9]

In the recent years, there has been a significant expansion in the application field of wireless sensor networks. Initially, these were mainly used for military applications such as surveillance but today such networks are highly used for industrial process monitoring, weather monitoring, agriculture, railways and so on. The main aim of the sensor network is to collect information from its environment. Fig1 shows network model of wireless sensors in which all sensor nodes can directly communicate to base station which further updates the information on internet for users.

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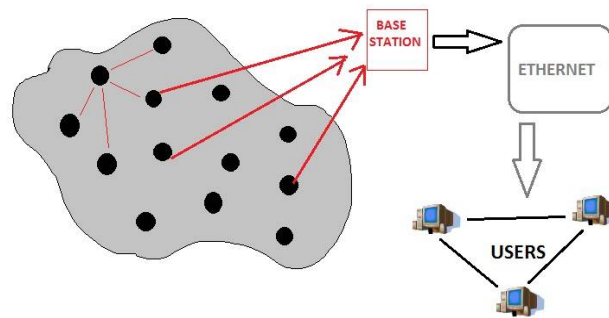


Fig 1: A wireless sensor network

There are different routing protocols which can be used for routing in a wireless sensor network. These can be categorised into hierarchical[12][13][14][15]based or data centric or location based protocols [4] [5].In this paper, we have worked on a hierarchy based protocol called as Low Energy Adaptive Clustering Hierarchy Centralised(LEACH C) protocol [2], in which cluster formation takes place and the cluster head after receiving the data from its nearby nodes, interacts directly with the base station [11].This helps in saving the energy of remaining nodes in that cluster. LEACH CS [3] protocol has been proposed which uses centralised cluster formation technique and intelligent sleeping mechanism for deciding which nodes to be switched to sleep mode on the basis of the significance of data. It has been explained in detail in Section III of this paper. Actually, LEACH CS protocol uses TDMA schedule for communication between cluster head of the nodes and the base station. In this paper, we have suggested a new scheduling algorithm STF(Shortest Time First) which helps to increase the lifetime of sensors by increasing their sleep time, based on the concepts of CPU scheduling (SJF-Shortest Job First) in Operating System [6].The proposed algorithm mainly focuses on the scheduling mechanism adopted by the cluster heads to transmit signals to the base station.

According to STF, the sensors will send their data according to a specific schedule which uses time and energy level of sensors to sort them in ascending order. This schedule will reduce the time for which a sensor may have to wait for its turn to send signal, called as waiting time. The life time of sensors can be increased by allowing sensors to sleep for more time and also helps to conserve energy of sensors.

A lot of research work has been done to minimise the energy consumption by sensors to prolong the lifetime of networks. This paper aims at conserving the energy of sensors by using an improved scheduling technique which has been explained in detail in Section III of this paper. Section II provides details of the related works which we have studied to propose our algorithm.

## II. RELATED WORKS

### Low Energy Adaptive Clustering hierarchy (LEACH)

Many research papers have been published for improving the efficiency of wireless sensor networks. HierarchicalBased protocol has always been considered as an efficient method of routing as it involves cluster formation which helps in saving the energy of sensor nodes [8].

The working of LEACH [1] can be divided into two phases, the *Setup phase* and the *Steady –State phase*. The selection of cluster head and the formation of clusters are done in the *setup phase*. Some of the sensor nodes are selected to serve as cluster head (CH )for the current round, and the remaining nodes form clusters by joining the proper cluster head on the basis of signal strength from CH's .In the setup phase, only a fraction of nodes, elect themselves as the cluster head for the current round. After that the cluster head has been selected, the remaining sensor nodes can start transmitting data to their corresponding cluster head according to TDMA schedule. All this comes under the *Steady state* operation

### LEACH Centralized (LEACH C)

Although LEACH[1] helps in increasing the energy efficiency of the network by using distributed cluster formation, it does not ensure that how these cluster heads will be uniformly distributed in the network. Thus, it may happen that



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selected CHs will be concentrated only in some part of network while the remaining nodes may not have any CH nearby to them.

This problem is solved by Centralized LEACH which ensures uniform distribution of CHs throughout the network. LEACH C [2] also performs operation similar to LEACH except for selection procedure of cluster heads. This means that only the setup phase of centralized LEACH is different from LEACH. In the *setup phase*, all the nodes send the information about their location and energy level to the base station. Now, the base station has a Global knowledge of all the nodes in the network. It then calculates the average node energy value and the nodes having energy value less than this value will not be allowed to become cluster head for the current round. Now, the base station uses remaining nodes as possible CH and finds optimal clusters using simulated annealing algorithm [2]. It minimizes the amount of energy that non-cluster head nodes require to transmit their data to the cluster head, by minimizing the total sum of 'squared distances' between all the non-cluster head nodes and the closest cluster head. After selection of cluster head and formation of clusters, the Base Station (BS) sends a message to each node to tell its cluster head ID. If it matches its own ID, then it is a cluster head and creates TDMA schedule to send to its member nodes. The *Steady state phase* is identical to that of LEACH.

## LEACH CS (Centralized Sleeping)

LEACH CS [3] is an extension to centralized LEACH which makes use of Intelligent Sleeping Mechanism (ISM), for choosing nodes that should be switched to the sleeping mode. This choice is based on the quality of data sensed by sensor nodes. The cluster heads aggregate the data after collecting it from nodes in the cluster. If this data appears insignificant in a period of time, these clusters are set to sleep mode till the next data round.

The nodes in the sleeping mode do not sense any data, hence a lot of energy is conserved which increase the lifetime of sensors. Before this algorithm, the choice of sleep nodes was done on the basis of node density in a region or the number of neighboring nodes. Thus, most of them were based on redundancy of data. Hence, ISM adds quality of service to LEACH C.

### Intelligent sleeping mechanism:

- 1) After the cluster head has been selected by the base station and optimal clusters have been formed using the centralized LEACH algorithm, the non-cluster head nodes may start sending the data according to the assigned TDMA slot.
- 2) Cluster head aggregates the data after receiving it from the nodes in its cluster.
- 3) This average value is sent in the form of packet to base station. The BS after receiving the first packet from a cluster head in the current round analyses this data for its quality.
- 4) BS compares the data of first frame to a defined threshold value, which may be specified by user. If its value is less than the threshold value, then the data is not going to be significant in this round.
- 5) This means that all the nodes of the corresponding cluster can be switched to sleeping mode for the remaining time of round. The rest of the clusters continue the round and send their data to the base station in their fixed TDMA slots consequently.

The threshold value can be set by the user according to sensed application. For example, it may be temperature in a cultivation field or level of oxygen in large water bodies or may be the value of humidity in agricultural land, all such parameters will not change instantly and hence if these are under a user defined threshold, then it will not reach a critical value in given data round and will remain stable for a while.

## III. PROPOSED ALGORITHM

The proposed algorithm focuses on the scheduling mechanism used by the cluster heads of each cluster for transmitting signals to the base station. The main problem in a sensor network is the short lifetime of sensors. The energy of sensors keeps on decreasing with time. Sensors need to sleep in order to save their energy. Waiting time is the time for which a sensor waits for its turn to send signals to base station. All sensor nodes follow some schedule which is considered efficient if the waiting time for each sensor is minimum. As the waiting time of sensors decrease, they will get more time to sleep.

A solution to above problem has been achieved using an efficient scheduling algorithm, whose concept actually came from CPU scheduling in Operating System [7]. In the proposed algorithm, we will be using the same mechanism for



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sleeping that was used in LEACH CS [3] based on the threshold value concept but instead of using the TDMA scheduling method for transmitting data, we will adopt *Shortest Time First (STF)* scheduling for data transmission to base station. The complete algorithm is described below:

**STEP 1:** This is about selection of cluster head and formation of clusters which is similar to that of LEACH CS algorithm [3]. There are several rounds of data transfer that continues until all the nodes in the network are dead.

- a) Before each round, every node in the network sends its energy level and location to the BS. Depending on these energy levels, the base station sorts the nodes in descending order. It also calculates the average node energy value and the nodes having energy value less than this value will not be allowed to become cluster head for the current round.
- b) So, out of the remaining nodes in descending order of their energies, the BS chooses only 5% of the alive nodes to be the cluster head of the current round.
- c) The remaining nodes are assigned by the base station to their corresponding cluster heads by choosing the minimum square distance between the node and the cluster heads. This would find optimal clusters in the network.

Now, the non cluster head nodes in the cluster can start transmitting the data to their corresponding cluster according to the TDMA schedule sent by CH to its cluster nodes. The CHs aggregate the data sent by other nodes in cluster. All the CHs in the network have a packet to send to base station and some schedule is again required to send these packets. This is the part of algorithm where we adopted an efficient scheduling method explained below:

**STEP 2:** This step describes the modification done in the existing algorithm LEACH CS [3] in its Scheduling part.

In Operating system, when there are more than one process ready for execution at the same time, then a specific schedule is followed. Similarly, in Wireless Sensor Networks, an efficient schedule needs to be followed by the cluster heads to transmit their data turn by turn in a given round. Two CHs cannot send data together as there are chances of collision. We use a technical term in CPU scheduling called as *Average Waiting Time (AWT)*. Average waiting time is the time for which a process or a job waits for its turn to get CPU time for execution. A scheduling is considered efficient if the waiting time for each process is minimum. As we know that in LEACH CS protocol, each cluster head needs to send aggregate data to an external Base Station. This is a high energy transmission as base station are usually far away from sensor nodes [10]. Table 1 shows a list of parameters used in the algorithm.

Table 1: Various parameters used in the algorithm

Notations	Description
$S(i).xd, S(i).yd$	Location of the $i$ th node
$S(n+1).xd, S(n+1).yd$	Location of the base station.
$t_i$	Time taken by $i^{\text{th}}$ node to transmit signal to base station.
tot_time	Total time of a round
N	Total number of rounds
Count_CH	Number of Cluster heads
Count_NCH	Number of non cluster heads
$E_i, E_a$	Energy level of $i^{\text{th}}$ node, average energy
P	Optimal number of cluster heads
dist_CH	Distance of cluster heads from base station
t (i-1)	Time taken by (i-1)th node to send signal to base station
Wt(i)	Waiting time of the $i^{\text{th}}$ node.
Sleep_time	Sleeping time of nodes in a round



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## SORTcluster heads according to the time taken and energy level in ascending order.

a) **Shortest distance:** The time taken for the signals to reach the base station depends upon the distance of cluster heads from base station. The Cluster Head closest to the base station will take minimum time for its data to reach BS and the Cluster Head farthest from the base station will take maximum time for its data to reach BS. Thus, we can sort the cluster heads according to the time taken for their data to reach base station which is indirectly based on the distance of CHs from BS as:

$$\text{Time taken by signals to reach BS} = \frac{\text{distance of CHs from BS}}{\text{Propagation speed of signals}}$$

Where, propagation speed of wireless signals is given by speed of radio waves. Radio frequency-based communication is mostly used in WSN applications and Speed of radio waves is equal to the speed of light. Considering the disturbance through space,

$$\begin{aligned} \text{Propagation speed} &= 2 \times 10^7 \text{ m/s and distance is calculated as} \\ \text{distance} &= \sqrt{(S(i).xd - S(n+1).xd)^2 + (S(i).yd - S(n+1).yd)^2} \end{aligned}$$

b) **Minimum Energy:** This is another parameter that can be used to increase energy efficiency of network. If two or more cluster heads have approximate equal or comparable distance from BS, then they will be next sorted according to their energy level. The base station already has the global knowledge of location and energy levels of each sensor nodes, it will ask that cluster head to send data first which has minimum energy as it would minimise its waiting time and hence it can sleep for more time and save energy. We can consider the time taken by CHs to transmit signals to BS similar to the CPU burst time and construct following Gantt chart by taking an example for 5 CHs.

Total time for a round (tot\_time) = sum of time taken by all CHs to transmit signal to BS =  $t_1 + t_2 + t_3 + t_4 + t_5$ .

Waiting time of a cluster head is the time for which it waits for its turn to send signals.

$$Wt(i) = Wt(i-1) + t(i-1)$$

Average Waiting Time of cluster heads:

$$\text{AWT per round} = \frac{\text{Sum of waiting time of all CHs}}{\text{Total number of CHs}}$$

SLEEPING TIME of node  $i$  = Amount of time which it gets for sleeping (if it sleeps [step 3])

$$= \text{tot\_time} - (\text{waiting time of node } i + t_i)$$

Along with the cluster heads, its member nodes will also go to sleep and have the same sleeping time as their CH.

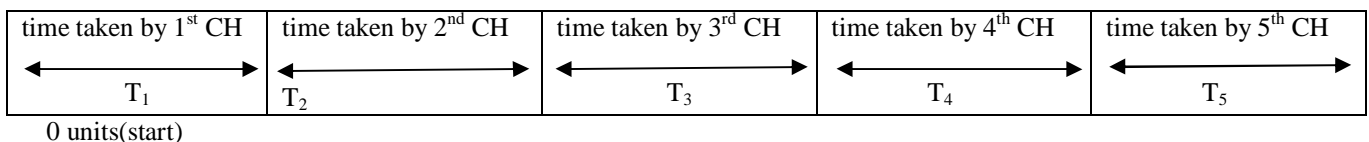


Fig2: Gantt Chart showing STF Scheduling mechanism

In Fig 2, the Gantt chart illustrates the time taken by each cluster head node to transmit signal to base station according to STF schedule. All CHs have data to send at time  $t=0$  and are waiting for their time frame. It shows the start and finish time of each signal.

STEP 3: This step is similar to that of LEACH CS[3] protocol. According to the schedule based on minimum distance and energy, the first cluster can transmit its data to base station. **Intelligent Sleeping Mechanism** of LEACH CS will be used to analyse the data and set the cluster to sleeping mode if the data sensed is insignificant that is if the aggregated value is less than the defined threshold value. The base station broadcasts packets with a flag set to 1 with the address of cluster nodes that will be switched to sleeping mode. The nodes with their address specified as destination receive the packet and stop sending data during this round. These nodes and their cluster heads then send



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their energy information to the base station finishing their role in the round. The rest of the clusters continue the round and send their data to the base station consequently.

## IV. PSEUDO CODE

```
Step 1: for i =1 to N do
Send energy_level and location to BS. // N is total no. of alive nodes
Step 2: Compute avg_energy  $E_a$ ;
Step 3: for each node i in network
if ( $E_i < E_a$ ) //  $E_i$  is the energy level of node i //
then node 'i' is not eligible to become CH for current round.
Step 4: Modify N by removing the above nodes from set of alive nodes.
Step 5: Sort remaining alive nodes (N) in descending order of Energies .
Step 6:  $p = N * 0.05$  // only 5% of nodes selected for optimal no. of CHs//
Step 7: select first p nodes out of N to be CH.
Count_CH = p; //Count_CH is no. of cluster heads//
Count_NCH = N - p; //Count_NCH is no. of non cluster heads//
Step 8: for i = 1 to Count_CH do
for i = 1 to Count_NCH do
/* non cluster heads are assigned by BS to
If  $\text{dist}(n1, n2) < \text{Min-square distance}$  appropriate CHs by finding minimum
Min-square distance =  $\text{dist}(n1, n2)$  square distance between them*/
cluster_head(n1) = n2
Step 9: NCHs send data to their CH.
Step 10: CH aggregates data. // all CHs now have to send their data to BS//
Step 11: STF Schedule
r = no. of rounds
dist_CH =  $\sqrt{(S(i).xd - (S(n+1).xd))^2 + (S(i).yd - (S(n+1).yd))^2}$ . /*distance of CHs from BS*/
s =  $2 * 10^7$  //s is speed of propagation of signals//
for every round do
time = t(i) = (dist_CH) / s
Sort((CH nodes) asc_time) asc_energy /*CHs with equal time will be further sorted
Acc. To their energies*/
avg = 0
for j = 1 to Count_CH
CH send data to BS
Wt(i) = Wt(i-1) + t(i-1) /* calculating waiting times of CHs*/
If (CH(aggregate_value) < threshold)
Send flag = 1 /* BS sends flag to CH */
Else
Send flag = 0

Step 12: if (flag = 1)
Sleep CH and cluster nodes
Sleep_time = tot_time - (waiting time of node i + t_i) /* calculating sleep time of nodes*/
Else
Continue for
```

## V. SIMULATION RESULTS

Simulation results are obtained using MATLAB. Following assumptions have been made for the simulation:

1) The nodes are immobile.

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- 2) The base station is situated at a far location from the sensor nodes in the network.
- 3) Speed of propagation of signals has been taken as  $2 \times 10^7$  m/sec which is slightly less than speed of light as we have considered disturbance through space.

Fig 3 simply shows the distribution of nodes. Fig 4 shows the average waiting time of nodes calculated in TDMA scheduling and STF scheduling. It can be clearly seen that values of average waiting time for STF scheduling are always less or equal as compared to normal TDMA scheduling. Fig 4 is obtained using 500 nodes and same comparison is done for 1000 nodes in Fig 5. Number of rounds equals 100 for every simulation. In Fig 6, we have done a comparison of sleeping time for both schedules TDMA and STF. Results shows that sleeping time of nodes in STF schedule is always more or equal as compared to TDMA schedule and hence it has been proved through simulation that decrease in waiting of sensor nodes certainly increases their sleeping time. Simulation results also showed that on average 100 rounds, the STF scheduling is approx. 20% more efficient as compared to TDMA.

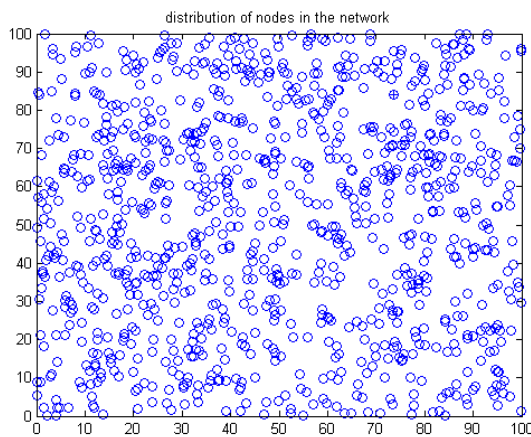


Fig 3: Distribution of nodes in the network

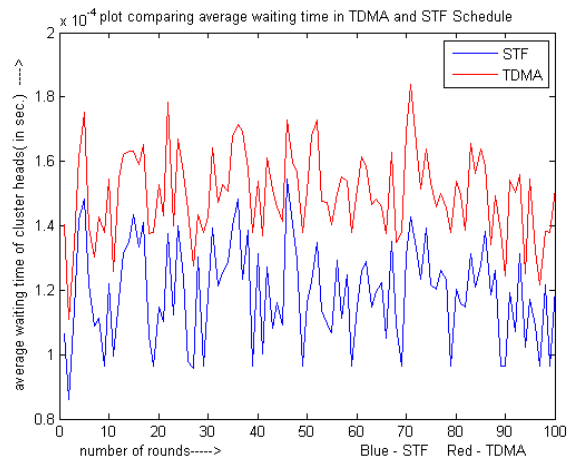


Fig 4: Average waiting time Vs no. of rounds Plot for 500 nodes

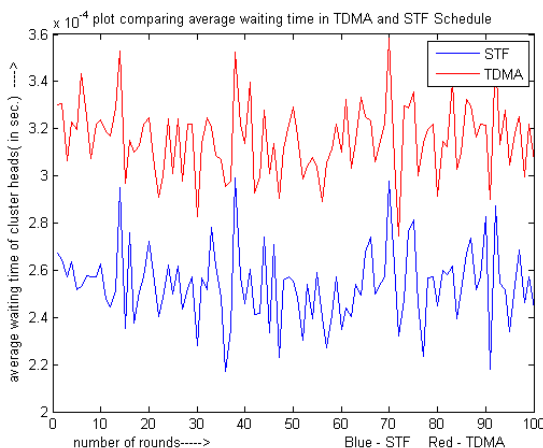
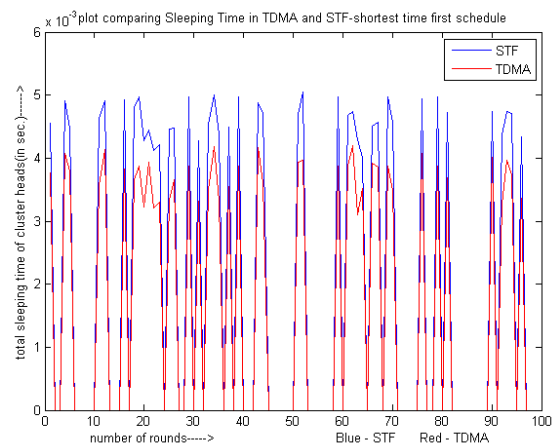


Fig 5: Average waiting time Vs no. of rounds Plot for 1000 nodesplot for 500 nodes



## VI. CONCLUSION AND FUTURE WORK

Form the simulation results, it can be concluded that our proposed STF (Shortest Time First) scheduling algorithm helps to increase the sleeping time of the sensor nodes and hence, the lifetime of sensor nodes also increases. The



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waiting time of sensors greatly affects their sleeping time and also the performance of network. The proposed algorithm STF increases the energy efficiency of wireless sensor network by allowing sensors to sleep for some more time that was wasted in waiting during transmission process that is used in a simple TDMA schedule.

In this paper, we have worked on stationary nodes in a wireless sensor network. In future we will be checking the proposed algorithm for MANET (Mobile Ad hoc network).

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