



International Journal of Innovative Research in Computer and Communication Engineering

(An ISO 3297: 2007 Certified Organization)

Website: www.ijirccce.com

Vol. 5, Issue 8, August 2017

An Algorithmic Study of Link Prediction in Social Network Sites

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ABSTRACT: In the area of social network web mining there can be different kinds of links or edges between the nodes are using. The example of them are social contacts, phone calls, hyper-reference etc. Actually the link prediction is the problem of predicting the edges and links that either don't yet exist at the given period of time t or exist, but have not been discovered are likely in near future. In our work we introduced approaches to link prediction based on various measures for proximity based approaches. For example we can consider a co-authorship network among scientists, e.g. two scientists who are close to the network will have colleagues in common, so they most likely collaborate near future. Our aim is to make betterment in creative notion precise and understandable which measures the most accurate outputs. There are three types of link prediction algorithms are used in social network link prediction. There are path based, node base and meta approaches. The node neighbour approach is based on network local features only. It is mainly focussing nodes structure only. In the case of local based measures Jaccards Coefficient, Adar/Adamic etc. The next type is path based algorithms like Sim-Rank, Katz, the rooted page rank etc. The meta approaches changes the data being passed to one the path based algorithms. The algorithms are clustering algorithms, LR-approximation algorithm, unseenbig rams etc.

KEYWORDS: Social Network, Link Prediction, Social Analysis, Similarity Metric, Learning Model.

I. INTRODUCTION

A Social Structure consists of nodes (Individuals or Organizations) and nodes are connected by different types of relationships. A set of social actors or nodes (such as individuals or organizations) and a set of the dyadic ties between these nodes constitute a social network. For example scientists in a discipline, employees in a large company, business leaders can be thought as nodes in a network and co-authors of a paper, working on a project, serve together on board can be thought as edges respectively. The idea behind Social Networks is to create opportunities to develop friendships, share information and promote business in a network. OSN like Facebook and Twitter have become important part of daily life of millions of people. The enormous growth and dynamics of these networks has led to several researches that examine the network properties i.e. structural and behavioural properties of large scale social networks. Social network is a term used to describe web-based services that allow individuals to create a public/semi-public profile within a domain such that they can communicatively connect with other users within the network [8]. Before the advent of social network, the homepages was popularly used in the late 1990s which made it possible for average internet users to share information [11].

II. LINK PREDICTION PROBLEM

Different kind of links or edges between the nodes exist in a social network. For example, social contacts, phone-calls or hyper-references. On analysis of social networks, there can be many information about the linkage between the nodes that are not discovered or unknown at a given point of time. Link Prediction is the problem of predicting links that either don't yet exist at the given time t or exist, but unknown up to this time. Given a picture of a social network (nodes and links) at time t , we need to predict accurately the links that will be added to the network during the interval from time t to a given future time $t+1$. In effect, the link prediction problem concentrates on to what extent can the evolution of a social network be modelled by using intrinsic features of the network itself? Let us consider a co-authorship network among researchers, for example, there are different reasons, outside to the network, why two researchers who have never written



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a paper together will do so in the next few years. Or, when one of the researchers changes institutions, they may come geographically very close. Such interactions are be hard to predict. But by studying the network characteristics, we can predict the possible links that are going to form. Our objective is to make this intuitive notion very exact, and to understand which measures of proximity in a graph lead to accurate predictions. Data mining provides a wide range of techniques for detecting useful knowledge from massive datasets like trends, patterns and rules [9]. The link prediction problem is also deals with the problem of getting missing links from a known network, in a number of fields. It involves prediction of additional links that are not directly visible currently, are likely to exist in a network based on observable data. It considers a static picture of the network, rather than taking network evolution and network dynamics. It also considers specific properties of the nodes in the network, rather than computing the power of prediction methods that focuses on the graph structure. Given a social network $G(V;E)$ in which an edge represents some kind of interactions between its vertices on nodes at a given time t . Suppose we have a snapshot of a social network at a given time. We choose four times $t_0 < t_{00} < t_1 < t_{01}$, and give our algorithm to predict links that are likely to be formed in the near future from the network $G[t_0; t_{00}]$. That results in predicting new links, not present in $G[t_0; t_{00}]$, that are expected to appear in the network $G[t_1; t_{01}]$. We refer to $[t_0; t_{00}]$ as the training interval and $[t_1; t_{01}]$ as the test interval [1]. The most basic approach for similarity between any pair of nodes is by taking the length of their shortest path in graph. We rank pairs of nodes in descending order of $score(x; y)$, where $score(x; y)$ is the negative of the shortest path length between x and y . We take a snapshot of a social network as training set and predict the interactions among the nodes of training set that are likely to occur in near future. The algorithms are classified as follows [2].

III. PROPOSED ALGORITHM

The link prediction algorithms based on user input as maximum path to be traversed predicts the probability of formation of link between any two nodes of the network by traversing all the paths of the network up to that certain input path length. It first traverses the path lengths of 2 i.e. the immediate neighbourhoods of the node. we can say it runs a neighbourhood algorithm on path length 2. It then produces a similarity list between every two nodes. When it traverses the graph for path length 3, it uses new paths that are made by path length 2 to get the new the length of the path of three and computes their similarity matrix and updates the similarity matrix in a cumulative way. This process continues till the graph is traversed up to the maximum path length to be traversed given by user. Let we get a path from A to B while traversing for the path length of $n-1$ with some similarity value, when we traverse the graph for path length value n , then we will check all the neighbours of B (i.e. C) to get path from A to C. We compute the similarity of each pair (A,C) and update the path list if their no direct link between A and C in the original graph. The inputs to the algorithm are the graph in terms of a list or adjacency array, the total number of nodes the graph, maximum length to be traversed, which determines how many time the algorithm will run and the path length for each specific traversal. The output of the algorithm is a similarity list containing the similarity value between every two nodes by traversing the path lengths of given maximum user input value. By observing the similarity matrix we can predict the future links. The high similar values have more probability to form links in near future and we can classify the values based on certain threshold value. The similarity values more than the threshold value are likely to form future links. This prediction can be compared with the test data to get the efficiency of the algorithm.

IV. DIFFERENT TYPES OF ALGORITHMS IN LINK PREDICTION

A. Node Neighborhood Algorithm

Node neighbourhood meaning the nodes directly connected to the two given nodes. It is simple technique which traverse only paths of length 2. For any node A it check the neighbour of A and computes their similarity with A. It considers only local features of a network, focusing mainly on the nodes structure (i.e. based on the number of common friends that two users share).

B. Common Neighbors

The Common Neighbours method provide a measure for similarity by calculating the intersection of the sets of neighbours of the nodes to predict future linkage. The Common Neighbors (CN) is defined as follows $CN(x,y) :=$



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$\Gamma(x) * \Gamma(y)$. This measurement is based on the idea that two nodes a and b have an increased probability to connect if they have a shared neighbor c. With a growing number of shared neighbours this probability grows even higher.

The weighted Common Neighbors (CNw) is defined as follows where $w(x,y)$ is the number of interactions between the nodes x and y. $CNw(x,y) := \sum_{z \in \Gamma(x) \cap \Gamma(y)} w(x,z) + w(y,z)$

C. Jaccard coefficient

Jaccard coefficient measures number of the features (neighbors) that are shared between two nodes commensurate to all features that either one of the nodes has. Jaccard coefficient [11] is a normalized variation of Common Neighbors [7]. The Jaccard coefficient is defined as follows

$$J(x,y) := \frac{|\Gamma(x) \cap \Gamma(y)|}{|\Gamma(x) \cup \Gamma(y)|}$$

This is the Common Neighbors measurement normalized by the union of the node neighborhoods.

D. Adamic/Adar

It is a measurement that compares how many attributes two nodes have in common. They rate items that are unique to a few users more heavily than items shared amongst a huge group of users. This measurement can easily be adjusted in the context of node neighborhood by looking at shared neighbors as an attribute. Therefore the sum over the shared neighbors inverse of the logarithms of their neighborhoods is proposed [3]. The Adamic/Adar is defined as follows

$$AA(x,y) := \sum_{z \in \Gamma(x) \cap \Gamma(y)} \frac{1}{\log|\Gamma(z)|}$$

The weighted Adamic/Adar (AAw) is defined as follows where $w(x,y)$ is the number of interactions between the nodes x and y [4].

$$AAw(x,y) := \sum_{z \in \Gamma(x) \cap \Gamma(y)} \frac{w(x,z) + w(y,z)}{\log|\Gamma(z)|}$$

E. Preferential Attachment

Preferential Attachment is based on the hypothesis that a node x will get new neighbors faster than a node y given y has less neighbors than x. So the probability that a node will form a new link varies with number of its present neighbors. The likelihood of two nodes being connected by an edge based on preferential attachment is measured by multiplying the number of their neighbors [5]. The Preferential Attachment is defined as follows

$$PA(x,y) := \Gamma(x) \cdot \Gamma(y)$$

The weighted Preferential Attachment (PAw) is defined as follows where $w(x,y)$ is the number of interactions between the nodes x and y: $PAw(x,y) := \sum_{x \in \Gamma(x)} w(x,x) \cdot \sum_{y \in \Gamma(y)} w(y,y)$

V. PATH BASED ALGORITHMS

Some measurements of link prediction take all paths between two nodes in consideration. The computation of graphs that take the entire graph in consideration is by nature much more complex than node neighborhood algorithms.

A. Katz

A measurement that takes all paths between two nodes in consideration while rating short paths more heavily. The measurement exponentially reduce the contribution of a path to the measure in order to give less weightage longer paths. Therefore it uses a factor of β^l where l is the path length.

The Katz is defined as follows

$$K(x,y) := \sum_{l=1}^{\infty} \sum_{paths \langle l \rangle_{x,y}} \beta^l$$

$$\beta^l \cdot |\text{paths} \langle l \rangle_{x,y}|$$

where $\text{paths} \langle l \rangle_{x,y}$ the set of all paths from source x to destination y that have the path length l.

B. Unweighted : $\text{paths} \langle l \rangle_{x,y} = 1$, if x and y have collaborated and 0 otherwise • **Weighted :** $\text{paths} \langle l \rangle_{x,y}$ is the number of times that x and y have collaborated

The β can be used to control how much the length of the paths should be considered. A very small β concludes to an algorithm where paths of length three or more are taken much less into account and therefore the algorithm converges node neighborhood algorithms. It has roughly cubic complexity as it requires matrix inversion [6].



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C. SimRank

If two nodes are referenced by more similar objects, then the two nodes have large similarity value. Every object obviously has a similarity score of 1 to itself. Node x and node y are then similar to the degree they are joined to similar neighbours [7]. The SimRank is defined as follows[12]

$$S(x,x) := 1 \quad S(x,y) := \gamma \cdot Pa\Gamma(x)Pb\Gamma(y) S(x,y) \Gamma(x).\Gamma(y)$$

γ is a constant with $\gamma \in [0,1]$. The constant can be thought of as a confidence level. If you consider a situation in which a and b are both neighbours to c , than obviously the similarity of c to itself is 1, but we do not want to conclude that $s(a,b) = s(c,c) = 1$. Instead we let $s(a,b) = \gamma * s(x,x)$ because we are not as confident about the similarity of a and b as we are about $s(x,x) = 1$.

D. Hitting Time and Commute Time

Starting from a node x a random walk on a given graph moves iteratively over the graph while choosing the next node each step at random. The expected number of steps to get from x to y via a random walk is defined as the Hitting Time $H(x,y)$. A short hitting time implies node similarity and therefore a heightened chance of future linking. The commute time $C(x,y)$ is a variant of Hitting time which is useful for undirected graphs, because the hitting time is not symmetric. Therefore it is defined as follows:

$$C(x,y) := H(x,y) + H(y,x)$$

The commute time can have high variance, hence, prediction by this feature can be poor. If z is a node with high stationary probability far o $\cdot x$ and y , then a random walker would probably reach the neighbourhood of z . To avoid that we can use reset the random walker to x with a fixed probability of α .

two normalized versions Hitting Time normalized (H_n) and Commute-Time normalized (C_n) are defined where π_x is the stationary probability of x to safeguard it against vertices with a very high π :

$$H_n(x,y) := -H(x,y).\pi_y \quad C_n(x,y) := -(H(x,y).\pi_y + H(y,x).\pi_x)$$

VI. THE PROPOSED ALGORITHM

The link prediction algorithms based on user input as maximum path to be traversed predicts the probability of formation of link between any two nodes of the network by traversing all the paths of the network up to that certain input path length. It first traverses the path lengths of 2 i.e. the immediate neighbourhoods of the node. we can say it runs a neighbourhood algorithm on path length 2. It then produces a similarity list between every two nodes. When it traverses the graph for path length 3, it uses new paths that are made by path length 2 to get the new path lengths of 3 and computes their similarity matrix and updates the similarity matrix in a cumulative way. This process continues till the graph is traversed up to the maximum path length to be traversed given by user. Let we get a path from A to B while traversing for the path length of $n-1$ with some similarity value, when we traverse the graph for path length value n , then we will check all the neighbors of B (i.e. C) to get path from A to C . We compute the similarity of each pair ($A;C$) and update the path list if their no direct link between A and C in the original graph.

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For each path length, we have to follow many steps:

- Calculate that path list with respect to the list with previous input path value
- Update the current adjacency list for the entries having nonzero path value
- Calculate the similarity measure with respect to the corresponding path list
- Update the similarity list by adding the new similarity measures to the list
- Increment the path length

This iteration stops when current path length exceeds the maximum value of path length to be traversed.

A. Algorithm Parameters



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a) Input Parameters

- A : adjacency matrix of undirected and unweighted graph
- n : total number of nodes of the graph
- l : max length of path to be explored in G
- m : the length of a path for current iteration

b) Output Parameters

- sim(i,j) : Similarity measure between nodes i and j

VII. ALGORITHM

Algorithm 1: MAIN FUNCTION

```
for m ← 2 to n do
  cpath(A,n,prev,or) sim ← simi(sim,path,n,m);
  path ← 0 ;
end
```

The main Program iteratively calls for checking the new collaboration between any two nodes for a specific path length through cpath function and computes the similarity between the new collaborations with exactly m path length and updates the similarity measures through sim function.

```
For i ← 1 to N do
  for j ← 1 to N do
    if i < j then
      if or (i,j) ≠ j then
        for k ← 1 to N do
          if prev (i,k) ≠ 0 then
            if A (i,k) ≠ 0 and or (k,j) ≠ 0 then
              path (i,j) ← path (i,j)+prev (i,k) * or (k,j) ;
            end
          end
        end
        path (j,i) ← path (i,j) ;
      end
    end
  end
  prev ← path;
  for i ← 1 to N do
    for j ← 1 to N do
      if path (i,j) ≠ 0 then
        A (j,i) ← 1 ;
      end
    end
  end
  return path and A;
```

The *cpath* function first checks whether there is path from any two nodes of m path length. It checks it by merging the new paths generated while traversing the previous path length (m-1) and their neighbours in the original graph. So pathmatrix contains new collaborations of path length exactly m.

Algorithm 2: FUNCTION SIMI

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```
for i ← 1 to n do
  for j ← 1 to n do
    lower ← 1;
    for k ← 1 to m do
      lower ← lower * (n-k);
    end
    sim (i,j) ← ( 1/(m-1)* sim (i,j) ) / lower;
  end end return sim;
```

The simi function finds the similarity measure for every pair of nodes which have path length exactly m. It then cumulates the similarity values till the path length of l for every two nodes. The pair of nodes having higher value of similarity are more likely to form link in the near future.

VIII. OUR PROPOSED METHOD IMPLEMENTATION

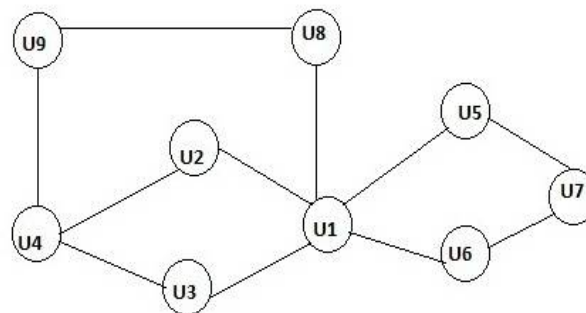


Figure 1 Link Prediction Analysis

Our proposed algorithm performs better than global approaches as it is based on user input bounded path traversal. It predicts the links by traversing to a certain path length given by the user. In global approaches, the time and space complexity are high due to consideration of all paths in the network. As our algorithm only traverse up to certain path length, so its complexity is low with comparison to global approaches. It also outperforms the node neighbourhood algorithms as it traverses more path lengths than the node neighbourhood algorithms. It considers the network characteristics around the target node to predict its future links. We are implementing our proposed method on a small network. The input to the algorithm is the network below and the output is a similarity matrix. It has been discussed that with path length of 3, U4 has more probability to form link with U1 with comparison to U7. When we consider only path lengths of 2, then U4 and U7 have the equal probability of forming links with U1 as they are connected with two different ways of path length 2. But when we consider the path length of 3, then U4 has more probability to form a link than U7. Because U4 is connected with three different paths with U1. If we have followed node neighbourhood techniques, then we will get equal probability of U4 and U7 getting connected with U1. But through our algorithm with max traversal length of 3, we can conclude that U4 has more chance to get connected with U1 than U7. So our algorithm performs better than node neighbourhood algorithms in terms of efficiency of result. It also performs better than global approaches in terms of time and space complexity.

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TABLE I

Algorithm	Condmnat	Disease-g
Page Rank	370 sec	530 sec
katz	430 sec	620 sec
Node neighbor	72 sec	230 sec
Proposed	Proposed	350 sec

Figure 3. Comparison Table

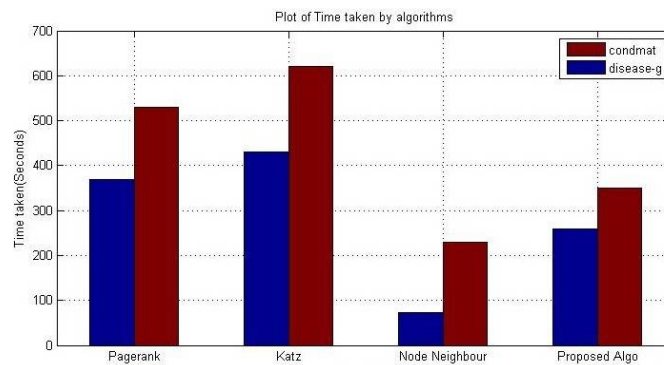


Figure4: Link Prediction Analysis

The algorithm is implemented using maximum length to be traversed as 4. This result shows that the similarity value of the node pair U1 and U4 is maximum

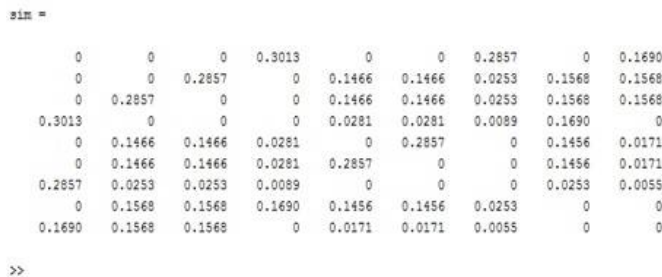


Figure 2: Similarity matrix

and more than that of the pair U1 and U7. So, there is more probability of formation of link between U1 and U4. Hence our proposed method works correctly.



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IX. CONCLUSION

Link Prediction is the method to predict the possible future interactions among the nodes in the near future. Our algorithm uses both global and local characteristics of the network to predict the links. Global approaches has the time constraint as they traverse all paths of network to predict the links and local approaches are less efficient as they consider only local features of the node. Our approach is compared with all the approaches and it provides efficient and accurate friend suggestions in a less interval of time.

X.FUTURE RESEARCH OPERTUNITIES

Link Prediction based on other features like photo, video tagging can be used for better prediction. As many features as we consider simultaneously, the prediction will be better because it gives information about many ways peoples may by connected. We can consider the positive as well as negative links in a network. If positive weight is for support, then negative weight should be for opposing it. As network is always dynamic, so we can consider network dynamics into consideration.

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BIOGRAPHY

B.Suryakumar received M.Sc and MCA degrees in Computer Science from Annamalai University, Tamilnadu. Currently he is doing PhD in Computer Science at Bharathiar University, Coimbatore. His research interest lies in the area of Data Mining ,Networking and Data security.

Dr.E.Rama Devi received PhD degree in Computer Science from Mother Teresa Womens University, Kodaikanal. Currently she is an Associate Professor in Computer science at NGM College, Pollachi, India. She has got 14 years of research experience and has more than 19 years of teaching experience. Her researchinterest includes areas like Data Mining, Knowledge base System, Intelligent and Control System and Fuzzy Logic. She has presented various papers in national and International Conferences and published 10 research papers on refereed journals.