



Area–Delay–Power Efficient Carry-Select Adder

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ABSTRACT: Carry Select Adder (CSLA) is faster than any other adders used in many data-processing processors to perform arithmetic functions speedily. In adder design carry generation is the critical path. To reduce the power consumption of data path we need to reduce Area and Delay of the adder. The proposed design is implemented without using multiplexer and RCA structure with $C_{in}=1$. Instead of multiplexer and RCA $C_{in}=1$ structure here we used simple combinational circuit which consists AND and XOR and OR gates. In the proposed scheme, the carry select (CS) operation is scheduled before the calculation of *final-sum*, which is different from the conventional approach. Bit patterns generated by two carry words and fixed carry c_{in} is used by CS unit. The proposed CSLA design involves significantly less area and delay than the recently proposed BEC-based CSLA. A theoretical estimate shows that the proposed SQR- CSLA involves nearly 35% less area–delay–product (ADP) than the BEC-based SQR- CSLA, which is best among the existing SQR- CSLA designs, on average, for different bit-widths. Carry Select Adder (CSLA) is one of the fastest adders used in many data-processing processors to perform fast arithmetic functions. The proposed design is synthesized and simulated in Xilinx ISE design suite 14.2 and is implemented on Spartan 6 XC6SLX16 CSG324 FPGA device.

KEYWORDS: Adder, arithmetic unit, SQR, CSLA, FPGA, BEC.

I. INTRODUCTION

Low power, area-efficient, and high-performance VLSI systems are increasingly used in electronic applications such as portable mobile devices, multi standard wireless receivers, and biomedical instrumentation [1], [2]. An adder is the main component of an arithmetic unit. A complex digital signal processing (DSP) system involves several adders. Applications where these are used are multipliers, DSP to execute various algorithms like FFT, FIR and ITR. In microprocessors, millions of instructions per second are performed. An efficient adder design essentially improves the performance of a complex (CS) methods have been suggested to reduce the CPD of adders. A conventional carry select adder (CSLA) is consist of an RCA–RCA configuration it uses multiple pairs of Ripple Carry Adders(RCA) to generate partial sum and carry by considering carry input $C_{in}=0$ and $C_{in}=1$, then the final sum and carry are selected by the multiplexers (mux)[3]. A ripple carry adder (RCA) uses a simple design, but carry propagation delay (CPD) is the main concern in this adder. A conventional CSLA has less CPD than an RCA, but the design is not attractive since it uses a dual RCA. Few attempts have been made to avoid dual use of RCA in CSLA design. Kim and Kim [4] used one RCA and one add-one circuit instead of two RCAs, where the add-one circuit is implemented using a multiplexer (MUX). He et al. [5] proposed a square-root (SQR)-CSLA to implement large bit-width adders with less delay. In a SQR CSLA, CSLAs with increasing size are connected in a cascading structure. The main objective of SQR-CSLA design is to provide a parallel path for carry propagation that helps to reduce the overall adder delay. Ramkumaret.al proposed BEC technique by using (n+1)-bit BEC instead of n-bit RCA with $C_{in}=1$ which reduces the area power-delay product and total power consumption with slight increase in delay [6].

The BEC- based CSLA involves less area and low power than the conventional CSLA, but it has marginally higher delay. A CSLA based on common Boolean logic (CBL) is also proposed in [7] and [8]. The CBL-based CSLA of [7] involves significantly less logic resource than the conventional CSLA but it has longer CPD, which is almost equal to that of the RCA. To overcome this problem, a SQR-CSLA based on CBL was proposed in [8]. However, the CBL-based SQR- CSLA design of [8] requires more logic resource and delay than the largely depends on availability of

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redundant operations in the formulation, whereas adder delay mainly depends on data dependence. In the existing designs, logic is optimized without giving any consideration to the data dependence.

Based on the proposed logic formulation, we have derived an efficient logic design for CSLA. Due to optimized logic units, the proposed CSLA involves significantly less ADP than the existing CSLAs. We have shown that the Sqrt-CSLA using the proposed CSLA design involves nearly 32% less ADP and consumes 33% less energy than that of the corresponding Sqrt-CSLA.

II. REGULAR/CONVENTIONAL Sqrt CSLA

A. Architecture of Regular Sqrt CSLA

The architecture of 16-bit regular Sqrt CSLA [3] is shown in figure 1. It is divided into five groups with different size of RCA. Two RCAs are required to add two n bit numbers in CSLA. The partial sum and carry are generated by one RCA with $C_{in}=0$ and another RCA with $C_{in}=1$. The mux selects the final sum and carry. The select line of mux is carry out of previous group. In Sqrt CSLA, if the number of FA increases, the complexity of the design increases and so does the power consumption.

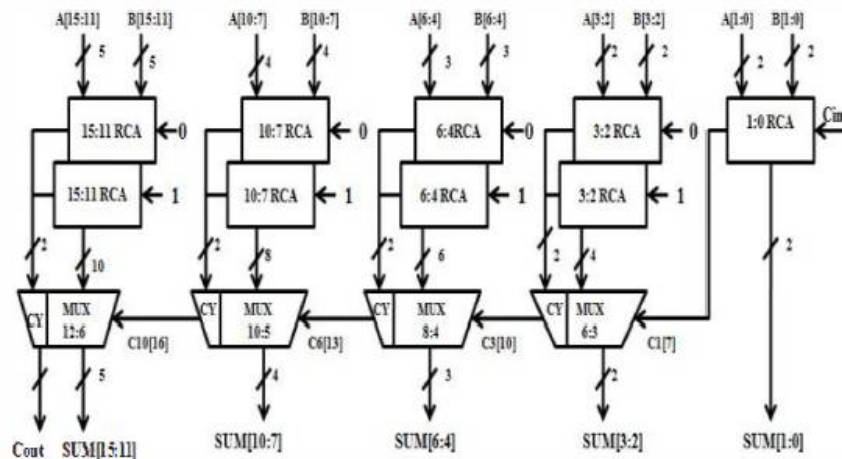


Fig.1: Regular 16 bit Sqrt CSLA

B. Area Calculation of Regular Sqrt CSLA

The group 3 of regular 16-bit Sqrt CSLA is elaborated in figure 2. It has two sets of 3-bit RCA. The 3-bit RCA with $C_{in}=0$ consists of two FAs and one HA whereas 3-bit RCA with $C_{in}=1$ consists of three FAs. Four 2:1 mux select the final sum and carry. RCA with $C_{in}=0$ and RCA with $C_{in}=1$ are the two inputs of mux. For group 3, 87 AOI gates are used as given below.

$$FA = 65 (5 \times 13)$$

$$HA = 6 (1 \times 6)$$

$$MUX = 16 (4 \times 4)$$

.....

$$\text{Total Gate} = 87$$

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Vol. 5, Issue 5, May 2017

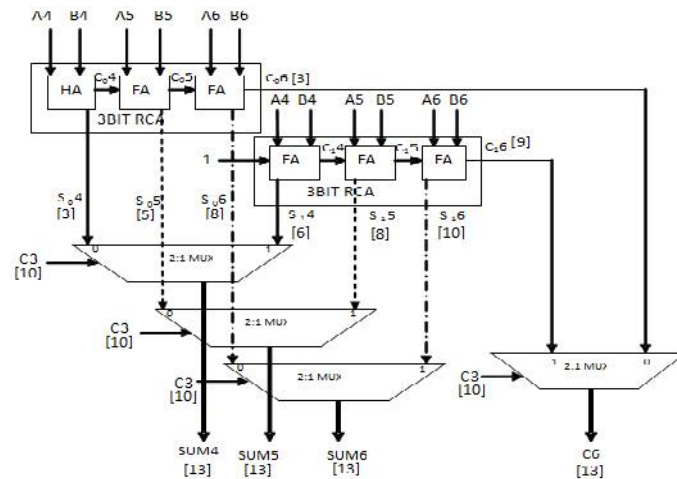


Fig. 2: Group-3 of 16 bit Regular Sqrt CSLA

III. MODIFIED Sqrt CSLA USING BEC

The area and power problem of regular Sqrt CSLA was overcome by BEC [6] technique. The truth table of 4-bit BEC is given in Table 1. The Boolean expressions (1) to (5) are derived from the truth table. Figure 3 illustrates the structure of a 4-bit BEC. Numbers of logic sources are less in BEC thus reducing area and power consumption than that of regular Sqrt CSLA. BEC is combination of NOT (\sim), AND (\cdot), XOR (\oplus) gates. P3, P2, P1, P0 are the inputs and Q3, Q2, Q1, Q0 are the outputs of BEC while CY is the final carry out.

Table 1: Truth Table of 4-Bit BEC

INPUT				OUTPUT				
P3	P2	P1	P0	CY	Q3	Q2	Q1	Q0
0	0	0	0	0	0	0	0	1
0	0	0	1	0	0	0	1	0
1	1	1	1	1	0	0	0	0
.
1	1	1	0	0	1	1	1	1

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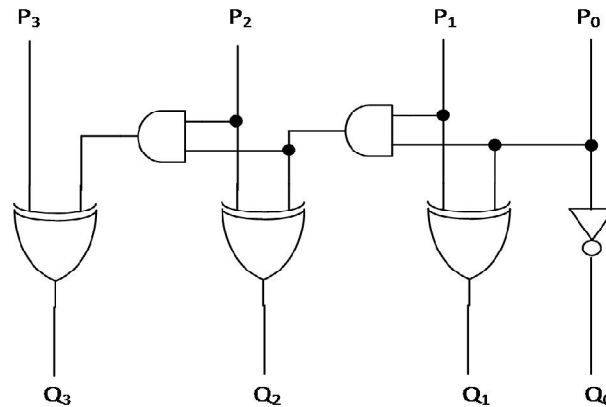


Fig. 3. 4-bit BEC

Boolean expression of 4 bit BEC Converter,

$$\begin{aligned}
 Q_0 &= \sim P_0 & (1) \\
 Q_1 &= P_1 \oplus P_0 & (2) \\
 Q_2 &= P_2 \oplus (P_1 \cdot P_0) & (3) \\
 Q_3 &= P_3 \oplus (P_2 \cdot P_1 \cdot P_0) & (4) \\
 CY &= P_3 \cdot P_2 \cdot P_1 \cdot P_0 & (5)
 \end{aligned}$$

A. Architecture of SQR T CSLA Using BEC

The architecture of 16-bit modified SQR T CSLA is shown in figure 4. The five groups consist of variable sized RCAs. CSLA consists of one RCA with $C_{in}=0$, BEC and mux. BEC logic is used instead of the RCA with $C_{in}=1$. A (n+1)-bit BEC is used to replace an n-bit RCA. In this design, the mux inputs are RCA (with $C_{in}=0$) output and BEC output. By comparing with regular 16 bit SQR T CSLA the area and power are reduced with a delay penalty.

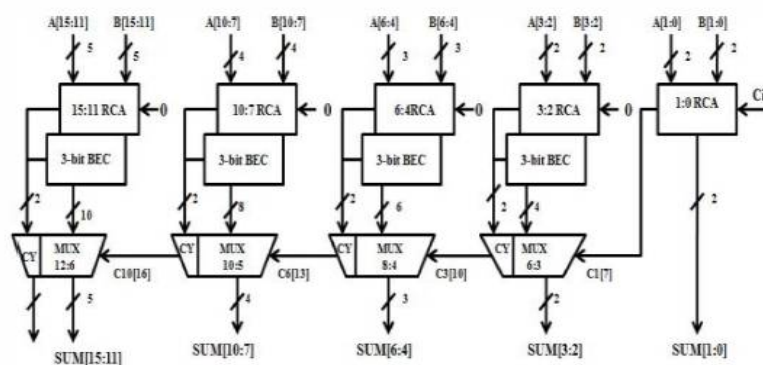


Fig. 4: Modified 16 bit SQR T CSLA

B. Area calculation of SQR T CSLA using BEC

The group 3 of modified 16-bit SQR T CSLA using BEC is elaborated in figure 5. The group 3 has one 3-bit RCA and one 4-bit BEC instead of another RCA. The 3-bit RCA with $C_{in}=0$ consists of two FA and one HA whereas a 4-bit BEC consists of three XOR gates, one NOT gate and two AND gates. Four 2:1 muxes are used to select the final sum and carry. The two inputs of mux are RCA with $C_{in}=0$ and BEC output. For group 3, 66 AOI gates are used as given below.

$$FA = 26 (2 * 13)$$

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HA = 6 (1*6)
XOR=15 (3*5)
NOT=1 (1*1)
AND=2 (2*1)
MUX=16 (4*4)
.....
Total Gate = 66

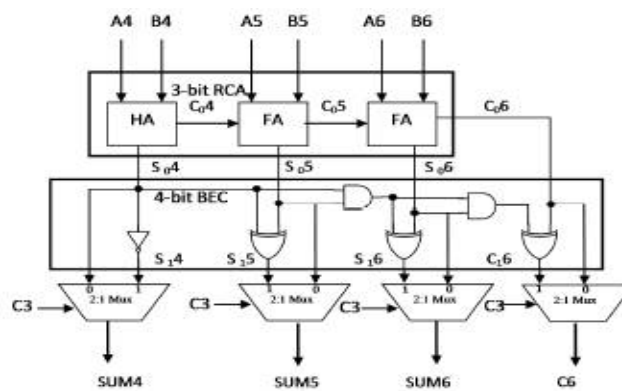


Fig. 5: Group3 of 16 bit Modified SQRCSLA

IV. PROPOSED ADDER DESIGN

As shown in figure 7 the proposed CSLA consist of a 4 basic units, 1]one HSG(Half Sum Generator)unit, 2]one FSG(Full sum generator) unit, 3]CG(Carry Generator)unit, 4]one CS(Carry Select)unit. The CG unit consist of two CGs (CG0 and CG1). One is corresponding to carry input '1' and another to carry input '0'.

A. Architecture of Proposed CSLA

The architecture 16 bit proposed CSLA is as shown in figure 6. It is consist of one RCA and 4 proposed CSLA. The HSG receives two n -bit operands (A and B) and generate *half-sum* word s_0 and *half-carry* word c_0 of width n bits each. Both CG0 and CG1 receive s_0 and c_0 from the HSG unit and generate two n -bit full-carry words c_{01} and c_{11} corresponding to input-carry '0' and '1', respectively. The CS unit selects one final carry word from the two carry words available at its input line using the control signal c_{in} . It selects c_{01} when $c_{in} = 0$; otherwise, it selects c_{11} .

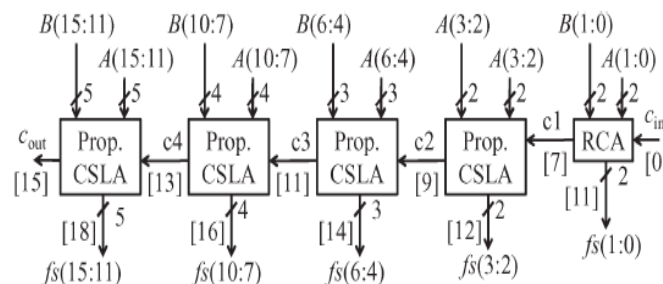


Fig. 6: Proposed SQRCSLA for $n = 16$. All intermediate and output signals are labelled with delay (shown in square brackets).

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B. Area calculation of Proposed CSLA

The group of proposed n-bit Sqrt CSLA is elaborated in figure 7. The n bit proposed Sqrt CSLA has 4 basic units and each unit is consist of different sets of AND, OR and XOR gates. HSG unit is combination of AND and XOR gate. CG0 and CG1 units includes AND and OR gate combinations unit consist of AND and OR gate in series combination and FSG unit consist of XOR gate. For group 3 64 AOI gates are used as given below

HA=1

FA=1

AND=22

OR=12(3*4)

XOR=28

.....

Total Gate=64

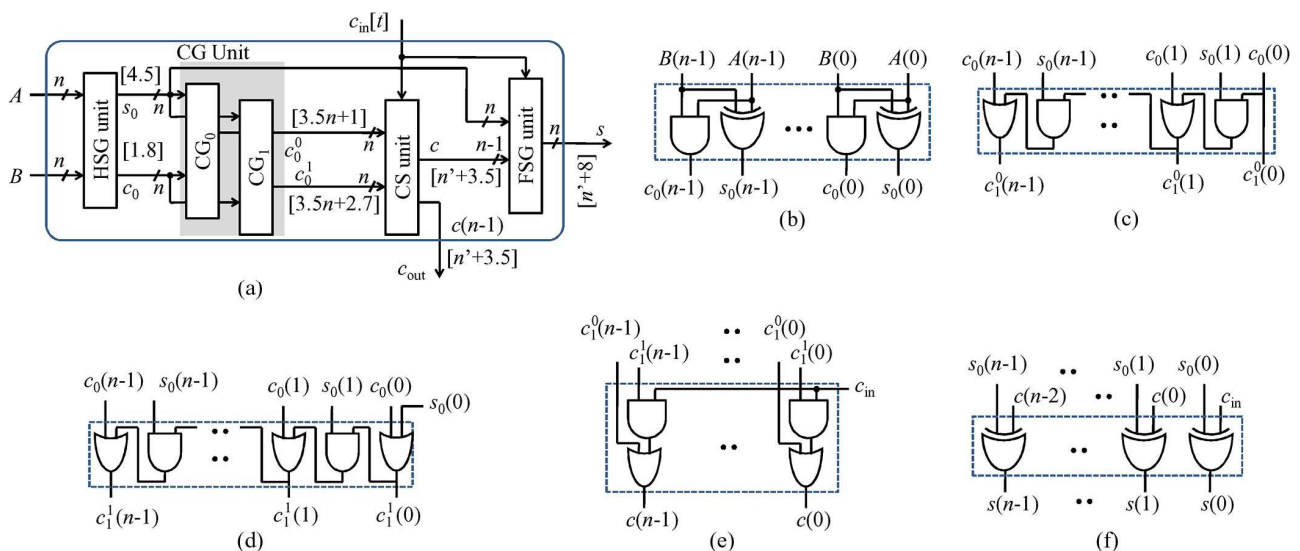


Fig. 7: (a) Proposed CS adder design, where n is the input operand bit-width, and $[*]$ represents delay (in the unit of inverter delay), $n = \max(t, 3.5n + 2.7)$. (b) Gate-level design of the HSG. (c) Gate-level optimized design of (CG0) for input-carry = 0. (d) Gate-level optimized design of (CG1) for input-carry = 1. (e) Gate-level design of the CS unit. (f) Gate-level design of the final-sum generation (FSG) unit.

V. TOOLS USED

This work evaluates the performance of the proposed designs in terms of area, power by hand with logical effort and through Xilinx ISE 14.2(Verilog HDL).This coding is implemented in FPGA (Spartan 6).The Xilinx 14.2 (Verilog HDL) provide all simulation results. Area, Delay requirement of all type of carry select adders along with proposed carry select adder, so that it's easy to compare and got conclusion that which system is good in all circumstances.

VI. SYNTHESIS RESULTS

We have coded the Sqrt-CSLA in VHDL using the pro-posed CSLA design and the existing CSLA designs of [6] and [7] for bit-widths 16, 32, and 64. This work evaluates the performance of the proposed designs in terms of area, power by hand with logical effort and through Xilinx ISE 14.2(Verilog HDL) and this will be implemented in FPGA (Spartan 6). As shown in Table 4, the proposed Sqrt-CSLA involves significantly less area and less delay than the existing designs. We can find from Fig. 7 that the proposed Sqrt. Table 2. Shows CSA 16 bit exist

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report and table 3. Shows CSA 16 bit proposed report which is run on Xilinx ISE 14.2(Verilog HDL)

Table 2: CSLA 16 Bit Exist Report Table 3:CSLA 16 Bit Proposed Report

rc Project Status (03/23/2017 - 18:55:19)			
Project File:	rc.vise	Parser Errors:	No Errors
Module Name:	rc	Implementation State:	Placed and Routed
Target Device:	xc6slx16-2zsg324	Errors:	No Errors
Product Version:	ISE 14.2	Warnings:	No Warnings
Design Goal:	Balanced	Routing Results:	All Signals Completely Routed
Design Strategy:	Xilinx Default (Unlocked)	Timing Constraints:	
Environment:	System Settings	Final Timing Score:	0 (Timing Report)

Device Utilization Summary				
Slice Logic Utilization	Used	Available	Utilization	Note(s)
Number of Slice Registers	0	18,224	0%	
Number of Slice LUTs	18	9,112	1%	
Number used as logic	18	9,112	1%	
Number using O6 output only	10			
Number using O5 output only	0			
Number using O5 and O6	8			
Number used as ROM	0			
Number used as Memory	0	2,176	0%	
Number of occupied Slices	11	2,278	1%	
Number of MIO I/Os used	0	4,556	0%	

prop Project Status (03/24/2017 - 15:01:28)			
Project File:	prop.vise	Parser Errors:	No Errors
Module Name:	prop	Implementation State:	Placed and Routed
Target Device:	xc7a100t-2zsg324	Errors:	No Errors
Product Version:	ISE 14.2	Warnings:	57 Warnings (57 new)
Design Goal:	Balanced	Routing Results:	All Signals Completely Routed
Design Strategy:	Xilinx Default (Unlocked)	Timing Constraints:	
Environment:	System Settings	Final Timing Score:	0 (Timing Report)

Device Utilization Summary				
Slice Logic Utilization	Used	Available	Utilization	Note(s)
Number of Slice Registers	0	126,800	0%	
Number of Slice LUTs	16	63,400	1%	
Number used as logic	16	63,400	1%	
Number using O6 output only	7			
Number using O5 output only	0			
Number using O5 and O6	9			
Number used as ROM	0			
Number used as Memory	0	19,000	0%	
Number used exclusively as route-thrus	0			
Number of occupied Slices	8	15,850	1%	

VII. SIMULATION RESULTS

Simulation result for 16 bit regular CSLA and 16 bit proposed CSLA is as shown in figure 8 and figure 9 respectively. Waveforms for input, output and carry signal are generated.

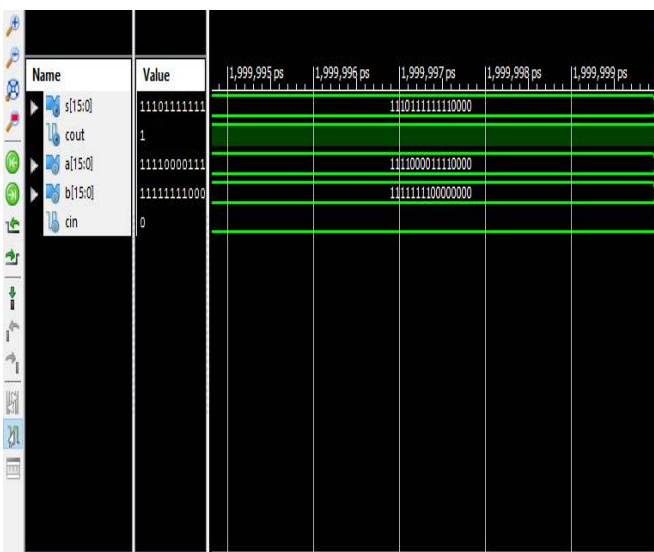


Fig.8 simulation result of 16 bit exist CSLA



Fig.9 simulation result of 16 bit proposed CSLA



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VIII. EVALUATION RESULTS

Table 4: Comparison of Regular and Modified Carry Select Adder

WORD SIZE	ADDER	DELAY(ns)	AREA(NO.OFLUTs)
16 bit	Regular CSLA	6.45	18
	Modified CSLA	5.688	16
32 bit	Regular CSLA	10.73	39
	Modified CSLA	8.859	32

As shown in Table 4, the proposed SQRT-CSLA involves significantly less area and less delay and consumes less power than the existing design which is run on Xilinx ISE 14.2 (Verilog HDL). This will be implemented in Spartan 6 kit.

IX. CONCLUSION AND FUTURE WORK

A simple approach is proposed in this paper to reduce the area and power of SQRT CSLA architecture. The reduced number of gates of this work offers the great advantage in the reduction of area and also the power. The modified CSLA architecture is therefore, low area, low power, simple and efficient for VLSI hardware implementation. I have analyzed the logic operations involved in the conventional and BEC-based CSLAs to study the data dependence and to identify redundant logic operations. I have eliminated all the redundant logic operations of the conventional CSLA and proposed a new logic formulation for the CSLA. In the proposed scheme, the CS operation is scheduled before the calculation of final-sum, which is different from the conventional approach. Carry words corresponding to input-carry '0' and '1' generated by the CSLA based on the proposed scheme follow a specific bit pattern, which is used for logic optimization of the CS unit. The proposed CSLA design involves significantly less area and delay than the recently proposed BEC-based CSLA. Due to the small carry-output delay, the proposed CSLA design is a good candidate for the SQRT adder. The ASIC synthesis result shows that the existing BEC-based SQRT-CSLA design involves 48% more ADP and consumes 50% more energy than the proposed SQRT-CSLA, on average, for different bit-widths.

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