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Security in Searching Shared and Encrypted Data

Chaudhari Ashwini , Darade Priyanka, Dhumal Sujata, Digraje Anuradha

Students, Dept. of CS, Pad.. Dr.D.Y.Patil Institute of Engineering and Technology, Pimpri, Pune, Savitribai

Phule Pune University, Pune, India

ABSTRACT: Cloud computing is a capable, evolving Internet computing of this era. It presents the users with a secure storage for storing the documents online wherein the users can take the benefit of freedom to access it remotely avoiding the usage of the data storage services. When it comes to cloud data security, new technique is required. Protecting data in the cloud can be similar to caring data within a traditional data center or enhanced data center like cloud. Authentication and uniqueness, accesscontrol, encryption, prote cted deletion, is a numerous authentication encryption term. For encryption-based data access control for cloud, in which it shows that the mechanism of security is dealing with revocation could achieve by the different security techniques. It demonstrates that a encryption method in cipher text updating key for authentication for trusted user, so a security susceptibility appears. A revoked user can still decrypt new cipher texts for that user want to requested for the new secret keys to access data.

KEYWORDS: Encryption, Trapdoor, Index, Searchable Encryption

I.INTRODUCTION

Cloud, also known as 'on-demand computing', is a class of Internet-based computing, where shared resources, facts and information are handle to computers and other devices on claim. Data security is the most important issues in cloud . To achieve high flexibility and to strong authentication for multiple data owners are outsourcing their data provides to private cloud. The data encryption reduces the data utilization. Consider large numbers of documents are outsourced on cloud by large number of cloud handler. It is mandatory for the search service to provide results similarity ranking to provide the exact results. Retrieving of all the data files having queried keyword will not be affordable in pay as peruse cloud model. The search techniques are shows that to solve the problem of multiuser data access over encrypted data using trusted third party in cloud . User will encrypt their data nearby. Before encrypting data, the index will be created. Trusted other party will use all these indexes to find data similar to the look for query of user. Using all the finding results, cloud server will send encrypted document to the user.

Data encryption makes effectual data consumption a very difficult task given that there could be a big amount of outsourced data files. In the Cloud, data owners may divide their outsourced data with a big number of authenticated users, who may want to only retrieve certain specific data files they are paying attention in through a given period. This keyword find technique allows users to selectively retrieve files of notice and has been widely useful inoriginal look for scenarios. The data encryption technique, which unauthorized user's ability to perform keyword look for and it demands the protection of data privacy, makes the traditional plaintext examine methods fail for encrypted cloud data.

II.RELATED WORK

In [1] authors initiated the investigation of searchable encryption schemes in the symmetric setting .Thissetting typically assumes one user and one server, where the user can generate searchable contents and stores them at the server, and later delegate the server to search on her behalf. To address our problem with any of these schemes, one needs to independently generate a key to protect each document and then shares the key with the users who will be authorized to search this document. As a result, the number of trapdoors required to search for a keyword will equal to the number of documents (or, indexes), so that the solution obviously does not satisfyour scalability criteria.In[2]authors proposed the concept of multi-user searchable encryption schemes, where a user can authorize multiple other users to search her encrypted data. However, the proposed primitive does not take into account the fact



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that the same user may also be authorized to search other users data and the corresponding security issues. As a result, the primitive from [8] offers a solution for a much more simplified problem than ours, and it seems not trivial to construct a scalable solution for our problem based on their scheme. In [4] authors introduced into the system, to manage multiple users' search capabilities (e.g. enable them to search each other's data). In this extension, the user manager needs to be fully trusted since it is capable of submitting search queries and decrypting encrypted data. This conflicts with our security criteria.

III.LITERATURE SURVEY

1)Selective Document Retrieval Scheme[1] SDR the scheme is secure in security model and can be adapted to support many useful search features, including collecting search results, associate conjunctive keyword search queries, advanced keyword search, search with keyword existence frequency, and search based on central product. These are the parameter are define the SDR parameter: Keygen, Build Index, Trapdoor, Search Index, Retrieve. Keygen(s): Run by a client, this algorithm takes security parameters as input, and outputs a secret key K. It may also produce some other public parameters such as a predicate set F. Build Index(K; d): Run by the client, this algorithm takes the key K and a document d 2 D as input, and outputs an index Id which encodes u(d) (i.e. all keywords from the document d).Trapdoor(K;f): Run by the client, this algorithm takes the key K and a predicate f 2 F as input, and outputs a trapdoor Tf . SearchIndex(Tf ; Id): Run by the server, this algorithm takes a trapdoor Tf and an index Id as input and returns an encrypted result to the client, where Rd implies whether u(d) satisfies the predicate f or not. Retrieve: Run between the client and the server, the client takes the secret key K and the encrypted search results as input and the server takes the encrypted database DB as input. At the establishment of the protocol, the client first decrypts and Decides which documents to retrieve, and at the end of the protocol the Client retrieves the documents user wants. This scheme provide flexible services to the trusted users but it is not efficient to provide Multi-User Authentication Services.[1]

2)Secure Inverted Index Scheme[2]An inverted index is a data structure loading words or numbers in a file along with its location. The determination of an inverted index is to progress the time of full text searches. An inverted index holds an index of keywords which stores a different list of terms finding the collection and, for individually term, a posting the updating list of documents that hold the keyword. An inverted index improves search effectiveness which is required for very large text files. An inverted index consists of a distinct terms and a posting list which stores the IDs of the documents that hold that term. In count to an ID, each posting holds list element gives the number of rates occurrences of that term in the document.



Structure of an Inverted Index

It provide good retrieval performance as well as better security for indexes. The major drawback of this process is that, It Track unnecessary network traffic for retrieval of data. [2]

3) Password-based Group Key Exchange in a Constant Number of Rounds.[3]In the password-based authorization setting, it assumes each player holds a password pw drawn consistently at random from the wordlist Password of size N. This secret of low-entropy (N is often assumed to be small, i.e. typically less than a million) can be used to authenticate the parties to each other unfortunately, one cannot prevent an rival to choose randomly a password in the vocabulary and to try to copy a player. However such online in-depth search (even if N is not so large) can easily be limited by requiring a slight time interval between successive failed attempts or securing an account after a beginning of failures. Security against such active attacks is measured in the number of passwords the rival can "erase" from the



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candidate list after a failure. Other hand, off-line full search cannot be limited by such practical performances or computational resources considerations. They can be prevented if the protocol is carefully designed and ensures that no data about the password can leak from passively listen in transcripts, but also from active attacks.[3]

4) Shared and Searchable encrypted data for untrusted severs.[4]An RSA-Based Proxy Encryption Scheme a proxy encryption scheme, a ciphertext encrypted by one key can be transformed by a proxy function into the matching ciphertext for another key without revealing any information about the keys and the plaintext. Applications of proxy encryption include: secure email lists, access control systems and attribute based publishing of data. It has feature like, the keys can be easily revoked without any overhead. But, Authorized user in the system has his own keys to encrypt and decrypt data.[4]

5)Search in encrypted data: Theoretical models and practical applications [6]It is the encryption block size, due to the fact that encryption techniques, either symmetric or asymmetric, 16 only takes a binary string of certain length as input. This means that a data item is potentially needed to be split into blocks before encryption is done in a block wise manner. For instance, in the SED scheme from , every word is treated as a block and only exactly matching is supported. In order to support search in a fuzzy manner, the data item may need to be encrypted letter-wise instead of word-wise. Referring to the analysis , the downgrade of encryption granularity may result in additional security issues, such as more severe statistical analysis. In practice, the encryption block size will not only be limited by the encryption algorithms, it also depends on the search criteria to be supported. This paper supports Security issues facing SED schemes which are provably secure in their respective security models.[6]

6)Dynamic Group Diffe-Hellman Key Exchange [7] Diffie-Hellman key exchange (D-H) is a encryption that allows two parties that have no prior knowledge of each other to equally establish a shared secret key over an insecure communications channel. This key can then be used to encrypt successive communications using a symmetric key cipher. This paper have multiple advantages like, It provide strong corruption of data. And It access single casting of data . But It does not support multi cast group.[7]

Sr.	Paper Name	Advantages	Disadvantages	Review
No				
1.	Selective Document Retrieval from Encrypted Database	Provide flexible services.	It is not efficient to provide Multi-user authentication Services.	Provide better security with the help of selective document retrieval.
2.	Privacy Preserving Keyword Search over Encrypted Cloud Data	 Good retrieval performance Provide better security for indexes. 	Track unnecessary network traffic for retrieval of data.	Provide security under different attackers model with high performance.
3.	Password-based Group Key Exchange in a Constant Number of Rounds	 constant-round password-based key exchange protocol for group, derived from the Burmester-Desmedt scheme The use protocol is secure against dictionary attacks under the DDH 	It only requires four rounds of communication and four multi- exponentiations per user.	Security is provide by using the protocol namely password- based constant-round group key exchange. Due to This communication can be happened in minimum rounds.

IV.LITERATURE REVIEW



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		assumption.		
4.	Shared and searchable encrypted data for untrusted servers.	The keys can be easily revoked without any overhead.	Authorized user in the system has his own keys to encrypt and decrypt data.	It provide security as well as revocation. when unauthorized people can access data then key can be revoke by different techniques.
5.	Search in encrypted data: Theoretical models and practical applications	Security issues facing SED schemes which are provably secure in their respective security models.		 Provide two orthogonal categorizations and review the related security models for each category of SED schemes. -Analyze the practical issues related to SED schemes and identify some future research directions.
6.	Dynamic Group Diffe-Hellman Key Exchange under Standard Assumptions	Strong corruption data	Do not support multicast group.	Identify future work for support multicast data.

V.CONCLUSION

In this review, all search scheme that provides both privacy protection capability with less overhead has been proposed. Results on an encrypted data and security analysis using different models show that data privacy can be preserved while retaining very good retrieval performance using enhanced algorithm. Future work will further improve the efficiency and security of search and secure data with the trusted user.

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