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# Urdhwa Multiplier using XOR-XNOR based 4:2 and 7:2 Compressors

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ABSTRACT: Digital signal processor forms a very important part for any signal processing application.

While designing a processor, high speed processing and low area design are the two key factors of concern in today's era. It is a well-known fact that a multiplier unit forms an integral part of processor design. Due to this regard, high speed multiplier architectures have become the need of the day. In this paper, we have introduced urdhawa multiplier, which is considered to be the fastest multiplier because of its feature of parallel calculation. Also, three designs for Urdhawa Multiplier are further developed and described. In first design, 5:3 compressors based on full adder and utilization in term of 42 delays and 21 areas is developed. In second design we have developed 5:3 compressors based on XOR gate and utilization in term of 36 delays and 24 areas. In third design we have developed 5:3 compressors based on full adder and utilization in term of 28 delays and 18 areas. All the designs and experiments were carried out on Xilinx Vertex-7 series of FPGA and the timing and area of the design on the same have been calculated.

**KEYWORDS**: 4:2 Compressor based on Full Adder, 4:2 Compressor based on XOR Gate, 4:2 Compressor based on XOR-XNOR Gate.

### I. INTRODUCTION

Digital signal processing (DSP) is generally defined as the mathematical manipulation of an information signal to modify, analyse or improve it in a particular way. It can be characterized either by the representation of discrete time, discrete frequency or other discrete domain signals by a sequence of numbers or symbols and the processing of these signals [1].

The main goal of DSP is usually to measure, filter and/or compress continuous real-world analog signals. The first step is usually to convert a signal from an analog to a digital form This is done by sampling and then digitizing it using an analog-to-digital converter (ADC), which in turnconverts the analog signal into a stream of numbers. Often the required output signal is another analog output signal, which requires another digital-to-analog converter (DAC). Though this process is more complex than analog processing and has a discrete value range, the application of computational power to digital signal processing allows many advantages over analog processing in many fields, such as error detection and correction in transmission as well as data compression. DSP algorithms were generally performed on standard computers, as well as on specialized processors called digital signal processor and also on purpose-built hardware such as application-specific integrated circuit (ASICs). Today, there are additional technologies used for digital signal processing that include more powerful general purpose microprocessors, field-programmable gate arrays (FPGAs), digital signal controllers (mostly for industrial apps such as motor control [2-3].

### II. URDHWA MULTIPLIER

Vedic mathematics is considered to be an ancient fast calculation mathematics technique which is taken from historical ancient book of wisdom. It is an ancient Vedic mathematics technique that provides the unique technique of mental calculation with the help of simple rules and principles. Swami Bharati Krishna Tirtha (1884-1960), former Jagadguru Sankaracharya of Puri culled set of 16 Sutras (aphorisms) and 13 Sub - Sutras (corollaries) from the Atharva Veda. He developed methods and techniques for amplifying the principles contained in the formulas and their sub-formulas, and called it Vedic Mathematics. According to him, there has been considerable literature on Mathematics in the Veda-sakhas.



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Vedic mathematics is combined from four Vedas (books of wisdom). It covers explanation of several modern mathematical terms including arithmetic, geometry (plane, co-ordinate), trigonometry, quadratic equations, factorization and even calculus.

#### • 4:2Compressor based on Full Adder

Compressor is generally used generally used instead of other adders to add binary numbers with minimal carry propagation. Compressor is a modern digital circuit used for high speed calculations with minimum gates required in the designing technique. Thus, a compressor becomes an essential tool for fast multiplication and adding technique by keeping an eye on fast processor and lesser area.

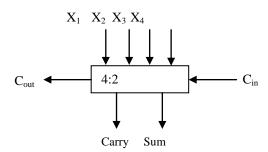


Figure 1: Block Diagram of 5:3 Compressors

4:2 compressors adds 4 bits and one carry, in turn producing a 3 bit output. The 5:3 compressors has 4 inputs  $G_1$ ,  $G_2$ ,  $G_3$  and  $G_4$  and 2 outputs Sum and Carry along with a Carry-in (Cin) and a Carry-out (Cout) as shown in Figure 1. The input Cin, shown in the figure 2 is the output from the previous lower significant compressor.

The Cout, again is the output to the compressor in the next significant stage. The final critical path calculated is smaller in comparison with an equivalent circuit to add 5 bits using full adders and half adders. The basic equation by which 5:3 compressors is governed by the basic equation

$$X_1 + X_2 + X_3 + X_4 + C_{in} = Sum + 2 * (Carry + C_{out}) (1)$$

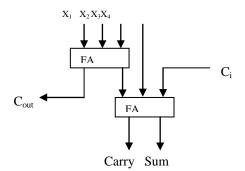


Figure 2: 4:2Compressors based on Full Adder

The standard implementation of the 5:3 compressors is done using 2 Full Adder cells as shown in Figure 2.

### • 5:3 Compressor based on XOR Gate

Figure 3 shows that 5:3 compressor can also be designed using XOR gates when the individual full adders are broken into their constituent XOR blocks. In general case, it can be observed that the overall delay is equal to 4\*XOR. Figure 3

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shows the existing architecture for the implementation of the 5:3 compressor resulting with a delay of 3\*XOR. The equations governing the outputs in the existing architecture are shown below

$$Sum = X_1 \oplus X_2 \oplus X_3 \oplus X_4 \oplus C_{in} \qquad (2)$$

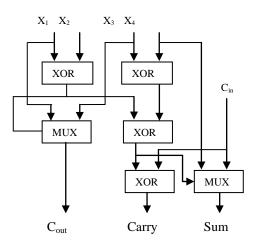


Figure 3: 4:2 Compressors based on XOR Gate

$$C_{out} = (X_1 \bigoplus X_2) \cdot X_3 + (\overline{X_1 + X_2}) \cdot X_3 \quad (3)$$

 $Carry = (X_1 \oplus X_2 \oplus X_3 \oplus X_4) \cdot C_{in} + (\overline{X_1 \oplus X_2 \oplus X_3 \oplus X_4}) \cdot X_4(4)$ 

### • Compressor based on XOR Gate

It is observed that replacing some XOR blocks with multiplexer's results in a significant improvement in delay. It is so because the MUX block at the SUM output gets the select bit before the inputs arrive and thus the transistors are already switched by the time they arrive. This minimizes the delay to a considerable extent. This is shown in figure 4.

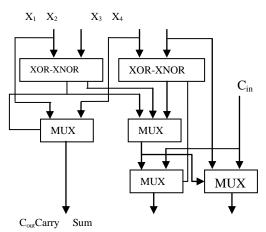


Figure 4: 4:2 Compressors based on XOR-XNOR Gate

The equations governing the outputs in the proposed architecture are shown below

 $Sum = (X_1 \oplus X_2) \cdot (X_3 \oplus X_1) \top (X_1 \oplus X_2) \cdot (X_3 \oplus X_1) \cdot C_{in}$ (5)  $C_{out} = (X_1 \oplus X_2) \cdot X_3 + (\overline{X_1 + X_2}) \cdot X_1$ (6)

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 $Carry = (X_1 \oplus X_2 \oplus X_3 \oplus X_4) \cdot C_{in} + (\overline{X_1 \oplus X_2 \oplus X_3 \oplus X_4}) \cdot X_4(7)$ 

#### • 7:2Compressor

A 7:2 compressors as shown in figure 5, is capable of adding 7 bits of input and 2 carry's from the previous stages, at a time, almost similar to 4:2 compressor which is capable of adding 4 inputs and resulting in 2 outputs.

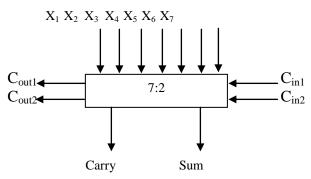


Figure 5: Block Diagram of 7:2 Compressors

In our implementation, we have designed a novel 9:4 compressor utilizing two 5:3 compressors, two full adders and one half adders. The architecture for the same has been shown in Figure 6.

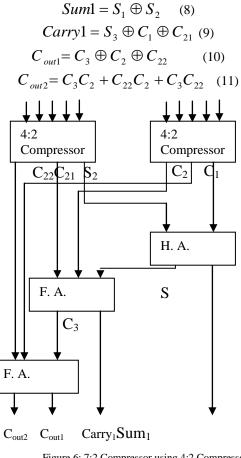


Figure 6: 7:2 Compressor using 4:2 Compressor DOI: 10.15680/IJIRCCE.2017. 0504337



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### III. DELAY AND AREA EVALUATION METHODOLOGY

As per the delay and area evaluation methodology, all gates to be made up of AND, OR, and Inverter (AOI), each one of them having delay equal to 1 unit and area equal to 1 unit.

Adder Blocks	Delay	Area
XOR	3	5
2:1 MUX	3	4
Half Adder	3	6
Full Adder	6	13

Table 1: Delay and Area Count of the Basic Blocks of Urdhwa Multiplier

We then add up the number of gates in the longest path of a logic block that contributes to the maximum delay. The area evaluation is done by counting the total number of AOI gates required for each logic block.

Architecture	No. of Gate	Delay
	Count	
5:3 Compressor based on	42	21
Full Adder		
9:4 Compressor	116	57
Compressor based	1132	591
Urdhwa Multiplier		
5:3 Compressor based on	36	24
XOR Gate		
Modified 9:4	104	63
Compressor		
Modified Compressor	1012	654
based Multiplier		
5:3 Compressor based on	28	18
XOR-XNOR Gate		
Proposed 7:3	88	51
Compressor		
Proposed	852	531
Compressor based		
Multiplier		

Table II: Delay and Area Count of the of Urdhwa Multiplier



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#### **IV. SIMULATION RESULT**

We functionally verified each unit presented in this paper including all three 4:2 Compressor, 7:2 Compressor, Compressor based Urdhwa multiplier. We have been found from the results shown in Table 3 respectively, that number of slices used is same in case of 4:2 compressor based on Full adder and 4:2 compressor based on XOR gate which is less than slices used in 4:2 compressor based on XOR-XNOR gate.

Design	No. of	No. of 4	MCPD
	slices	input LUTs	(ns)
Urdhwa Multiplier	108	190	25.816
(Full Adder)			
Urdhwa Multiplier	88	158	21.011
(XOR Gate)			
Urdhwa Multiplier	78	143	20.951
(XOR-XNOR Gate)			

Table 3: Device utilization summary (Vertex-4) of Urdhwa multiplier

#### V. CONCLUSION AND FUTURE WORK

Among all three designs, proposed Urdhwa multiplier based on XOR-XNOR gate provides the least amount of Maximum combinational path delay (MCDP). Also, it takes least number of slices i.e. occupy least area among all three design.

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