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A System to Counteract Sniffing Assaults Over System

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ABSTRACT: by and large Malicious clients utilize various assaults at various dimensions to take diverse dimension of information. A portion of the sniffing assaults that can be utilized in various dimensions of systems administration/transmission are Media Access Control (MAC) Flooding, Dynamic Host Configuration Protocol (DHCP) Attacks, DHCP Starvation Attack, Rogue DHCP Server Attack, Address Resolution Protocol (ARP) Spoofing, MAC ridiculing and Domain Name Server (DNS) Poisoning. In this paper, a relative report has been finished with the previously mentioned sniffing assaults and the dimension of recuperation that should be possible with each sniffing assault.

KEYWORDS: Sniffing, MAC Flooding, DHCP Attacks, ARP Spoofing

I. INTRODUCTION

Packet sniffing is a form of wiretapping applied to the Networks. Through this, attacker can monitor and capture data from networks. Attacker can capture the following information: 1.File Transfer Protocol (FTP) Passwords, 2.Router Configuration, 3.Telnet Passwords 4.DNS traffic 5.Email Traffic 6.Web Traffic 7.Chat Sessions.

Computers can be connected by bus or switch. If the network is connected by bus then it is termed as shared Ethernet. If it is connected by switch then it is termed as switched Ethernet. In Shared Ethernet environment if a packet is send means it will broadcast the packet to all machines. Only destination machine will acknowledge. Other Machines will ignore the packet. This rule is ignored by attacker. In switched Ethernet the switch maintains a table which contains the MAC address of all computers in the Network. So the message is sent to the destination machine only. Though Switch is more secure than Hub, sniffing can be done even in switch.

Sniffingis categorized into various types:1.MACflooding,2.DNS poisoning3.ARP Poisoning4.Dynamic Host Configuration Protocol (DHCP) Attacks5.Password sniffing



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Figure 1: Sniffer in a Local Area Network

The Protocols vulnerable to sniffing are1.Telnet 2.HTTP, 3.POP, 4.IMAP, 5.SMTP, 6.FTP

Application	Application Stream	Application
	POP3, IMAP,	
resentation	IM, SSL, SSH	Presentation
Session	mis	Session
Transport	Protocols/Ports	Transport
Network	IP Addresses	Network
Data Link	Initial Compromise	Data Link
	Physical Links	

Figure 2: A Sniffer monitors a network and sniffs data

Sniffers operate at the Data link layer. If that layer is sniffer then the upper layers are compromised.

II.MAC FLOODING

Media Access Control (MAC) is a physical address which identifies each node in a computer network. Content Address Memory (CAM) stores the MAC address in switch. The size of CAM is fixed. If the CAM table is flooded with MAC address beyond its capacity, the switch is turned into a hub. Then attackers can easily sniff data



Figure3: MAC Flooding

In figure 3, the attacker uses MAC Flooding concept to turn the switch into hub and thus can easily steal sensitive data. [1]



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III.DYNAMIC HOST CONFIGURATION PROTOCOL (DHCP) ATTACKS

DHCP is a client server protocol which provides IP address to a host. It also provides configuration related information like default gateway, subnet mask. A DHCP relay agent, which passes DHCP requests from one LAN to another so that there need not be a DHCP server on every LAN. It involves the following steps

- 1. Client broadcast DHCPDISCOVER request asking for DHCP configuration
- 2. DHCP Relay agent unicast this message to DHCP server
- 3. DHCP server unicasts DHCPOFFER, which contains client and server's MAC address
- 4. Relay agent broadcast DHCPOFFER in the client's network
- 5. Client broadcast DHCPREQUEST asking for DHCP configuration
- 6. Server unicasts DHCPPACK which contains configuration information.

IV. DHCP STARVATION ATTACK

In this type, the hacker requests large number of DHCPREQUEST and uses all the available IP Address. As a result, the DHCP server cannot issue any more IP Address and in turn leads to Denial of Service (DOS) attack.



Figure 4: DHCP Starvation Attack

In Figure 4: Attacker Requests a large number of IP addresses to the DHCP server which results in Denial of Service to the other users.

V. ROGUE DHCP SERVER ATTACK

In this attacker will introduce a rogue server. DHCP Server and rogue server both will respond to the DHCP request of the client. The server which responds will be taken by first. The rogue server will respond first. Client will send data to the rogue server which in turn will send to actual server. Thereby attacker monitors and captures all sensitive data. The client will be unaware of all these attacks.

VI. DEFENSE AGAINST DHCP ATTACKS

Port security will limit the maximum number of MAC addresses on the switch port, thereby DHCP attacks can be avoided. DHCP snooping feature is available on switches. In order to defend against rogue DHCP servers, DHCP snooping is configured on the port on which valid DHCP server is connected. Thereby the switch will not allow the other ports to respond to DHCP request. [1]

VII. IV.ARP SPOOFING

Address Resolution Protocol (ARP) is a stateless protocol which resolves IP address to MAC Address. If the MAC address is not present in the ARP table, the node will broadcast ARP request. All the nodes will compare their IP



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address with this. Only one of the node identifies this and respond with this. ARP provides no means of authenticity checking.

The attacker can send any arbitrary IP and Mac address. The victim's computer ARP table store this malicious content.



Figure5: ARP Poisoning

Figure 5 explains the concept of ARP Poisoning. The node 10.0.0.7 stores the IP address of 10.0.0.1 and attacker's MAC Address. Similar ARP Poisoning is done for 10.0.0.1 also.

The threats of ARP Poisoning include:1. Packet Sniffing, 2.SessionHijacking, 3.Manipulating Data, 4.Man in the Middle Attack, 5.Connection Hijacking, 6.Connection Resetting, 7.Stealing Passwords, and 8.Denial of Service (DOS) Attacks. [3]

VIII. DEFENSE AGAINST ARP POISONING DYNAMIC ARP INSPECTION (DAI) VALIDATES THE ARP

Packets in a network. DHCP snooping must be enabled prior to DAI. DAI performs IP to MAC address binding inspection stored in the DHCP snooping database. If any invalid binding is found then discard the ARP Packets. Thus Man in the Middle Attacks is eliminated. [2]

IX. V. SPOOFING ATTACKS

Spoofing allows the hacker to pretend to be authorized user. The types are 1.MAC Spoofing 2.ICMP Router discovery protocol (IRDP)

X. MAC SPOOFING

MAC addresses are used for authorization. These are permanent by design and can be changed by physical hardware.MAC spoofing is nothing but forging MAC address.MAC duplicating or spoofing is sniffing one of the client's MAC address and reusing it. Malicious user can listen and receive all the traffic to the legitimate user.

Ways to defend MAC Spoofing:

1. Port security enabling is a way to mitigate MAC spoofing

2. By using DHCP snooping binding table. This table contains the legitimate user's MAC address and IP address. It acts as a firewall between legitimate and malicious users.

3. Dynamic ARP inspection: It verifies MAC address and IP address for all the packets. If any invalid address is found then they are dropped.

XI. IRDP SPOOFING

IRDP is an extension of ICMP protocol. IRDP protocol allows a router to identify the active router's IP address by its advertisements. It allows nodes to listen "Router Advertisement". When the node receives, this may lead to change



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in the routing table. The node does not check for authenticity of the message. A malicious user can spoof the router itself. A Hacker may change the default route provided by server. Thus the attacker can sniff data



XII. DNS POISONING

Domain Name Server is a Protocol which translates the Domain name to IP address. It maintains a DNS table which contains domain name and its IP address. In DNS Poising/DNS spoofing can be done by manipulating DNS table. So the victim is redirected to malicious server instead of actual server. Once connected it can sniff data. There are 4 ways of DNS Poisoning is there.

1. Intranet DNS spoofing, 2.Internet Spoofing, 3.Proxy server DNS Poisoning, 4.DNS cache poisoning

XIII. INTRANET DNS POISONING

When DNS poisoning is done in Local Area Network then it is known as Intranet DNS Poisoning. The attacker poisons the router in order to get DNS request to his machine. Then DNS response is sent to the legitimate user and thus legitimate user is redirected to fake website which is created by malicious user instead of actual website. Thus hacker can sniff important data like passwords. So once sniffing is done, the client is again redirected to actual website and client is totally unaware of the scenario.



Figure7: DNS Poisoning

XIV. INTERNET DNS SPOOFING

It is otherwise known as remote DNS Poisoning. Attacker sends a Trojan to legitimate user which changes the DNS IP address to attacker's IP address. Thus sniffing confidential data is done.



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XV. PROXY SERVER DNS POISONING

Attacker sends a Trojan to change Proxy settings of the user to malicious user's fake website. Thus data can be sniffed.

XVI. DNS CACHE POISONING

The DNS system contains cache which contains domain names and respective IP addressed. When a request comes it first checks DNS cache. If the entries are found then it is redirected accordingly. This DNS resolver cache is modified by attacker. If the DNS resolver cannot validate the DNS responses then it is redirected to malicious server. [4]

XVII. SNIFFING TOOLS

The various sniffing tools used by hackers are as follows:

- ξ Wireshark allows us to capture live network data. This network data can be filtered by IP address, Protocols and Ports.
- ξ Capsa Network Analyzer captures all data transmitted over the network. It captures IP address and MAC address of each host in the network.
- ξ MSN Sniffer 2 captures MSN chats across all computers in the same LAN.
- ξ Colasoft Packet is a packet generator or Packet Editor Tool. Attacker can use this tool to create malicious network packets to carry out the attack on the network.

XVIII. STEPS USED BY SNIFFERS

Step 1: Hacker finds out the switch of the network and connects to one of its ports.

- Step 2: Once connected, discovers network topology using network discovery hacking tools like wireshark
- Step 3: By analyzing the network, sniffer gets the IP of victim's machine using tool like Capsa Network Analyzer
- Step 4: Once he knows IP of the victim, he sends fake messages using the tool like Colasoft Packet
- Step 5: The Previous step results in the Man in the Middle Attack (MITM)
- Step 6: Now the attacker succeeds in sniffing the packets.

XIX. DEFENSE AGAINST SNIFFING

The countermeasures are as follows:

- ξ Use Encryption to protect confidential information.
- ξ Use static IP addresses and static ARP tables to prevent attackers from adding the spoofed ARP entries.
- ξ Restrict the network to authorized users only
- ξ Use IPV6 instead of IPV4
- ξ Use encrypted session such as SSH instead of telnet.
- ξ Use HTTPS instead of HTTP to protect usernames and passwords.
- ξ Use of switches is better than Hubs as switches deliver data only to the recipient.
- ξ Password authenticate the shared folders and services
- ξ Encrypt the communication between PC and access point to prevent MAC address spoofing

XX. SNIFFER DETECTION TECHNIQUE

The sniffer leaves no trace, since it does not transmit data. Systems should be checked for promiscuous mode. Promiscuous mode is a mode of the Network Interface Card (NIC) that allows all the packets without validating its destination address. Stand alone sniffers are difficult to detect as they do not send packets. The reverse DNS Lookup



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can be used to detect non stand-alone sniffers. An Intrusion Detection System is a security mechanism that helps us to detect sniffing activities.

Ping Method: A Ping request has to be sent to the suspect machine with incorrect MAC address. Only the sniffer will respond to this.

ARP Method: ARP should be added to all the nodes. The node which runs in promiscuous mode will cache the MAC address. The Ping message should be sent with different MAC. Only the sniffer node will be able to respond.

Promiscuous Detection Tool: Promqryui tool will helps us to detect the node which is in promiscuous mode.

XXI. PROPOSED SOLUTION

To detect Sniffing is to do sniffing pen testing

1. Using the tool macof, the switch should be checked for failopen mode where it broadcasts the data and this may lead to sniffing. If so the system is prone to MAC flooding.

2. Using the tool Gobbler, DHCP starvation attack can be checked.

3. WinArpattacker tool is used to find ARP poisoning. It changes the ARP table by manipulating MAC address

4. SMAC is a tool used for MAC spoofing and it changes the MAC address in the network card

Once all the tests are performed, it is analyzed and has to take countermeasures to fill the gap in security

XXII. RESULT AND ANALYSIS

Packet sniffing and network capture is done with wireshark. In figure 8, wireshark displays the TCP conversations and thus sniffers can identify port number, sequence number and IP address of a legitimate user.

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8 10.200764	192,168.1.7	209, 85, 227, 18	10	\$2338 > http [404] Sep-1 Ark-1 wir-1/200 ten-0
8 30,204038	392.168.3.2	209.85.227.33	10	(TCP segment of a reassantiled POu)
10 10, 204049	192,168.1.2	209.85.227.19	ALL N	POST_UNETR/Tut=261k=F7bcFFb0058v1ex+c14start=08mm+708axto=58xer+x0x
12 10.267626	209,85.227.29	162.268.1.2	10	http > 52206 [ACK] 580-L Ack+3431 Win+8040 Len+0
12 10.270788	209.45.227.19	392.188.1.7	10	SEED > S2336 [ACR] SHOP ARA-D738 WINH13458 18040
13 10.422381	209.85,227.18	182.166.1.2	107	(TCP segment of a reassestilled PDU)
14 10.423515	209.83.227.10	197.166.1.2	104	(TCP segment of a reassandled PDU)
15 10.423582	192,166.1.2	209.45.227.19	102	52136 > 1020 [ACK] 5604273E A0Xx2724 with427160 Lenve
18 10, 425534	209,85.227.29	192.168.1.2	10	[TCP segment of a reassenbled PO0]
17 10,426887	209.85.227.19	192.248.3.2	10	[TCP segment of a reassanbled POU]
18 10,426926	192,188.1.2	209.85.227.19	107	52336 > Http [ACR] Sep-1788 Ack+5384 win+17160 Let+0
19 20, 473452	209.45.227.19	192.148.1.7	107	[TCP segment of a remisenbled PDu]
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Figure8: Sniffing Network

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inging 192.168.1.1 with 32 bytes of data: ply from 192.168.1.1: bytes=32 time=4ms TTL=254 ply from 192.168.1.1: bytes=32 time=1ms TTL=254 ply from 192.168.1.1: bytes=32 time=1ms TTL=254 ply from 192.168.1.1: bytes=32 time=1ms TTL=254	
ing statistics for 192.168.1.1: Packets: Sent = 4, Received = 4, Lost = 0 (0% proximate round trip times in milli-seconds: Minimum = 1ms, Maximum = 4ms, Average = 1ms	loss),

Figure9: Ping Method

In the above figure ICMP traffic is done using Ping command



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_	2 0.002540	AskeyCom_ed:\$5:9a	LiteonTe_a9:f5:5e	ARP	192.168.1.1 1s at 00:16:e3:ed:55:9a					
-	3.0.007560	3.927.1642.17.2	192.165.1.1	TCHP	Echo (ping) request					
	4 0.004065	192.105.1.1	192.108.1.2	TCMP	acho (ping) reply					
	6 1.002305	102 168 1.1	192, 168, 1, 3	TCMP	Echo (ping) reply					
	7 2,001938	192,168,1,2	192,168,1,1	ICMP	Echo (ping) request					
	8 2,003382	192.168.1.1	192.168.1.2	ICMP	Echo (ping) reply					
	9 3,004165	192.168.1.2	192.168.1.1	ICMP	Echo (ping) request					
	10 3,005770	192.168.1.1	192.168.1.2	ICMP	Echo (ping) reply					
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Figure10: ARP Poisoning

Figure 10 explains ARP Poisoning, ICMP traffic is monitored

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Figure 11: Spoofing a MAC using SMAC

Figure 11 explains spoofing a Mac using smac tool.



Figure 12:Mac list



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Figure 12 in the network which can be spoofed using smac.

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Figure 13:Colasoft capsa Network Analyzer

Figure 13 explains the network traffic which can be spoofed using colasoft

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Figure 14:Promiscuous mode detection

Figure 14 explains promiscuous mode detection using promqry tool

CONCLUSION

This paper proposes an approach to detect packets through packet sniffing. Sniffer is not only used for hacking purpose but also it is used for network traffic analysis, packet/traffic monitoring, troubleshooting and other useful purposes. Packet sniffers can be used in intrusion detection. There exist some tools also that can be used for intrusion detection. Packet sniffing is a technique through which an intrusion can be created and through which an intrusion can be detected.

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