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A Survey on Online Friend Matching Protocol

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ABSTRACT: Mobile social networking (MSN) constitute an assuring Global Positioning System. Connecting remote host within a local physical distance by the wireless use or mobile phone communication. Anyhow in MSN, smart phone users may face the possibility of releasing their personal details and locality information. Here in the study, we diagnose a new security risk originating from private friend discovery protocols, devised runaway attack, that could introduce extreme unfairness problem. to prevent this problem, We proposed a unique Blind Vector Transformation techniques(BT),which can conceal relation between the existing vector and transformed output. On this basis, we introduce our privacy preserving and fairness aware interest and profile matching protocol, which allow one user to match its interest with the other user without exposing its real interest and profile. The precise security evaluation in addition to real world executions exhibit the adequacy and the productivity of the proposed protocol.

KEYWORDS: Epreserving privacy, Discovery of Friend, Mobile portable Networks.

I. INTRODUCTION

Now In the last decade, the multiple users of online networking sites and mobile phone services has increase rapidly. The popular online social networking site, Facebook, has more than 500 million active users, and more than 50% of its active users log on to Facebook at least once per day. In terms of smart phone services, there were 4.1 billion mobile cellular subscribers in total in March 2009 with the propagation of mobile devices, MSNs are becoming devoted part of our life's. Exploiting networked mobile devices such as smart phones and personal digital assistant (PDA) as platforms, MSN not only enables to use their existing (OSNs) at everywhere, but also introduces infinite numbers of mobility-oriented applications, such as location-based services and augmented reality. Among them, an important service is to make new social connections/friends within local physical distance based on the matching of personal profiles. For example, MagnetU is a MSN application that matches one with nearby people for dating or friend-making based on common interests. In such an paper, a user only needs to input some (query) attributes in her profile, and the system would automatically find the persons around with similar profiles. The scopes of these applications are very broad, since people can input anything as they want, such as hobbies, phone contacts and places they have been to. The latter can even be used to find "lost connections" and "familiar strangers".

Privacy preservation is a significant research issue in social networking. The social net-working platforms are extended into the mobile environment, users require more extensive privacy-preservation because they are unfamiliar with the neighbours in close vicinity who may store, and correlate their personal information at different time periods and locations. Once the personal information is correlated to the location information, the behaviour of users will be completely disclosed to the public

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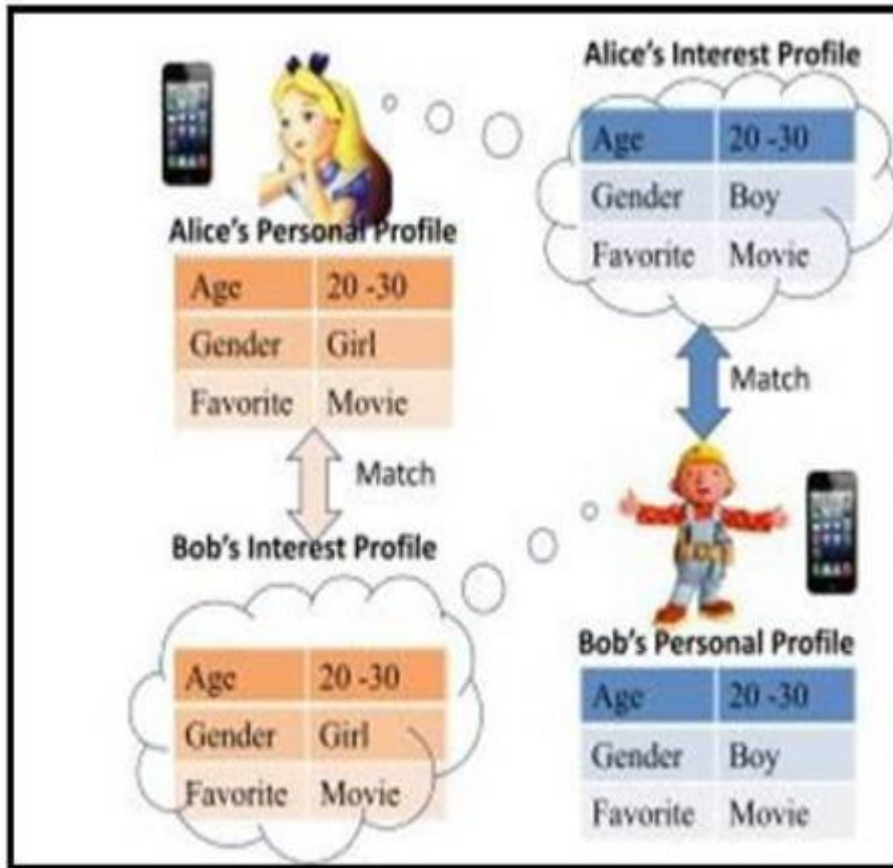


Fig:1.Friend Discovery in Social Network

The content-sharing applications, all of which provide no feedback or control mechanisms to users and may cause inappropriate location and identity information disclosure. Alice has her own profile and Bob has her own. Alice is interested in finding a boy with similar profile and bob is also interested in finding a girl with same profile. A successful matching could be achieved in case that Alice's profile matches Bob's interest while, at the same time, Bob's profile matches Alice's interest. The existing model support this type of mapping process. Further, the existing systems are one-way only and profile matching requires running a protocol twice, with reversed roles in the second run. A malicious attacker or disgraced user launch the runaway attack to exploit this two-pass protocol. This runaway attack in-curs a serious unfairness issue. The new system for Secure Profile Matching, which allow multiple users to compare their profiles without leaking private information to each other.

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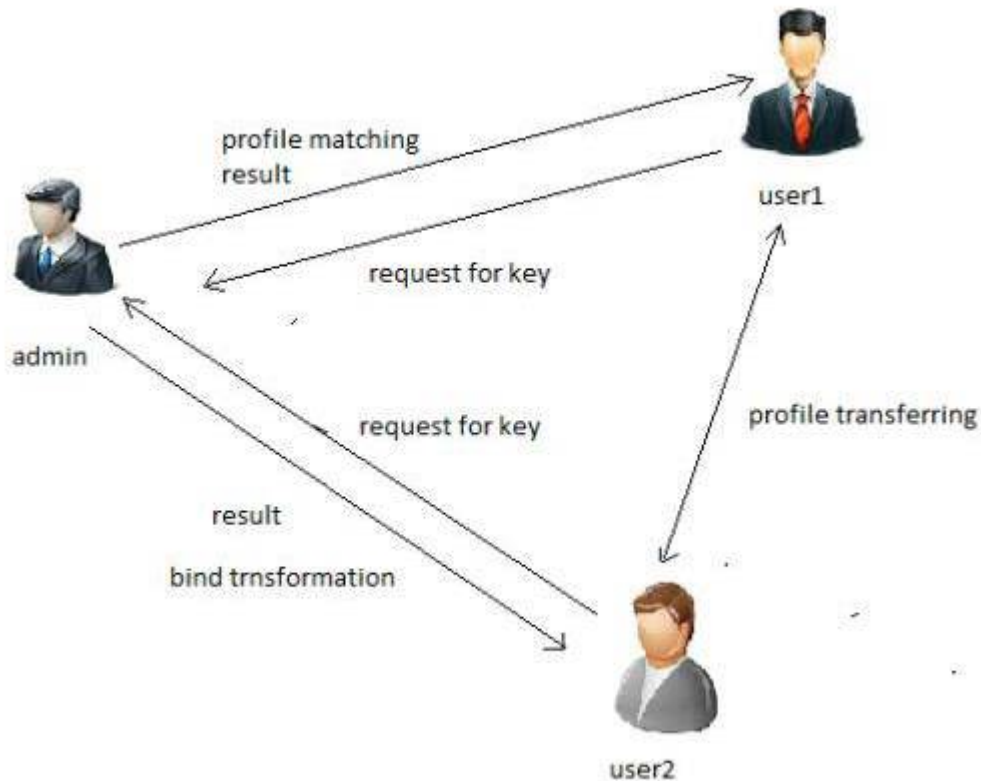


Fig2: System architecture

The secure profile matching problem could then be converted into Private Set Intersection or Private Set Intersection Cardinality. To gain this aim, a novel Blind Vector Transformation Techniques introduce, which could hide the relation between the initial vector and the transformed result. Based on it, the privacy-preserving and fairness-aware friend matching protocol is proposed, which enables one party to match its interest with the profile of another, and vice versa, without revealing their real interest. By the help of Blind Vector Transformation Technique, the interest as well as profile private during match its interests with another's pro-file is kept. To prevent runaway attack, a lightweight verifier checking technique is enable that verify the matching at the minimized overhead and prevent from launching the runaway attack.

II. RELATED WORK

In this part we will present the details of our protocol.

2.1. Protocol Overview

There are two proposed protocols:

1. Blind Vector Transformation
2. Fairness aware and Collusion-free matching protocol

We define five primitive operation as follows:

Encrypt: Given vector performs paillier encryptions with public key pk to get ciphertext. This operation denoted by Encrypt

Vec_add: Given two vectors p and q both are encrypted by paillier encryption, The vec_add operation will be executed to perform add operation

Vec_ext: Given vector, the operation vec_ext could change real value of by diffusion operation. Append some dummy vector to to obtain $p||q$



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Vec_shuffle: Given vector ,the operation vec_shuffle) could change the real value of by performing confusion operation which shuffle the elements in vector .

Vec_rev: Given vector ,the operation Vec_rev could change the real value of by changing value of k element in vector P

Algorithm1:The Blind Transformation Algorithm

Input: P1 \square U1's profile encrypted under his public key pk1, I2 \square interest, n \square the number of interest U2 considers in I2 and s \square a security parameter .

Output:P22 \square The Blind-transform profile vector for U1 , I22 \square the transform interest vector for U2 and rb \square the actual matching result for U2.

Function:

Blind-Transformation(P1, pk1, I2, n, s)

r \square random vector of length n1 = ||P1||

r1 \square Encrypt (r, pk1)

P1' \square VecAdd (P1, r1)

I' \square VecAdd (I2, r)

y \square random vector of length s

y1 \square Encrypt (y, pk1)

P11 \square VecExt (P1', y1)

Rn \square random number between [1,s]

yi \square VecRev (y, Rn)

I21 \square VecExt (Ii, yi)

I22 \square VecShuffle(I21)

P22 \square VecShuffle(P11)

rb \square n + s - Rn

return P22, I22, rb

end function

2.2 The Proposed Blind Transformation Protocol:

The basic idea of Blind Vector Transformation protocol is allowing two untrusted parties to transform two vectors in the blind one by following series of private and identical steps, eg. Adding random vector, shuffling in same order, VecExt and VecRev. Blind Add: U2 generates random vector r and then perform $r1 = \text{Encrypt}(r, pk1)$ after the, U2 calculate $P1' = \text{VecAdd}(P1, r1)$ and $I2' = \text{VecAdd}(I2, r)$ by adding r1 and r to P1 and I2 ,respectively. Blind Append: U2 generates random vectory of length s, where s is predetermined security parameter, then perform $y1 = \text{Encrypt}(y, pk1)$ to get $P11 = \text{VecExt}(P1', y1)$. Blind Reverse: U2 randomly selects $Rn \in \{1, 2, \dots, s\}$ and perform $yi = \text{VecRev}(y, Rn)$ then obtain $I21 = \text{VecExt}(Ii, yi)$. Blind Shuffle: U2 performs $I22 = \text{VecShuffle}(I21)$ and $P22 = \text{VecShuffle}(P11)$ with the same order.

2.3 The proposed Fairness-aware and Collusion free Matching Protocol:

In this proposed protocol, instead of directly sending $h1, h2$ (number of matched entries) to the verifier, an additional blind linear transformation is used To check if two parties' interests match their equivalent profiles, we have the following verification protocol. First-ly, U1 obtains $P'1 = \text{Decrypt}(P22, sk1)$ by performing the decryption operation with his own secret key $sk1$. After obtaining $P'1$, U1 compares it with U2's blinded interests I2 to get the number of matched entries $s'2$ while U2 could get $s'1$ similarly. To verify if their interests and the profiles match or not, U1 sends $h1 = H(s'1 || s'2)$ whereas U2 sends $h2 = H(s'1 || s')$ to a randomly chosen verifier. The verifier could verify whether $h1 || h2$. If $h1 || h2$, the match succeeds, otherwise, it fails.

III. CONCLUSION

Equalize In this work, we have develop a proposed protocol that will ensure the fairness aware and privacy preserving profile matching process and also provide more security in social network.

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