



## International Journal of Innovative Research in Computer and Communication Engineering

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# A Dynamic and Efficient Multi-Keyword Ranked Search over Encrypted Cloud Data Using Secure Searching Encryption Algorithm

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**ABSTRACT:** By the advantages of the cloud computing software as a service, many companies and large scale organization are motivated to outsource their valuable data to the cloud databases for its costless service, easily scalable, and able to access anytime from anywhere, minimizes the cost of management. Due to privacy issues of data in public cloud databases, data should be encrypted before outsourcing to cloud environment. However, searching on outsourced data is obsolete. Data utilization requires keyword based document retrieval. In this, paper we proposed a scheme to support Multi-keyword Ranked Search and simultaneously allows the data owner to perform dynamic operations like insertion and deletion of documents. For that, we use the TF&IDF model in index construction and query generation. Searchable encryption is the cryptographic method used for providing security, for that we explore some effective cryptographic algorithm is CRSA and construct a special tree called balanced M-Way search tree to propose “Greedy-Depth First Search” process to provide Multi-keyword Ranked Search, meanwhile calculating the relevance score between encrypted index and query vectors to providing Top-k Ranked results.

**KEYWORDS:** Cloud computing, Searchable encryption, Multi-keyword ranked search, Dynamic operations, Depth-first search technique (DFST), Commutative RSA (CRSA).

### I. INTRODUCTION

Cloud computing is the embryonic technology; it is a new model of enterprise IT infrastructure. Cloud computing providing several resources (databases, cpu cycles, network bandwidth, applications, software's) as services over internet. It provide convenient and on demand service to both individuals and organizations for reduces the economic overhead, its scalability and less hands on management. By these appealing features, more and more organizations and individuals are motivated to outsource their sensitive information such as e-mails, health records, and company financial data, documents related to new product releases and government documents, certificates into cloud servers.

There are many cloud service providers in the market to providing the on demand service in pay-as you-use model with less cost, for example Amazon web services, Google app engine, Microsoft azure, Sales force, IBM cloud services are substantially adding to their databases. These services are shared among many users and CSP's (Cloud Service Providers). They keep sensitive and user's outsourced data for user availability with effective recovery mechanisms. To providing confidentiality for user's documents, data should be encrypted with different security keys before outsourcing cloud environment, because we don't say all the data users are honest or a user unfotunatly may lead to modify the content of other user's document. To check user authority we preferably use the authentication mechanisms like password based user logins.

To address these problems, researchers are proposed encryption methods to contributing searching process in cloud. Those searching algorithms are referred as searchable encryption (SE) schemes; they have made specific contributions in terms of efficiency, functionality, security and searching. It should not leak any keyword information to the cloud server. It is very suitable in “pay-as-use” model. In this paper, proposed methods are intended to solve



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some problems of multi-keyword ranked search over encrypted cloud data and security by preserving strict system wise privacy in cloud environment.

The remainder of this paper is organized as Related work/literature survey is discussed in section 2, section 3 describes the problem formulation with system models and section 4 contains proposed search, encryption schemes and section 5 shows the results, section 6 gives conclusion and future work, references are discussed in later section and figure 1 shows the architecture and figure 2,3 shows the results.

## II. LITERATURE SURVEY

The data encryption is used to protect the data confidentiality. Searchable encryption schemes enable the data owner's to store encrypted in cloud server and execute keyword search over cipher text domain. When come to searching on cipher text, the efficiency gets low. This inefficiency made lead to leakage of information to unauthorized people. So, previously many research works have proposed to words for efficient searching over encrypted cloud data.

For the first time, *Song et al.* [1] proposed the first symmetric searchable encryption (SSE) scheme, in this file is encrypted word by word. The drawback of this scheme is word frequency is relieved and search time is linear to size of documents collection. *Change et al.*[2] proposed scheme index is build for each document. This scheme is more secure than Goh's scheme but less efficient and does not support arbitrary update of new words. *Curtmola et al.*[3] for the first time proposed two searchable encryption schemes to achieve the optimal search time. *Kamara et al.* proposed extension to SSE schemes called dynamic SSE, where insertion and deletion of documents can perform in index table, but all these supports only single keyword search. Later there are many searching methods implemented in cloud, such as single keyword Boolean search, which are very simple in functionality, but not efficient from different threat models. After abundant works have been proposed against different threat models to achieve various search functionality, such as single keyword search, exact or fuzzy keyword search, similarity search [1], [2], [3], [4], [5], [6], [7], ranked search [8], [9], [10].

All these proposed multi-keyword ranked search schemes retrieve search results based on the existence of keywords, which cannot provide acceptable result ranking functionality. Ranked search can returning most relevant document. Some early works [8],[9],[10], realized that ranked search using order-preserving techniques, but they support single keyword search. *N.Cao et al.*[11] scheme, in which documents and queries are represented as vector of dictionary size with "coordinate matching", the documents are ranked according to the number of matched query keywords, but this scheme is not accurate enough and search efficiency is  $O(n)$ . *Sun et al.*[12] scheme supports similarity-based ranking. In this, constructed a searchable index tree based on vector space model and adapted cosine measure together with  $TF \times IDF$  to provide ranking. This algorithm achieves better than linear search time but results in precision loss. *C.Orencik et al.*[13] proposed a search method, which utilizes the local sensitive hash(LSH) functions to cluster the similar documents. This scheme suitable for similarity search but results cannot give exact ranking. In [14] *Zhang et al.* proposed a method to deal with secure multi-keyword ranked search scheme, in this different data owners use different secret keys to encrypt their documents.

Although researchers from many universities had been investigating to identify a suitable privacy preserving search techniques for cloud domain. There exist wide range of research challenges in searching and privacy for encrypted data, however, we intended to choose the work to meeting their challenges.

## III. PROBLEM FORMULATION

### A. Notations and Preliminaries:

- D: This is collection of plaintext documents stored in cloud server  $D = \{d_1, d_2, \dots, d_n\}$ . Each document is considered as sequence of keywords.
- n: This is the total number of documents in D.
- T: It is the unencrypted form of index tree for the whole document collection D.
- I: Is the searchable encrypted tree index generated from T.
- Q: It is the query vector for keyword set  $W_q$ .
- TD: This is referred as trapdoor for search request, it is the encrypted form of Q.
- W: This is the total number of keywords in W.



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- $W$ : It is the dictionary, namely the set keywords, denoted as  $W=\{w_1,w_2,\dots,w_m\}$ .
- $M$ : This is total number of keywords in  $W$ .
- $W_q$ : This is the subset of  $W$ , representing keywords in the query.
- $P_u$ : Is the index vector stored in tree node  $u$ , whose dimensions is equal to cardinality of the dictionary  $W$ .
- *Note*: 'u' can be either a leaf node or internal node in the tree.
- $C_u$ : This is encrypted form of  $P_u$ .

### B. The Vector Space Model for Relevance Score:

The collection of documents can be represented as vectors in vector space model with TF×IDF model supports ranked multi-keyword search efficiently. Here,  $TF_{(w,d)}$  is the term frequency defined as the number of times keyword (term) 'w' appears in the document 'd'. the number of documents that contains keyword w is known as document frequency(DF). The inverted document frequency ( $IDF_w$ ) of a keyword 'w' obtained from the total number of documents is divided by document frequency. Every document is denoted with a vector, whose elements are the normalized TF values of keywords. The query is also denoted with a vector Q, whose elements are normalized IDF values of query keywords. Naturally, the size of query vector and document vector are equal to the total number of keywords in the dictionary.  $Q \times I_u$  required to calculate the relevance score between the query and corresponding documents.

- $N_{d,w_i}$ : The number of keywords  $w_i$  in document D.
- $N$ : The total number of documents.
- $N_{w_i}$ : The number of documents that contains keyword  $w_i$ .
- $TF_{d,w_i}^1$ : The TF value of  $w_i$  in document d.
- $IDF_{w_i}$ : The IDF value of  $w_i$  in document collection.
- $TF_{u,w_i}$ : The normalized TF value of the keyword  $w_i$  stored in the index vector  $P_u$ .
- $IDF_{w_i}$ : The normalized IDF value of keyword  $w_i$  in document collection.

Relevance score of query vector Q and document vector or index vector  $P_u$  of node u is calculated as  $Rscore(P_u, Q) = P_u \cdot Q$

$$= \sum_{w_i \in W} TF_{u,w_i} \times IDF_{w_i} \quad (1)$$

If u is an internal node of the tree,  $TF_{u,w_i}$  is calculated from index vectors in the child nodes of u, if u is a leaf node, the  $TF_{u,w_i}$  is calculated as:

$$TF_{u,w_i} = TF_{d,w_i}^1 \div \sum_{w_i \in W} w_i \in W (TF_{d,w_i}^1)^2, \text{ Where } TF_{d,w_i}^1 = 1 + \ln N_{d,w_i} \quad (2)$$

In the search query vector Q,  $IDF_{w_i}$  can be calculated as:

$$IDF_{w_i} = IDF_{w_i}^1 \div \sqrt{\sum_{w_i \in W} w_i \in W (IDF_{w_i}^1)^2}, \text{ Where } IDF_{w_i}^1 = \ln(1 + N/N_{w_i}) \quad (3)$$

### C. Keyword-based Binary Search Tree:

In this section, we first construct BMs tree based on the vector space model, it is widely used to deal with optimization problems and is a dynamic data structure whose node stores the vector P, and these values are normalized TF values. Later, in algorithm 2 we discuss DFST (depth first search technique) searching algorithm for BMS tree.

### D. BMS Tree Index Construction:

In the process of index tree construction, we generate node for each document. These nodes are act as leaf nodes in the tree. Then, the internal nodes are formed based on these leaf nodes. The data structure of the node is defined as  $u=(ID, D, child[], DID)$  (4)

Where ID is unique id generated using GenID() function, D I index vector, child[] is pointers to children of the node and DID is a document ID.

In the algorithm 1,  $children.D[i] = \max\{u.P_1 \rightarrow D[i], u.P_r \rightarrow D[i]\}$  Where  $i=1, \dots, m$ ; (5)

#### a) Algorithm 1: BuildBMSIndexTree (D)

- *Input*: The document collection  $D=\{d_1,d_2,\dots,d_n\}$  with the identifiers  $DID=\{DID \setminus DID=1,2,\dots,n\}$ .
- *Output*: The index tree T.
- for each document  $d_{DID}$  in D do
- Construct a leaf node u for  $d_{DID}$ , with  $u.ID=GENID(), u.P_1=u.P_r=null$ ,  $u.DID =DID$ , and  $P[i]=TF_{d_{DID},w_i}$ , for  $i=1,2,\dots,m$ ;
- Insert u to CurrentNodeSet;



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- end for
- while the number of nodes in CurrentNodeSet is larger than 1 do
- if the number nodes in CurrentNodeSet is even, i.e 2h then
- for each pair of nodes u' and u" in CurrentNodeSet do
- Generate a parent node u for u' and u", with  $u.ID=GenID( ), u.P_1=u', u.P_r=u''$ ,  $u.DID=0$  and  $P[i]=\max\{u'.P[i], u''.P[i]\}$  for each  $i=\{1,2,\dots,m\}$
- Insert u to TempNodeSet;
- end for
- else
- for each pair of nodes u' and u" of the former (2h-2) nodes in CurrentNodeSet do
- Generate a parent node u for u' and u";
- Insert u to TempNodeSet;
- end for
- Create a parent node  $u_1$  for the  $(2h-1)^{th}$  and  $2h^{th}$  node, and then create a parent node u for  $u_1$  and the  $(2h+1)^{th}$  node;
- Insert u to TempNodeSet;
- end if
- Replace CurrentNodeSet with TempNodeSet and then clear TempNodeSet;
- end while

## E. Search Process using DFST:

The search process of multi-keyword ranked search scheme is the recursive function upon the BMs tree name as Depth First Search Technique (DFST). We create a result documents as "Ranked list", whose element is denoted as (Score, DID). The elements of Ranked list are in descending order according to score function during the search process. The DFST algorithm is discussed in algorithm 2.

### a) Algorithm 2: DFST (IndexTreeNode u)

- if the node u is not a leaf node then
- if  $RScore(P_u, Q) > k^{th}$  score then
- DFST(u.hchild);
- DFST(u.lchild);
- else
- return
- end if
- else
- if  $RScore(P_u, Q) > k^{th}$  score then
- Delete the element with the smallest relevance score from Ranked list;
- Insert a new element  $RScore((P_u, Q), u.DID)$  and sort all the elements of Ranked list;
- end if
- return
- end if

## F. System Models:

In this paper, we explore cloud computing system model involves three different entities. The responsibility of each entity is discussed as follows:

### 1) Data Owner(DO):

Data owner has a collection of documents  $D=\{d_1, d_2, d_3, \dots, d_n\}$  with valuable information to be outsourced to the cloud server. To provide data security, the documents are encrypted before outsourcing. Data owner creates the dictionary based on keywords extracted from each document and creates the query vector by using the keywords entered by the data user is called trapdoor. To provide user privacy, query is encrypted and sends it to the data user. The data owner sends access control keys to the authorized user.

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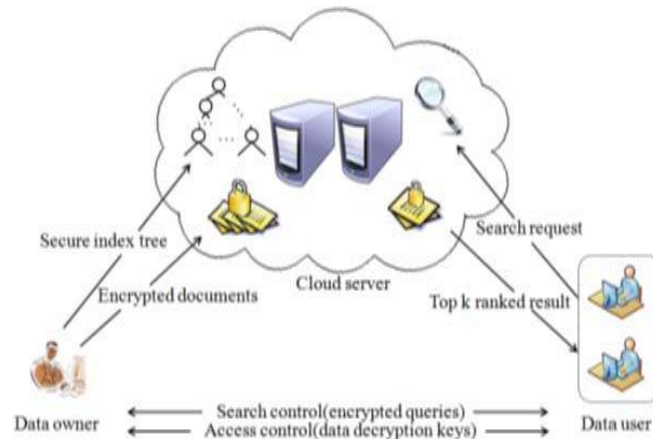


Figure 1: Architecture of multi-keyword ranked search over encrypted cloud data.

## 2) Data Users:

Data users are the authorized users who accessing cloud services and sensitive data from the cloud. The data user create the request for certain documents and send it to data owner, then data user receives trapdoor according to search query request and access control from the data owner. Data user sends these trapdoor and access controls to the cloud server to retrieve documents from the cloud.

## 3) Cloud Service Provider:

Cloud server receives the encrypted documents and encrypted index vectors from data owner and stores into the data owners cloud database. Cloud sever having the capability to take the data request from users and checks the access controls and searches keywords related documents. It will retrieve the documents from cloud storage depending upon the privileges. To increasing document retrieval accuracy from cloud server, the top scored (ranked) documents return to data user. Fig.1 shows the architecture of multi-keyword ranked query search over encrypted cloud data.

## G. Threat Models:

The cloud server is measured as “honest-but-curious” in our proposed scheme. We considered two threat models for identify the system security and performance in different attack capabilities that is as follows:

a) *Known cipher text model*: In this model, the cloud server knows only about encrypted documents and encrypted index vectors, which are outsourced from data owner.

b) *Known background model*: As compared with known cipher text model, the cloud server is much stronger because it knows about term frequency (TF) statistics of specific keywords in the document collection.

1) *Design Goals*: To enable system secure, our proposed system has following design goals.

a) *Dynamic*: The proposed scheme is designed to not only supporting multi-keyword query search and accurate result ranking, but it also supporting dynamic update on document collection.

b) *Search efficiency*: The scheme aims to achieve sub-linear search efficiency by exploring a special tree based index and an efficient search technique.

c) *Privacy-preserving search*: The scheme is designed to prevent the loud server from learning additional information about the document collection, index tree and query.

The specific privacy requirements are summarized as follows:

a) *Index confidentiality and query confidentiality*: the underlying plaintext information including keywords in the index and query, TF values of keywords in index and IDF values of query keywords should be protected from cloud server.

b) *Trapdoors unlinkability*: the cloud server should not be able to determine whether two encrypted queries (trapdoor) are generated from the same search request.

c) *Keyword privacy*: the cloud server could not identify the specific keywords in the query request and index or document collection by analyzing the statistical information like term frequency.





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## IV. PROPOSED SEARCH SCHEMES

### A. Search Module:

In our proposed scheme, Balanced M-way Search Tree (BMS) is used to perform searching. Let's suppose cloud server has received 'n' encrypted documents  $D=\{d_1, d_2, \dots, d_n\}$ . So, if user wants to search with 't' keywords, DO generates a secret trapdoor encrypted using CRSA. The trapdoor containing the encrypted keywords is sent as a token to the server. Then the server uses this trapdoor to match the encrypted keywords in the index tree node. If a match is found, it stores the pointer to that document in the encrypted database. The following Algorithm 3 gives the stepwise information about how search is to be done on B-Tree.

#### 1) Algorithm 3: Search\_query (root, trapdoor)

- *Input:* root, trapdoor containing keyword to be searched.
- *Output:* pointer to the documents containing the keywords, NULL if non-exist.
- $NODE\_x = \text{Disk\_Read}(\text{root})$ .
- if  $NODE\_x$  is an index node
  - (a) If there is an object O in  $NODE\_x$  such that  $O:\text{key}=\text{keyword}$ , return  $O:\text{value}$ .
  - (b) Find the child pointer  $x:\text{child}[i]$  whose key range contains key.
  - (c) Return  $\text{Search\_Query}(NODE\_x:\text{child}[i], \text{key})$ .  
Else if there is an object O in  $NODE\_x$  such that  $O:\text{key}=\text{keyword}$ , return  $O:\text{value}$ .  
Otherwise, return NULL.  
End if.

### B. Ranking Module:

In large databases, it is quite common that the keyword might be matching with more number of documents. It is very difficult to decrypt every document and go through all documents. Therefore there is a need to rank the documents based on their relevance score to the keywords. In our scheme we used  $TF \times IDF$  to rank the documents. The proposed schemes consist of the following phases:

- 1) *Setup:* The proposed system is initialized with this step. In this step, the data owner generates a secret key. DO generates an n-bit random vector S and two  $n \times n$  invertible matrices  $M_1, M_2$ . Hence the secret key SK is a form of three tuples as  $SK=\{M_1, M_2, S\}$ .
- 2) *Generate Index(DC, SK):* Data owner (DO) extracts keywords from document collection and creates a dictionary with keywords and then creates an index vector for each document. DO constructs an unencrypted index tree by calling  $\text{BuildBMSIndexTree}(D)$  algorithm. Now split each index vector  $P_u$  of node u as two random vectors  $P_{u1}$  and  $P_{u2}$  using S-bit vector. If  $S[i]=0$  then  $P_{u1}[i]$  and  $P_{u2}[i]$  will be set to the same value of  $P_u[i]$ . If  $S[i]=1$  then  $P_{u1}[i]$  and  $P_{u2}[i]$  will be set to two random values whose sum is equal to  $P_u[i]$ . The encrypted index vector for the index tree is  $C_u=\{M_1^T \cdot P_{u1}, M_2^T \cdot P_{u2}\}$ .
- 3) *Trapdoor(W):* The interested keywords are entered by the user. It generates a bit vector Q based on keywords and synonyms d, if keywords are available in d then set the IDF values of terms to a specific dimension in Q respectively, otherwise set to 0. So, apply the same splitting technique which we have applied earlier to Q, divide Q into Q' and Q'' to provide search privacy, then encrypt it. The encrypted query vector is denoted as  $TD=\{M_1^{-1}Q', M_2^{-1}Q''\}$ .
- 4) *Query(TD, C\_u):* The query is processed using encrypted keywords of the trapdoor TD and encrypted index vector  $C_u$ , the score of the documents is stored in node u. The same process is applied to all nodes in the index tree. The relevance score calculation of the document is as

$$\begin{aligned}
 \text{Score} &= C_u \cdot T \\
 &= \{M_1^T \cdot P_{u1}, M_2^T \cdot P_{u2}\} \times \{M_1^{-1}Q', M_2^{-1}Q''\} \\
 &= M_1^T \cdot P_{u1} \cdot M_1^{-1} + M_2^T \cdot P_{u2} \cdot M_2^{-1} \cdot Q'' \\
 &= P_{u1}^T \cdot M_1 \times M_1^{-1} \cdot Q' + P_{u2}^T \cdot M_2 \times M_2^{-1} \cdot Q'' \\
 &= P_{u1}^T Q' + P_{u2}^T Q'' \\
 &= P_u^T Q
 \end{aligned}$$

The above used scheme is used to provide the security against known cipher text model. To prevent the cloud server from data access patterns.



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## C. Dynamic Update Operations:

After insertion or deletion of document, need to synchronously update the index by updating the nodes using the document identifiers in the BMS index tree. The process is as follows:

For updating:  $\{I_s, c_i\} \leftarrow \text{GenUpdateInfo}(\text{SK}, T_s, i, \text{updtype})$

Here updtype denotes the insertion or deletion operations of document  $d_i$ . The notation  $T_s$  denotes the set of consisting of the tree nodes that needs to be changed during the update.

## V. RESULTS



Fig 2. Searching files using secret key and filename after successful login.

### Search Results

java.txt  
java.txt  
java.txt  
java1.txt  
java2.txt

Fig 3. Users gets ranked documents and decrypt files with their decryption keys sent to email.

## VI. CONCLUSION AND FUTURE WORK

In this, our work uses CRSA asymmetric algorithm for encrypted data files and index tree based on B-tree. CRSA increases the data privacy by its commutative nature; data in a file can be updated dynamically without affecting the overall performance of searching on BMS tree. If encrypted data is modified (insert and delete), then re-encryption of whole data is not needed. This is desirable feature it reduces the computation time. The future work will be concentrate for better data security in multi-user model for performance improvement, because we cannot say every data user is honest and may leads to any security problems means they can distributes the decrypted documents and share secret key to unauthorized users.

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