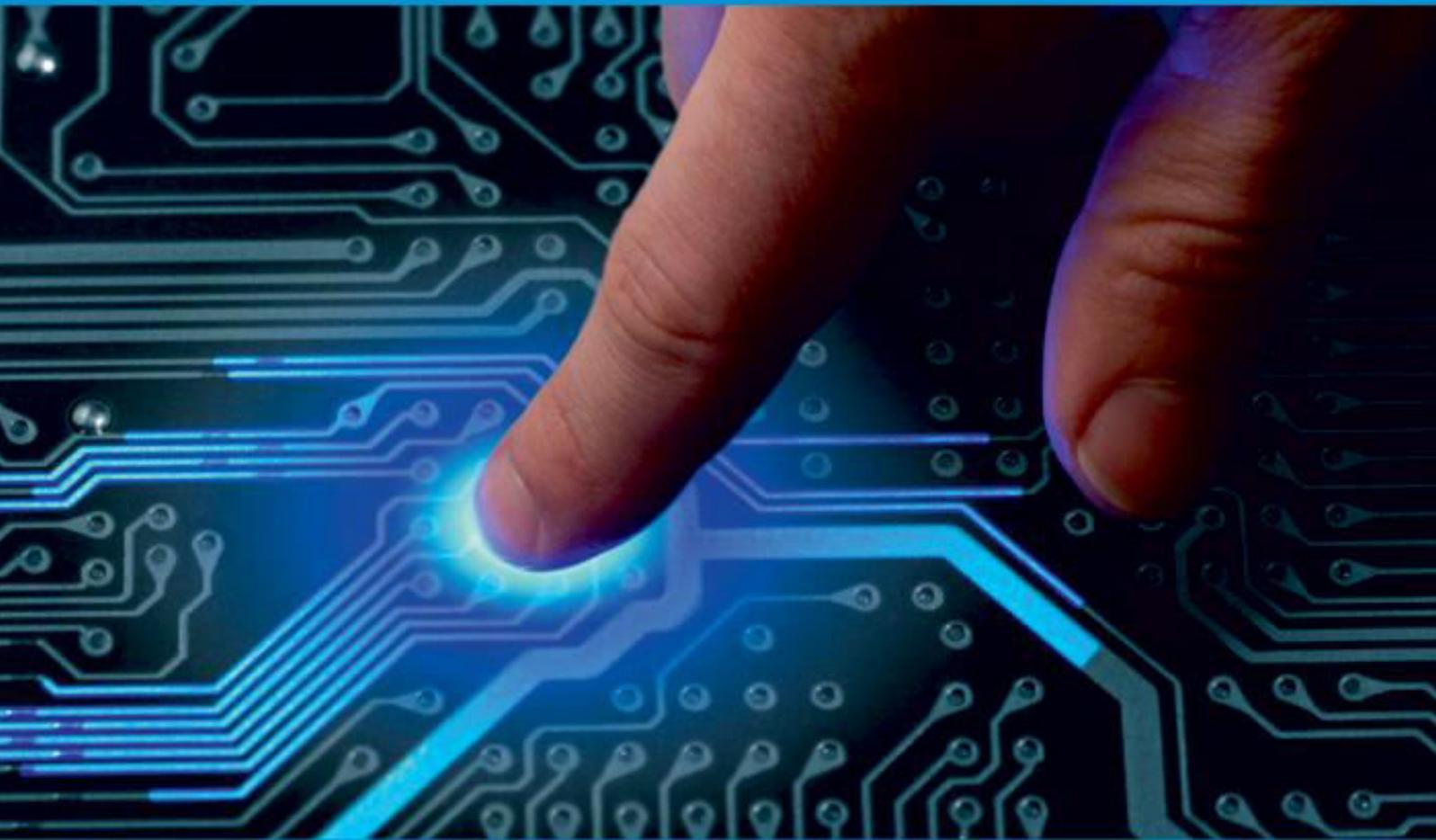




**IJIRCCCE**

e-ISSN: 2320-9801 | p-ISSN: 2320-9798



# INTERNATIONAL JOURNAL OF INNOVATIVE RESEARCH

IN COMPUTER & COMMUNICATION ENGINEERING

**Volume 9, Issue 10, October 2021**

**ISSN** INTERNATIONAL  
STANDARD  
SERIAL  
NUMBER  
INDIA

**Impact Factor: 7.542**



9940 572 462



6381 907 438



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www.ijircce.com

# MPLS-VPN Configuration of Service Provider

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**ABSTRACT:** This paper gives us comprehensive performance analysis of MPLS enable network and IP enable network. It states the behaviour of MPLS protocol with OSPF protocol. We have analysis these two on basis of latency, utilization in the network with the help of graphical network simulator-3. We have used nine cisco Series 3745 router with IOS version 12.4 for testing network performance with MPLS and traditional IP routing. Results obtain in these testing shows how service provider can benefit from MPLS services with increasing network latency and additional benefits obtain from MPLS.

**KEYWORDS:** Internet Protocol, Multiprotocol Label Switching, Virtual Private Network, Open Shortest Path First, Transmission Control Protocol

## I. INTRODUCTION

Computer network were circuit switched, where continuous bit streams carried over the physical links. This was well suitable for voice and data unicast communications. This leads to some severe consequences in case of failure. All the communications over the failed link are interrupted in such situation. These days packet switched networks are used in which data is divided into small chunks called as a packet and these packets are routed over the communication links. Different packets can take different paths. In case of link failure, the packets can be rerouted through alternate available path to avoid failed link and hence communication is not interrupted. This feature makes packet switched networks more reliable but on other hand as packets are routed individually, it is difficult to manage flow of data. Traditional IP networks offer little predictability of service, which is undesirable for applications such as telephony, and for rising and future real-time applications. IP networks are frequently layered over ATM networks, which is very expensive in terms of overhead (adding 25 percent or more of overhead to every IP packet), but had one great advantage, IP networks have no means of tagging or monitoring the packets that cross them. The history tells us the upper limit of transmittable bandwidth doubles and sometimes quadruples every nine to twelve months. We need matching data transferring topologies as well as improved system reliability [2].

Multiprotocol Label Switching is a tool applied in distinguished performance telecommunications networks that carries materials from one complex over to the next. Originally MPLS created by a crew of engineers that were consumed with improving the quickness of routers nevertheless from the time it has emerged as a classic in today's telecommunications. There have been a multitude number of attempts at developing many technologies with the identical goals, to date none have reached the position of success that we now see with MPLS. Appreciating MPLS means looking to some of the parts that exist concerning to the MPLS such as the label which is a locally significant identifier that is allotted to a packet. Every label contains four fields, a label value, traffic class field which determines the quality of service, bottom of stack label which is not always set but when it is it signifies that the label is currently the last in the stack and finally there is the "time to live" (also referred to as TTL) field which is the limit of time that data can experience before it will be discarded. To realize the magnitude of MPLS one just has to measure it against some earlier technologies that are similar like the frame relay which focused on making previously existing physical resources more adequate. In recent days the use of frame relay has been given a poor name in several markets because of overdone bandwidth used by some companies hence making the use of MPLS much more alluring. One more similarity would be that between ATM (also referred to as Asynchronous Transfer Mode) MPLS when comparing the two have many differences both offer connection oriented service to allow for transporting data across networks.

An MPLS connection shows the most significant difference in its approach as they are able to work with various lengths of packets where as an ATM is only capable of dealing with a fixed length. The most favorable difference you

will find between the two is MPLS configuration which was developed specifically for internet protocol. MPLS are just being used only with internet protocol networks and are standard. It can connect to two facilities or can control thousands of locations simultaneously. MPLS does not compete with IP forwarding but it complements IP forwarding. MPLS technology works to solve those flaws of IP, encapsulating IP packets within labels.

Emergence of MPLS is not for replacing IP, it is designed to add a set of rules to IP so that traffic can be classified, marked, and policed. MPLS (Multiprotocol label switching) as a traffic-engineering tool has emerged as an elegant solution to meet the bandwidth management and service requirements for next generation Internet Protocol (IP) based backbone networks. An MPLS network can offer the quality of service guarantees that data transport service like frame relay (FR) or ATM give, without requiring the use of any dedicated lines. The availability of traffic engineering has helped MPLS reach critical mass in terms of service provider mind share and resulting MPLS deployments. Most carriers run MPLS underneath a wide range of services, including FR, wide-area Ethernet, native IP, and ATM. Advantages accrue primarily to the carriers. User benefits include lower cost in most cases, greater control over networks, and more detailed Quality of Services.

## II. MULTI-PROTOCOL LABEL SWITCHING

Multiprotocol Label Switching (MPLS) is a data-carrying mechanism, in computer networking and telecommunications, which is highly scalable and protocol agnostic. Often referred to as "Layer 2.5 protocol" MPLS technology operates between the Data Link layer (Layer 2) and the Network Layer (Layer 3) of the OSI Model. MPLS is part of the family of packet-switched networks. It was designed primarily to provide a unified data-carrying service for Circuit-based as well as Circuit-switching clients. Both the clients offer a datagram service model[2].

Multiprotocol Label Switching enables to carry diverse types of traffic such as Asynchronous Transfer Mode (ATM), Internet Protocol (IP) packets, Synchronous Optical Networking (SONET), and Ethernet frames. Labels are assigned to the data packets in an MPLS network. Based on the label contents, packet-forwarding decisions are made, without necessitating examination of the data packets. Through this feature, end-to-end circuits may be created using any protocol over any type of transport medium. MPLS technology is beneficial as it helps to eliminate the dependence on ATM, Frame relay, SONET, Ethernet, etc., which are Layer 2 technologies. It also does not require multiple data link layer networks to gratify different traffic types. In MPLS technology, a specific path is set up for a given sequence of data packets. These packets are identified by the packet label, thereby saving the time that a router takes to search the address where the packet should next be forwarded. MPLS is referred to as "multiprotocol" since it closely works with IP, ATM, and frame relay network protocols. The major benefits of MPLS networks include:

Traffic Engineering - The capacity to determine the path that the traffic will take through the network.

MPLS VPN - Service providers can create IP tunnels all over their networks using MPLS, which does not necessitate encryption or end-user applications.

Layer 2 services (ATM, Ethernet, frame relay) can be carried over the MPLS core. Simplified network management through elimination of multiple layers. MPLS has become popular due to its capability to form multi-service networks with high speed. It can support pre-provisioned routes that are virtual circuits known as Label-Switched Paths (LSPs), across the network. Provision for backing up multiple service categories containing different forwarding and drop priorities, is also available with this technology. Multiprotocol label switching addresses common networking problems such as scalability, speed, Quality of Service (QoS), and traffic engineering, and provides them a viable and effective solution. Owing to its versatility, MPLS has emerged as a solution capable of meeting bandwidth and other service requirements for IP-based networks. Scalability and Routing -based issues can be resolved by MPLS technology, which also has the capacity to exist over existing ATM and Frame relay networks. Considering the positive points and shortcomings of ATM, MPLS technologies were designed to provide more leverage to network engineers and to be deployed flexibly. The marketplace is constantly being replaced with new technologies and technology devices. MPLS came to the forefront when there was a requirement for a protocol that needs less overhead and at the same time provides connection oriented-services for frames of variable length. Technology such as ATM and frame relay has been replaced in many areas by MPLS technology, which combines many options to satisfy the MPLS has dispensed cell-switching and signaling protocol used by ATM. Concurrently, Multiprotocol label switching technology continues to maintain the traffic engineering and bandwidth control, which was popularized by ATM and frame relay in large-scale networks. Migration to MPLS technology is beneficial especially since the benefits of traffic management are important. Performance level increases and so does reliability.

Currently, MPLS is used in large "IP only" networks. It is mainly used for forwarding Ethernet traffic and IP datagrams. MPLS VPN (Virtual Private Network) and traffic engineering are the major application areas of MPLS technology. MPLS IP VPN, a layer 3 VPN technology, is used to check, classify, and monitor IP packets. It is based on the service provider, to secure overlay VPN solutions. MPLS IP VPN is distinguished for its flexibility in networking modes, and features such as network scalability, QoS and traffic engineering. Today's business operations employ diverse applications across the Wide Area Networks (WANs) and it is essential to manage and prioritize traffic over the networks securely. This necessitates the use of technology such as MPLS IP VPN, which is a proven method for traffic engineering and network security.

A. *Label switch router (LSR):*

It refers to any router that has awareness of MPLS labels. The entry and exit routers of an MPLS network are called edge LSR (or label edge routers – LER), which, respectively, inject (push) an MPLS label onto an incoming packet (label assignment) and remove (pop) it off the outgoing packet (label removal). An edge LSR is often a high-speed router device in the core of an MPLS network that participates in the establishment of Label Switched Paths (LSP) using the appropriate label signaling protocol and high-speed switching of the data traffic based on the established paths.

B. *Label Switched Path (LSP):*

It is path defined by labels assigned between end points. An LSP can be dynamic or static. Dynamic LSPs are provisioned automatically using routing information. Static LSPs are explicitly provisioned.

C. *Label Virtual Circuit (LVC):*

It is a hop-by-hop connection established at the ATM transport layer 0 implement an LSP.

D. *LFIB:*

Used by the core MPLS routers (which are not ingress and egress MPLS routers). They compare the label in the incoming packet with the label they have in their LFIB. If a match is found, the routers forward that packet based on that match. If not, the packet will be dropped.

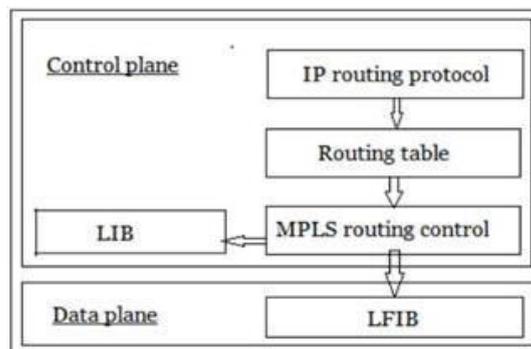


Fig. 1: Planes of router

E. *Label Distribution Protocol (LDP):*

It communicates labels and their meaning among LSRs. It assigns labels in edge and core devices to establish LSPs in conjunction with routing protocols such as Open Shortest Path First (OSPF), Intermediate System to Intermediate System (IS-IS), Routing Information Protocol (RIP), Enhanced Interior Gateway Routing Protocol (EIGRP), or Border Gateway Protocol (BGP)[3].

F. *IP Based Routing*

In traditional IP routing, each router in the network has to make independent routing decisions for each incoming packet. When a packet arrives at a router, the packet is stored in data plane of router. Each port of router is in its data plane. Now first layer 2 processing will be done on packet to check whether the packet is destined for that particular MAC of router. If yes then now layer 3 processing of packet is performed. Layer 3 process will check routing table, which is in control plane, the router to find the next hop for that packet based on the packets destination address in the packets IP header (longest match prefix lookup). Each router runs IP routing protocols like Border Gateway Protocol (BGP), Open Shortest Path First (OSPF) or Intermediate System-to- Intermediate System (IS-IS) to build the routing

table. Now if next hop is available then again layer 2 processing will be done to change the destination MAC of the packet and then the packet is forwarded to required port. Now routing table, layer 2 and processes are present in control plane of router. Each time for each packet which in data plane, each router performs the same steps of finding the next hop for the packet. The main issue with conventional routing protocols is that for entire decision making process, there will be transfer of processing from control plan to data plan many times. So this is time consuming process. Also IP routing is performed at each hop of the packets path in the network. Entire IP header analysis is done at each hop which is time consuming[2].

### III. MPLS BASED IP ROUTING

Multiprotocol label switching (MPLS) is an addition to the existing Internet Protocol (IP) architecture. By adding new capabilities to the IP architecture, MPLS enables support of new features and applications. In MPLS short fixed-length labels are assigned to packets at the edge of the MPLS domain and these pre assigned labels are used rather than the original packet headers to forward packets on pre-routed paths through the MPLS network.

In MPLS, the route the packet is forwarded through the MPLS domain is assigned only once i.e., when the packet enters the domain. Before a router forwards a packet it changes the label in the packet to a label that is used for forwarding by the next router in the path. MPLS unicast IP forwarding logic forwards packets based on the labels, however when choosing the exit interfaces, MPLS considers only the routes in the unicast IP routing table. This results in the packet flows over the same path as it would have even if MPLS was not used. Using MPLS labels does not add any benefit by itself, but it essentially enables the MPLS traffic engineering in an MPLS network, and therefore a critical feature of the MPLS.

MPLS still requires the use of control plane protocols such as OSPF and LDP to learn the labels and relate those labels to particular destination prefixes for building correct forwarding tables. MPLS also requires a fundamental change to the data plane's core forwarding logic, it defines a completely different packet-forwarding logic. In an MPLS network, the hosts should not send and receive labeled packets. All labeled packets are only for the routing and only routers should be sending and receiving the labeled packets in an MPLS network. Here when packet arrives at a router, it is stored in data plane. Now to take forwarding decision, router refers the LFIB table which in data plane itself. Now decision will be done on basis of LFIB and taken in data plane only. Labeled packet is switched to required port and as it does not involve processing of control plane, the process is faster.

The principal difference between a lookup in the routing table and the MPLS LFIB is that the routing table lookup is concerned with longest prefix match, i.e. having potentially many (imprecise) matches and selecting the one that most closely resembles the destination IP address. On the other hand, the MPLS LFIB always performs lookups on fixed-length values and with equality operation, not with prefix-based logic. Hence, at least in theory, a routing table lookup is algorithmically more complex than a lookup in the LFIB, as finding a longest prefix match is more computationally intensive than simply finding a single matching value. Therefore the LFIB lookups should be faster. MPLS forwards packets based on the MPLS labels, instead of using the packet's destination IP address.

Advantage of using labels and not the destination IP address is that packet forwarding decision can be made on the other factors such as traffic engineering and QoS requirements. In MPLS the first device does a routing lookup, just like in traditional IP routing. But instead of finding a next-hop, it finds the final destination router. And it finds a pre-determined path from current router to that final router. The router applies a "label" (or "shim") based on this information. Future routers use the label to route the traffic without needing to perform any additional IP lookups. At the final destination router, the label is removed and the packet is delivered via normal IP routing. Therefore in an MPLS network, data packets are assigned labels. Packet-forwarding decisions are made solely on the contents of this label, without the need to examine the packet itself. FEC is forward equivalence class which means providing the type of behavior to reach the destination. Whatever is the type of traffic (unicast or multicast), the mechanism used and forwarding algorithm used to take decision is same. Hence MPLS is faster.

IV. EXPERIMENTAL SETUP

It is important to know before the configuration it has no connection to customer sites and is deployed only in ISP networks.

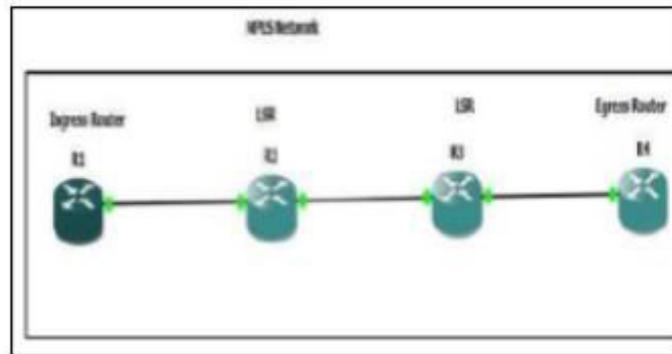


Fig. 2: MPLS Network in MPLS-VPN

MPLS based VPN has great importance in recent years. MPLS is technology used in WAN. It is deployed by ISPs in their cloud. IT has no direct linkage with the customer’s network. MPLS VPN is a VPN network construction based on the MPLS core network[6]. A MPLS based VPN is the implementation of VPN using the MPLS cloud. All the customer sites communicate with each other using the MPLS enabled provider network. MPLS label make a tunnel in this scenario.

The configuration is carried out on the Graphical Network Simulator-3 (GNS3). It is a GUI- based open source network simulator. The task is implemented in a cisco environment. The scenario is in figure 4.

Routers; Cisco 3745

IOS Version: 12.4

Router R1- R5 constitutes the MPLS network. It is also called the provider’s network. MPLS is running on this network. In the context of MPLS VPN, routers R1, R2 and R5 are called Provider Edge (PE) routers. They are the devices that have direct connectivity with a customer’s network.

redistribute the routes from IGP to MP-BGP to share the VPN routes among the PE devices. This increases the complexity in configuration at PE devices. Hence, BGP is used because it shares the routes by default with MP-BGP and no routes re- distribution is required.

Provider network OSPF is configured as routing protocol in the provider network. The MPLS is enabled on all provider network routers. MPLS labels are assigned based on all provider network routers. MPLS labels are assigned based on routers of OSPF MPLS doesn’t work without a routing protocol in a network. It can work with any IGP running in the network.

MP-BGP session; It is possible that some VPN have exactly the same IP address. To overcome this problem, VPNv4 addresses are used. In vpnv4 RD is added to the IP address to make a unique 96 bit long address. But the issue arising is that it no longer remains an IPV4 or IPV6 address. A normal routing protocol cannot carry this routing information. Hence, MP-BGP is used to carry the VPNv\$ addresses to other PE devices. In this scenario, MP-BGP sessions are established from Router R1 and Router R2 to Router R5. As it hub and spoke topology, We don’t need MP- BGP connectivity between R1 and R2. This situation is shown in figure 4.

Whereas routers R6-R9 are called Customer Edge (CE) routers. They are gateways of customer’s network and only device having connectivity with an ISP’s network. The whole customer’s network is called C-network.

Configuration at CE devices; At CE devices no special configuration is required. The only requirement is to assign IP addresses to interfaces and enable any IGP to carry the customer routes to connected PE devices.

Configurations at PE devices; In the context of MPLS VPN, most important configuration are done in PE devices. All the parameters should be configured carefully to establish the VPN connectivity. One of the most important parameters

is the configuration of virtual Routing and Forwarding (VRF) instances. Inside, VRFs Route Distinguishers (RD) and route targets (export/imports) are defined.

RD is the unique ID given to a particular VPN site. It must be unique in the whole network, as a customer site is defined based on RD. It is a 64 bit long address and mainly has three formats which are used to assign RDs to a customer site by ISP's as shown in figure 5. To established connectivity to a particular customer site, route targets exposed from one VRF must be imported into the VRF of another customer site and Vice versa.

PE-CE routing; PE-CE routing that achieved by using a BGP protocol. Any another Interior Gateway Protocol (IGP) like RIP, EIGRP or static routing can be used instead of BGP. If we use any another IGP, then we have to

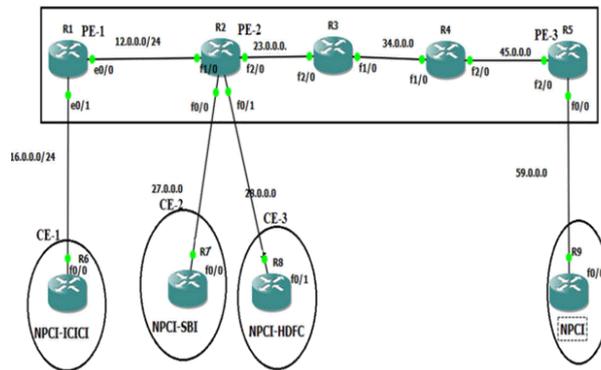


Fig. 4: Topology of MPLS-VPN and Traditional IP Network Our focus for this experimental setup is to analyze

network behaviour in congestion with different traffic flows[1]. We test IP(OSPF) and MPLS based network with following parameters,

- Latency in the network
- Utilization

A. Latency in the Network:

Latency and throughput are the two most fundamental measures of network performance. They are closely related, but whereas latency measures the amount of time between the start of an action and its completion, throughput is the total number of such actions that occur in a given amount of time. Latency is a networking term to describe the total time it takes a data packet to travel from one node to another. In other contexts, when a data packet is transmitted and returned back to its source, the total time for the round trip is known as latency.

B. Utilization:

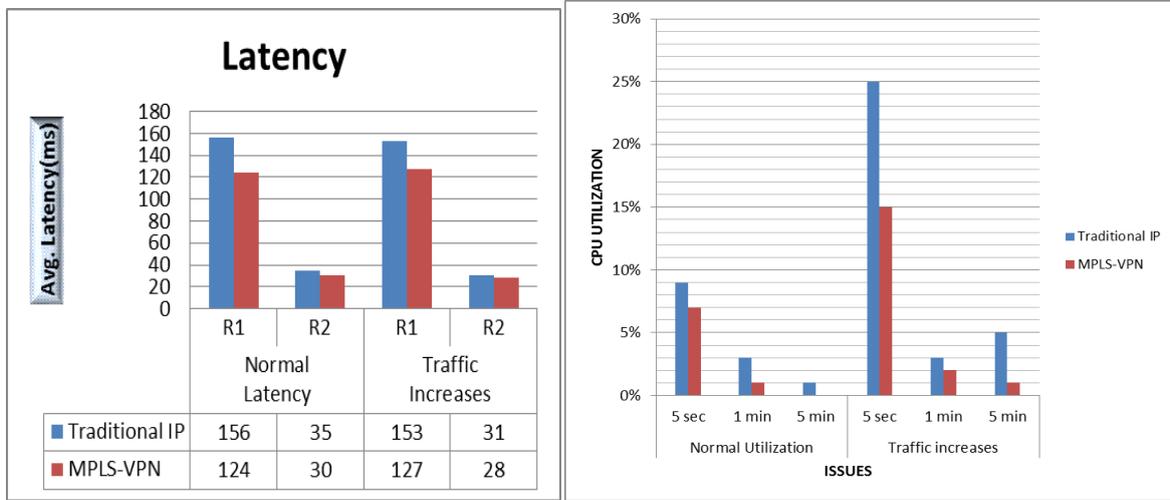
Network utilization is the amount of traffic on the network compared to the peak amount that the network can support. This is generally specified as a percentage. There are various times throughout the normal course of business when a network is busier, i.e., the network utilization is high. As a result, users experience a slow down when the network utilization is high enough. Response times grow greater than expectations preventing normal business processes from operating efficiently. Performance degradations are generally a nuisance but can become significant enough to result in lost revenues. It is important to understand the factors that can cause high network utilization and how to manage the network preventing it from negatively impacting the business.

V. EXPERIMENTAL RESULT

As per above readings, When traffic increases latency also increases in traditional IP routing and it effects on performance. In MPLS-VPN traffic increases in latency not that much varied as compared to normal latency.

| Issues         | Normal Latency |    | Traffic Increases |    |
|----------------|----------------|----|-------------------|----|
| Router         | R1             | R2 | R1                | R2 |
| Traditional IP | 156            | 35 | 153               | 31 |
| MPLS-VPN       | 124            | 30 | 127               | 28 |

Table 1: Latency Readings



**VI. CONCLUSION**

This paper has been prepared based on the traffic flow over both conventional and MPLS network, where network topology and other experimental parameters are chosen as common to establish the performance of MPLS network over traditional network.

Based on the comparison of MPLS and OSPF protocol OSPF chooses next hop on the basis of bandwidth as a cost of network. As higher is the bandwidth lower is the cost and the lower cost path is preferred. In the MPLS L3 VPN case between Hub and spoke, OSPF run as IGP (Interior gateway protocol). So, based on the comparison of signaling protocols, it can be found that using additional features of MPLS like MPLS TE with RSVP or CR-LDP protocol we can increase the network performance by diverting roots of different traffic flows and by setting traffic flows to different paths.

| Issues          | Normal Utilization |       |       | Traffic increases |       |       |
|-----------------|--------------------|-------|-------|-------------------|-------|-------|
| Time stamp      | 5 sec              | 1 min | 5 min | 5 sec             | 1 min | 5 min |
| Traditi onal IP | 9%                 | 3%    | 1%    | 25%               | 3%    | 5%    |
| MPLS- VPN       | 7%                 | 1%    | 0%    | 15%               | 2%    | 1%    |

Table 2: Utilization Reading

As per above reading of utilization, we conclude that after increasing the traffic in MPLS-VPN utilization is still normal.

The results are obtained after some experimentation and calculation with network scale (number of nodes, link capacity and delay) and traffic arrangements (sources and packet sizes, and rates). As expected, packet transmissions (in terms of both latency and loss) are improved in MPLS network. Throughput is also increased in MPLS enabled network. Although the chosen parameters can be disputed depends up traffic congestion to its extreme, the traffic engineering mechanism and setting MPLS experimental bits can enhance the performance of the service provider network.

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