



# Operations on Digraphs and Digraph Folding

E. EL Kholy, H Ahmed

Department of Mathematics, Faculty of Science, Tanta University, Tanta, Egypt

Department of Mathematics, Faculty of Shoubra Engineering, Banha University, Banha, Egypt

**ABSTRACT:** In this paper we examined the relation between digraph folding of a given pair of digraphs and digraph folding of new digraphs generated from these given pair of digraphs by some known operations like union, intersection, joins, Cartesian product and composition. We first redefined these known operations for digraphs, then we defined some new maps of these digraphs and we called these maps union, intersection, join, Cartesian and composition dimaps. In each case we obtained the necessary and sufficient conditions, if exist, for a dimap to be digraph folding. Finally we explored the digraph folding, if there exist any, by using the adjacency matrices.

**KEY WORDS:** Digraphs, adjacency matrices, digraph folding, union, intersection, join the Cartesian product and the composition of digraphs

## I. INTRODUCTION

Graph folding is introduced by E.EL-Kholy and A.AL-Esway [3]. The notion of digraph folding is introduced by E.EL-Kholy and H.Ahmed [4]. Definitions (1.1)

(1) A digraph  $D$  consists of elements, called vertices and a list of ordered pairs of these elements, called arcs. The set of vertices is called the vertex set of  $D$ , denoted by  $V(D)$ , and the list of arcs is called the arc list of  $D$ , denoted by  $A(D)$ . If  $v$  and  $w$  are vertices of  $D$ , then an arc of the form  $vw$  is said to be directed from  $v$  to  $w$ . The digraph with no loops is called simple. Two or more arcs joining the same pair of vertices in the same direction is called multiple arcs [5].

(2) Let  $D_1$  and  $D_2$  be digraphs and  $f: D_1 \rightarrow D_2$  a continuous function. Then  $f$  is called a digraph map if,

(i) For each vertex  $v \in V(D_1)$ ,  $f(v)$  is a vertex in  $V(D_2)$ .

(ii) For each arc  $e \in A(D_1)$ ,  $\dim(f(e)) \leq \dim(e)$  [4].

(3) Let  $D_1$  and  $D_2$  be simple digraphs, we call a digraph map  $f: D_1 \rightarrow D_2$  a digraph folding if  $f$  maps vertices to vertices and arcs to arcs, i.e., for each  $v \in V(D_1)$ ,  $f(v) \in V(D_2)$  and for each  $e \in A(D_1)$ ,  $f(e) \in A(D_2)$ . (4) If the digraph contains loops, then the digraph folding must send loops to loops but of the same direction. The set of digraph folding between digraphs  $D_1$  and  $D_2$  is denoted by  $\mathcal{D}(D_1, D_2)$  and from  $D$  into itself by  $\mathcal{D}(D)$ .

(5) Let  $D$  be a digraph without loops, with  $n$  vertices labeled  $1, 2, 3, \dots, n$ . The adjacency matrix  $M(D)$  is the  $n \times n$  matrix in which the entry in row  $i$  and column  $j$  is the number of arcs from vertex  $i$  to vertex  $j$  [5].

### A. Proposition

Let  $D$  be a connected digraph without loops with  $n$  vertices. Then a digraph folding of  $D$  into itself may be defined, if there is any, as a digraph map  $f$  of  $D$  to an image  $f(D)$  by mapping:

(i) The multiple arcs into one of its arcs.

(ii)(a) The vertex  $v_i$  to the vertex  $v_j$  if the numbers appearing in the adjacency matrix in the  $i^{\text{th}}$  and  $j^{\text{th}}$  rows (or columns) are the same.

(b) The vertex  $v_i$  to the vertex  $v_j$  if the entries of the  $i^{\text{th}}$  and  $j^{\text{th}}$  rows are zeros and if the  $i^{\text{th}}$  and  $j^{\text{th}}$  columns are the same, or there exists a row  $k$  which has numbers 1 in the  $i^{\text{th}}$  and  $j^{\text{th}}$  columns.

(iii)(a) The arc  $(v_i, v_k)$  to the arc  $(v_j, v_k)$  if the  $i^{\text{th}}$  and  $j^{\text{th}}$  rows (or columns) are the same.

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(b) The arc  $(v_i, v_j)$  to the arc  $(v_i, v_k)$  if the  $j^{\text{th}}$  and  $k^{\text{th}}$  columns (or rows) are the same. In general the arc  $(v_i, v_j)$  will be mapped to the arc  $(v_k, v_i)$  if  $v_i$  mapped to  $v_k$  and  $v_j$  mapped to  $v_i$ , [4].

## II. UNION OF DIGRAPHS

In the following we redefine the known operation, union, given for two simple graphs [3], for digraphs.

Definition (2-1)

Let  $D_1=(V_1,A_1)$  and  $D_2=(V_2,A_2)$  be simple digraphs. Then the simple digraph  $D=(V,A)$  where  $V=V_1 \cup V_2$  and  $A=A_1 \cup A_2$  is called the union of digraphs  $D_1$  and  $D_2$  and is denoted by  $D_1 \cup D_2$ . When  $D_1$  and  $D_2$  are vertex disjoint  $D_1 \cup D_2$  is denoted by  $D_1 + D_2$ , and is called the sum of digraphs  $D_1$  and  $D_2$ .

Definition (2-2)

Let  $D_1=(V_1,A_1)$  and  $D_2=(V_2,A_2)$  be simple digraphs. Let  $f: D_1 \rightarrow D_1$  and  $g: D_2 \rightarrow D_2$  be digraph maps. By the union dimap of the digraph maps  $f$  and  $g$ ,  $f \cup g$ , we mean a digraph map from the digraph  $D_1 \cup D_2$  into itself.  $f \cup g$  is defined by

$f \cup g: D_1 \cup D_2 \rightarrow D_1 \cup D_2$  such that  $(f \cup g)(v) = g(v)$ , for all  $v \in V_1 \cap V_2$ ,  $(f \cup g)(e) = f(e)$ , for all  $e \in A_1 \cap A_2$  defined by

$$(i) \text{ For each } v \in V_1 \cup V_2, (f \cup g)(v) = \begin{cases} f(v), & \text{if } v \in V_1 \\ g(v), & \text{if } v \in V_2 \end{cases}$$

$$(ii) \text{ For each } e \in A_1 \cup A_2, (f \cup g)(e) = \begin{cases} f(e), & \text{if } e \in A_1 \\ g(e), & \text{if } e \in A_2 \end{cases}$$

A. Theorem

Let  $D_1=(V_1,A_1)$  and  $D_2=(V_2,A_2)$  be simple connected digraphs. Let  $f: D_1 \rightarrow D_1$  and  $g: D_2 \rightarrow D_2$  be digraph maps. Then the union dimap  $f \cup g$  is a digraph folding if  $f$  and  $g$  are digraph foldings. In this case  $(f \cup g)(D_1 \cup D_2) = f(D_1) \cup g(D_2)$ . The proof is almost as in [2]. Example 2.4

Let  $D_1=(V_1,A_1)$ , where  $V_1=\{v_1, v_2, v_3, v_4\}$  and  $A_1=\{e_1, e_2, e_3, e_4, e_5\}$ .

Let  $D_2=(V_2,A_2)$ , where  $V_2=\{v_1, v_2, v_3, v_5\}$  and  $A_2=\{e_1, e_2, e_6\}$ , see Figure1.

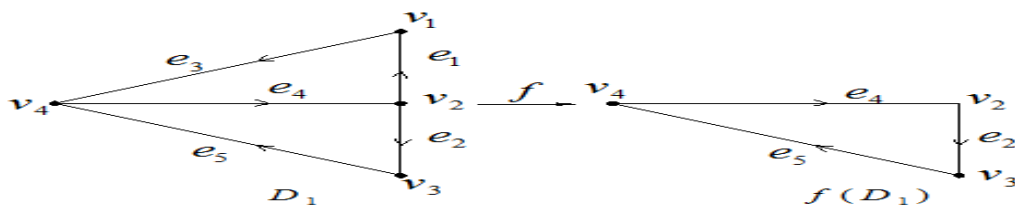
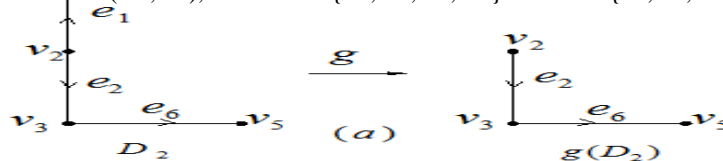


Figure 1:  $D_1=(V_1,A_1)$ , where  $V_1=\{v_1, v_2, v_3, v_4\}$  and  $A_1=\{e_1, e_2, e_3, e_4, e_5\}$

$D_2=(V_2,A_2)$ , where  $V_2=\{v_1, v_2, v_3, v_5\}$  and  $A_2=\{e_1, e_2, e_6\}$



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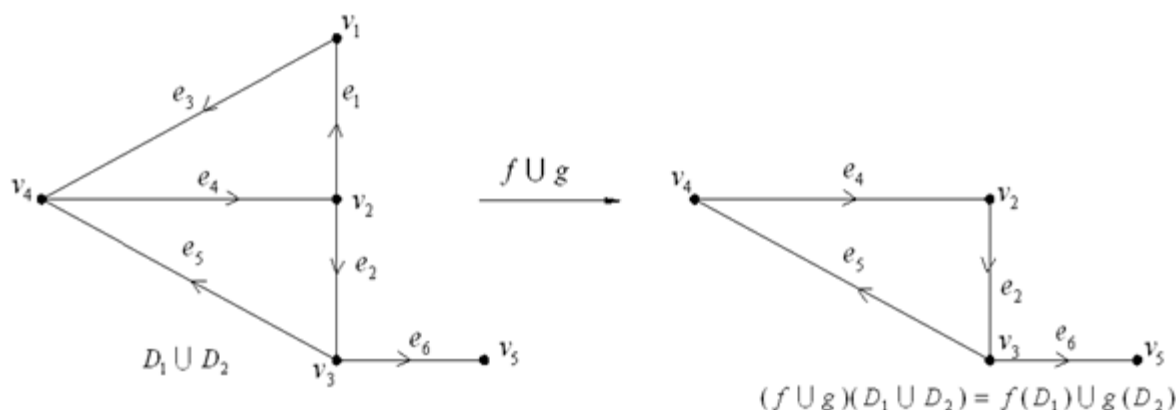


Figure 2:  $f \cup g: D_1 \cup D_2 \rightarrow UD_2$

Now let  $f \in \mathcal{D}(D_1)$  be a digraph folding defined by  $f\{v_1\} = \{v_3\}$  and  $f\{e_1, e_3\} = \{e_2, e_5\}$ , where through this paper the omitted vertices and arcs will be mapped to themselves. Also, let  $g \in \mathcal{D}(D_2)$  be a digraph folding defined by  $g\{v_1\} = \{v_3\}$  and  $g\{e_1\} = \{e_2\}$ , see Figure1. The union  $\text{map} f \cup g: D_1 \cup D_2 \rightarrow D_1 \cup D_2$  defined by  $(f \cup g)\{v_1\} = \{v_3\}$  and  $(f \cup g)\{e_1, e_3\} = \{e_2, e_5\}$  is a digraph folding, see Figure2. The adjacency matrices of  $D_1, D_2$  and  $D_1 \cup D_2$  are as follows:

$$M(D_1) = \begin{matrix} v_1 \\ v_2 \\ v_3 \\ v_4 \end{matrix} \begin{bmatrix} v_1 & v_2 & v_3 & v_4 \\ 0 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \end{bmatrix}, M(D_2) = \begin{matrix} v_1 \\ v_2 \\ v_3 \\ v_5 \end{matrix} \begin{bmatrix} v_1 & v_2 & v_3 & v_5 \\ 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \text{ and } M(D_1 \cup D_2) = \begin{matrix} v_1 \\ v_2 \\ v_3 \\ v_4 \\ v_5 \end{matrix} \begin{bmatrix} v_1 & v_2 & v_3 & v_4 & v_5 \\ 0 & 0 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

By using only these adjacency matrices we can define the digraph folding. For example, by using the adjacency matrix  $M(D_1)$  we can easily see that the vertex  $v_1$  will be mapped to the vertex  $v_3$  since the first and third rows of  $M(D_1)$  have the same entries. Also the arc  $(v_1, v_4) = e_3$  will be mapped to the arc  $(v_3, v_4) = e_5$  since the 1<sup>st</sup> and 3<sup>rd</sup> rows are the same, finally the arc  $(v_2, v_1) = e_1$  will be mapped to the arc  $(v_2, v_3) = e_2$  since the 1<sup>st</sup> and 3<sup>rd</sup> columns are the same. Again by using  $M(D_2)$  and  $M(D_1 \cup D_2)$  we can describe the digraph folding of both  $D_2$  and  $D_1 \cup D_2$ .

### III. INTERSECTION OF DIGRAPHS

#### A. Definition

Let  $D_1 = (V_1, A_1)$  and  $D_2 = (V_2, A_2)$  be simple digraphs. Then the simple digraph  $D = (V, A)$  where  $V = V_1 \cap V_2$  and  $A = A_1 \cap A_2$  is called the intersection of digraphs  $D_1$  and  $D_2$  and is denoted by  $D_1 \cap D_2$ .

#### B. Definition

Let  $D_1 = (V_1, A_1)$  and  $D_2 = (V_2, A_2)$  be simple digraphs. Let  $f: D_1 \rightarrow D_1$  and  $g: D_2 \rightarrow D_2$  be digraph maps. If  $f$  and  $g$  agree on  $V_1 \cap V_2$  and  $A_1 \cap A_2$  then by the intersection  $\text{map}$  of the digraph maps  $f$  and  $g$ ,  $f \cap g$ , we mean a digraph map  $f \cap g: D_1 \cap D_2 \rightarrow D_1 \cap D_2$ , where  $V_1 \cap V_2 \neq \emptyset$  defined by:

- (i) For all  $v \in V_1 \cap V_2$ ,  $(f \cap g)(v) = f(v) = g(v)$
- (ii) For all  $e \in A_1 \cap A_2$ ,  $(f \cap g)(e) = f(e) = g(e)$

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## THEOREM

Let  $D_1=(V_1,A_1)$  and  $D_2=(V_2,A_2)$  be simple connected digraphs .Let  $f:D_1 \rightarrow D_1$  and  $g:D_2 \rightarrow D_2$  be digraph maps .Then the intersection  $\text{dimap } f \cap g$  is a digraph folding if  $f$  and  $g$  are digraph folding. In this case  $(f \cap g)(D_1 \cap D_2) = f(D_1) \cap g(D_2)$ . The proof is easy.

let  $D_1=(V_1,A_1)$ , where  $V_1=\{ v_1,v_2,v_3,v_4,v_5,v_6 \}$  and  $A_1=\{e_1,e_2,e_3,e_4,e_5,e_6,e_7,e_8\}$ . Let  $D_2=(V_2,A_2)$ , where  $V_2=\{ v_1,v_5,v_6,v_7 \}$  and  $A_2=\{ e_5,e_8,e_9,e_{10},e_{11} \}$ , see Figure 3.

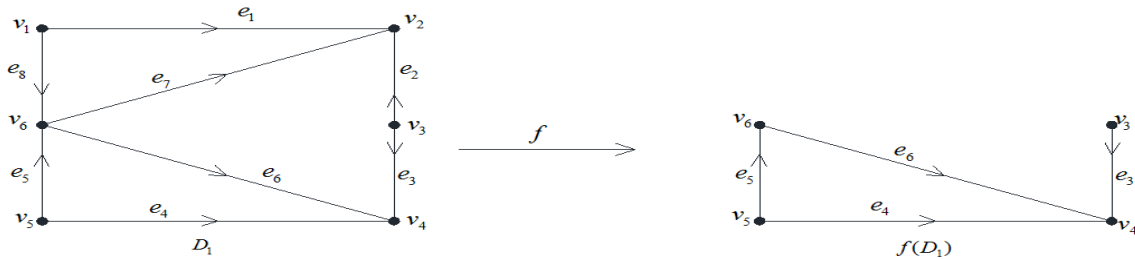


Figure 3:  $D_1=(V_1,A_1)$ , where  $V_1=\{ v_1,v_2,v_3,v_4,v_5,v_6 \}$  and  $A_1=\{e_1,e_2,e_3,e_4,e_5,e_6,e_7,e_8\}$

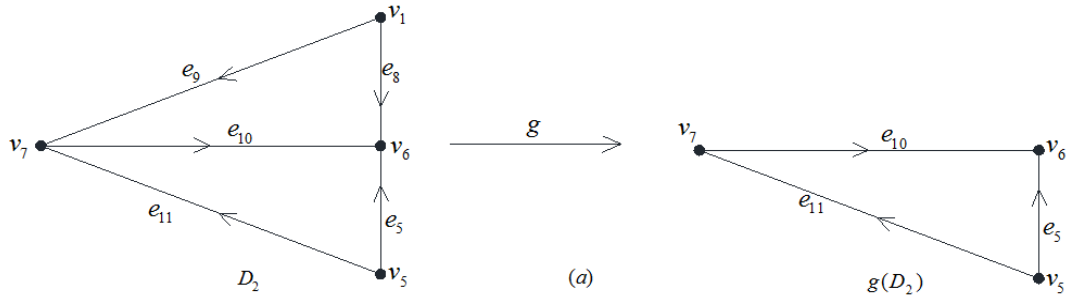


Figure 4:  $D_2=(V_2,A_2)$ , where  $V_2=\{ v_1,v_5,v_6,v_7 \}$  and  $A_2=\{ e_5,e_8,e_9,e_{10},e_{11} \}$

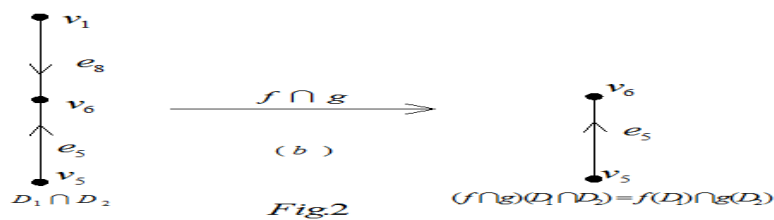


Figure 5:  $(f \cap g)\{v_1\}=\{v_5\}$  and  $(f \cap g)\{e_8\}=\{e_5\}$

Now let  $f \in \mathcal{D}(D_1)$  be a digraph folding defined by  $f\{v_1,v_2\} = \{v_5,v_4\}$  and  $f\{e_1,e_2,e_7,e_8\}=\{e_4,e_3,e_6,e_5\}$ . Also , let  $g \in \mathcal{D}(D_2)$  be a digraph folding defined by  $g\{v_1\}=\{v_5\}$  and  $g\{e_8,e_9\}=\{e_5,e_{11}\}$  , see Figure 1. The intersection  $\text{dimap } f \cap g : D_1 \cap D_2 \rightarrow D_1 \cap D_2$  defined by  $(f \cap g)\{v_1\}=\{v_5\}$  and  $(f \cap g)\{e_8\}=\{e_5\}$  is a digraph folding, see Figure 2. The adjacency matrices of  $D_1, D_2$  and  $D_1 \cap D_2$  are as follows:

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$$M(D_1) = \begin{matrix} v_1 \\ v_2 \\ v_3 \\ v_4 \\ v_5 \\ v_6 \end{matrix} \begin{bmatrix} v_1 & v_2 & v_3 & v_4 & v_5 & v_6 \\ 0 & 1 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 & 0 & 0 \end{bmatrix}, M(D_2) = \begin{matrix} v_1 \\ v_5 \\ v_6 \\ v_7 \end{matrix} \begin{bmatrix} v_1 & v_5 & v_6 & v_7 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix} \text{ and } M(D_1 \cap D_2) = \begin{matrix} v_1 \\ v_5 \\ v_6 \end{matrix} \begin{bmatrix} v_1 & v_5 & v_6 \\ 0 & 0 & 1 \\ 0 & 0 & 1 \\ 0 & 0 & 0 \end{bmatrix}$$

By using only these adjacency matrices we can define the digraph folding. For example, by using the adjacency matrix  $M(D_1 \cap D_2)$  we can easily see that the vertex  $v_1$  will be mapped to the vertex  $v_5$ , since the first and second rows have the same entries. Also, the arc  $(v_1, v_6) = e_8$  can be mapped to the arc  $(v_5, v_6) = e_5$ , since the first and second rows are the same. Again by using  $M(D_1)$  and  $M(D_2)$  we can describe the digraphfoldingof both  $D_1$  and  $D_2$ .

## IV. JOIN OF DIAGRAPHS

### A. Definition

Let  $D_1$  and  $D_2$  be vertex dis joint digraphs. Then we define the join digraph,  $D_1 \vee D_2$ , to be the digraph in which each vertex of  $D_1$  or  $D_2$  is adjacent to the vertices of  $D_2$  (or  $D_1$ ).

### B. Definition

Let  $D_1=(V_1,A_1)$ ,  $D_2=(V_2,A_2)$ ,  $D_3=(V_3,A_3)$  and  $D_4=(V_4,A_4)$  be simple digraphs. Let  $f: D_1 \rightarrow D_3$  and  $g: D_2 \rightarrow D_4$  be digraph maps. By a join dimap, we mean a digraph map,  $f \vee g: D_1 \vee D_2 \rightarrow D_3 \vee D_4$  defined by

- (i) For each vertex  $v \in V_1 \cup V_2$ ,  $(f \vee g)(v) = \begin{cases} f(v), & \text{if } v \in V_1 \\ g(v), & \text{if } v \in V_2 \end{cases}$
- (ii) For each arc  $e=(v_1,v_2), v_1 \in V_1$  and  $v_2 \in V_2$ ,  $(f \vee g)\{e\} = \{f(v_1), g(v_2)\} \in A_3 \vee A_4$ .
- (iii) If  $e=(u_1,v_1) \in A_1$ , then  $(f \vee g)\{e\}=(f \vee g)\{(u_1,v_1)\}=\{f(u_1), f(v_1)\}$ , Also if  $e=(u_2,v_2) \in A_2$ , then  $(f \vee g)\{e\}=(f \vee g)\{(u_2,v_2)\}=\{g(u_2), g(v_2)\}$

Note that If  $f\{u_1\} = f\{v_1\}$ , then the image of the join dimap  $(f \vee g)\{e\}$  will be a vertex of  $D_3 \vee D_4$  and thus is not a digraph folding.

### THEOREM

Let  $D_1, D_2, D_3$  and  $D_4$  be digraphs, let  $f: D_1 \rightarrow D_3$  and  $g: D_2 \rightarrow D_4$  be digraph maps. Then  $(f \vee g) \in \mathcal{D}(D_1 \vee D_2, D_3 \vee D_4)$  is a digraph folding if  $f$  and  $g$  are digraph folding.

Proof: Suppose  $f$  and  $g$  are digraph folding. Then  $(f \vee g)\{V_1 \cup V_2\} = \{f(V_1) \cup g(V_2)\}$ . But  $f(V_1) \in V(D_3)$ ,  $g(V_2) \in V(D_4)$ . Thus  $\{f(V_1) \cup g(V_2)\} \in V(D_3 \vee D_4)$ , i.e.,  $f \vee g$  maps vertices to vertices. Now, let  $e \in A(D_1 \vee D_2)$ . Then either  $e \in A(D_1)$  or  $e \in A(D_2)$  or  $e$  is an arc joining a vertex of  $D_1$  (or  $D_2$ ) to a vertex of  $D_2$  (or  $D_1$ ). In the first two cases and since each of  $f$  and  $g$  is a digraph folding,  $(f \vee g)\{e\} \in A(D_3 \vee D_4)$ . Now, if  $e=(v_1,v_2), v_1 \in D_1$  and  $v_2 \in D_2$ . Then  $(f \vee g)\{e\}=(f \vee g)\{(v_1,v_2)\}=\{f(v_1), g(v_2)\}=\{(v_3,v_4)\} \in A(D_3 \vee D_4)$ . Thus  $f \vee g$  maps arcs to arcs and hence the join digraph map is a digraph folding. The converse is guaranteed by the definition of the join digraph.

### C. Example

Let  $D_1=(V_1,A_1)$ , where  $V_1=\{v_1,v_2,v_3,v_4\}$  and  $A_1=\{e_1,e_2,e_3,e_4\}$  and  $D_2=(V_2,A_2)$ , where  $V_2=\{v_5,v_6,v_7\}$  and  $A_2=\{e_5,e_6\}$ , see Figure 3.

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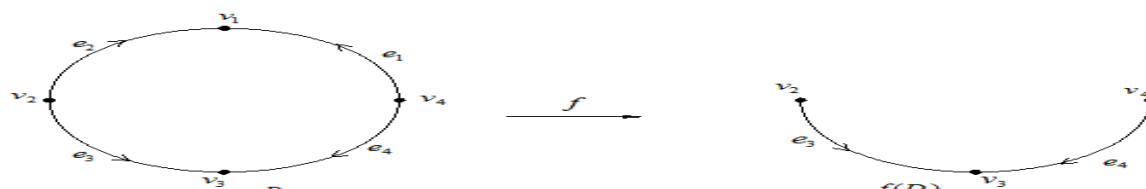


Figure 6:  $D_1=(V_1,A_1)$ , where  $V_1=\{ v_1,v_2,v_3,v_4,v_5,v_6 \}$  and  $A_1=\{e_1,e_2,e_3,e_4,e_5,e_6,e_7,e_8\}$

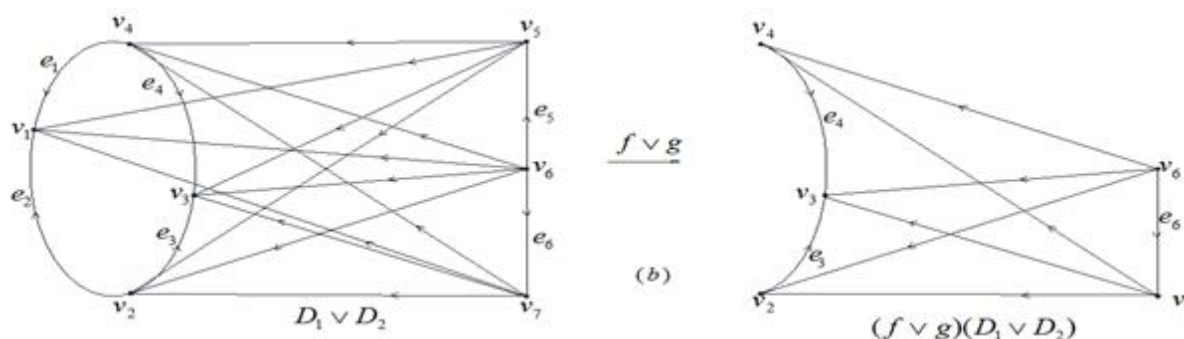


Figure 7:  $D_1= (V_1, A_1)$ , where  $V_1= \{v_1, v_2, v_3, v_4\}$  and  $A_1= \{e_1, e_2, e_3, e_4, e_5\}$   
 $D_2= (V_2, A_2)$ , where  $V_2= \{v_1, v_2, v_3, v_5\}$  and  $A_2= \{e_1, e_2, e_6\}$

Let  $f \in \mathcal{D}(D_1)$  be defined by  $f\{v_1\}=\{v_3\}$  and  $f\{e_1,e_2\}=\{e_4,e_3\}$ . Also, let  $g \in \mathcal{D}(D_2)$  be defined by  $g\{v_5\}=\{v_7\}$  and  $g\{e_5\}=\{e_6\}$ . The join dimap  $f \vee g: D_1 \vee D_2 \rightarrow D_1 \vee D_2$  is defined by  $(f \vee g)\{v_1,v_5\}=\{v_3,v_7\}$  and  $(f \vee g)\{e_1\}=(f \vee g)\{(v_4,v_1)\}=\{(v_4,v_3)\}=\{e_4\}$ , also,  $(f \vee g)\{e_5\}=(f \vee g)\{(v_6,v_5)\}=\{(v_6,v_7)\}=\{e_6\}$  and  $(f \vee g)\{(v_5,v_1)\}=\{(v_7,v_3)\}$ , and so on, see Figure 4. The adjacency matrices of  $D_1, D_2$  and  $D_1 \vee D_2$  are as follows:

$$M(D_1) = \begin{matrix} & \begin{matrix} v_1 & v_2 & v_3 & v_4 \end{matrix} \\ \begin{matrix} v_1 \\ v_2 \\ v_3 \\ v_4 \end{matrix} & \begin{bmatrix} 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 \end{bmatrix} \end{matrix}, M(D_2) = \begin{matrix} & \begin{matrix} v_5 & v_6 & v_7 \end{matrix} \\ \begin{matrix} v_5 \\ v_6 \\ v_7 \end{matrix} & \begin{bmatrix} 0 & 0 & 0 \\ 1 & 0 & 1 \\ 0 & 0 & 0 \end{bmatrix} \end{matrix} \text{ and } M(D_1 \vee D_2) = \begin{matrix} & \begin{matrix} v_1 & v_2 & v_3 & v_4 & v_5 & v_6 & v_7 \end{matrix} \\ \begin{matrix} v_1 \\ v_2 \\ v_3 \\ v_4 \\ v_5 \\ v_6 \\ v_7 \end{matrix} & \begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 1 & 0 & 0 & 0 \end{bmatrix} \end{matrix}$$

By using these adjacency matrices we can describe the digraph folding. The adjacency matrix  $M(D_1)$  suggests that the vertex  $v_1$  can be mapped to the vertex  $v_3$  since the first and third columns of  $M(D_1)$  have the same entries. Also the arc  $(v_4,v_1)=e_1$  can be mapped to the arc  $(v_4,v_3)=e_4$  and the arc  $(v_2,v_1)=e_2$  can be mapped to the arc  $(v_2,v_3)=e_3$  since the 1st and 3rd columns are the same. Again by using  $M(D_2)$  and  $M(D_1 \vee D_2)$  we can describe the digraph folding of both  $D_2$  and  $D_1 \vee D_2$ .

## V. THE CARTESIAN PRODUCT OF DIAGRAPH

### A. Definition

The Cartesian product  $D_1 \times D_2$  of two simple digraphs is a simple digraph with vertex set  $V(D_1 \times D_2) = V_1 \times V_2$  and arc set  $A(D_1 \times D_2) = [(A_1 \times V_2) \cup (V_1 \times A_2)]$  such that two vertices  $(u_1, u_2)$  and  $(v_1, v_2)$  are adjacent in  $D_1 \times D_2$  if, either

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- (i)  $u_1 = v_1$  and  $u_2$  is adjacent to  $v_2$  in  $D_2$ , or
- (ii)  $u_1$  is adjacent to  $v_1$  in  $D_1$ ,  $u_2 = v_2$ .

**B. Definition**

Let  $D_1, D_2, D_3$  and  $D_4$  be simple digraphs. Let  $f: D_1 \rightarrow D_3$  and  $g: D_2 \rightarrow D_4$  be digraph maps. Then by the Cartesian product  $\text{dimap} f \times g: D_1 \times D_2 \rightarrow D_3 \times D_4$ .

We mean a dimap defined as follows:

- (i) If  $v = (v_1, v_2) \in V_1 \times V_2$ ,  $v_1 \in V_1, v_2 \in V_2$ , then  $(f \times g)(v) = (f \times g)(v_1, v_2) = (f(v_1), g(v_2)) \in V_3 \times V_4$
  - (ii) If the arc  $e = ((v_1)_i, \{v_2\}_j), ((v_1)_i, \{v_2\}_k)$ , where  $\{v_1\}_i \in V(D_1)$  and  $\{v_2\}_j, \{v_2\}_k \in V(D_2)$ , then  $(f \times g)\{e\} = (f \times g)\{((v_1)_i, \{v_2\}_j), ((v_1)_i, \{v_2\}_k)\} = \{(f\{v_1\}_i, g\{v_2\}_j), (f\{v_1\}_i, g\{v_2\}_k)\}$ .
  - (iii) If the arc  $e = ((v_1)_i, \{v_2\}_j), ((v_1)_k, \{v_2\}_j)$ , where  $\{v_1\}_i, \{v_1\}_k \in V(D_1)$  and  $\{v_2\}_j \in V(D_2)$ , then  $(f \times g)\{e\} = (f \times g)\{((v_1)_i, \{v_2\}_j), ((v_1)_k, \{v_2\}_j)\} = \{(f\{v_1\}_i, f\{v_2\}_j), (f\{v_1\}_k, f\{v_2\}_j)\}$ .
- Note that if  $g\{v_2\}_j = g\{v_2\}_k$  or  $f\{v_1\}_i = f\{v_1\}_k$ , the image of the arc will be a vertex

**THEOREM**

Let  $D_1, D_2, D_3$  and  $D_4$  be digraphs, let  $f: D_1 \rightarrow D_3$  and  $g: D_2 \rightarrow D_4$  be digraph maps. Then  $(f \times g) \in \mathcal{D}(D_1 \times D_2, D_3 \times D_4)$  is a digraph folding iff  $f \in \mathcal{D}(D_1, D_3)$  and  $g \in \mathcal{D}(D_2, D_4)$  are digraph foldings. In this case  $(f \times g)(D_1 \times D_2) = f(D_1) \times g(D_2)$ .

**Proof:** Suppose  $f$  and  $g$  are digraph folding. Then for each vertex  $(v_1, v_2) \in (D_1 \times D_2) = V_1 \times V_2$ ,  $(f \times g)\{(v_1, v_2)\} = \{(f(v_1), g(v_2))\} = (v_3, v_4) \in V(D_1 \times D_2) = V_3 \times V_4$ , i.e., vertices to vertices. Now, let  $e \in A(D_1 \times D_2)$ , then if  $e = ((v_1)_i, \{v_2\}_j), ((v_1)_k, \{v_2\}_j)$ , where  $\{v_1\}_i$  is adjacent to  $\{v_1\}_k$  in  $D_1$  and  $\{v_2\}_j \in D_2$ , then  $(f \times g)\{e\} = (f \times g)\{((v_1)_i, \{v_2\}_j), ((v_1)_k, \{v_2\}_j)\} = \{(f\{v_1\}_i, f\{v_2\}_j), (f\{v_1\}_k, f\{v_2\}_j)\}$ , since  $\{v_1\}_i$  is adjacent to  $\{v_1\}_k$  and  $f$  is a digraph folding, then  $f\{v_1\}_i \neq f\{v_1\}_k$ . Thus  $(f \times g)\{e\} \in A(D_3 \times D_4)$ . By the same procedure, if  $e = ((v_1)_i, \{v_2\}_j), ((v_1)_i, \{v_2\}_k)$ , where  $\{v_1\}_i \in V(D_1)$  and  $\{v_2\}_j$  is adjacent to  $\{v_2\}_k$  in  $D_2$ , then  $(f \times g)\{e\} \in A(D_3 \times D_4)$ . i.e.,  $f \times g$  maps arcs to arcs and hence the Cartesian product dimap is a digraph folding. To prove the converse suppose that  $(f \times g)$  is a digraph folding and for  $g$ , is not a digraph folding. In this case  $f$  or  $g$ , will maps an arc to a vertex, say  $f\{(u_1, v_1)\} = \{u_3\} \in V(D_3)$ . Then  $(f \times g)\{((u_1, \{v_2\}_j), ((u_1, \{v_2\}_j))\} = \{(f\{u_1\}, \{v_2\}_j), (f\{u_1\}, \{v_2\}_j)\} = \{(u_3, \{v_2\}_j), (u_3, \{v_2\}_j)\} \in V(D_3 \times D_4)$ . This contradicts the assumption and thus each of  $f$  and  $g$  must be a digraph folding.

**A. Examples**

- (a) Let  $D_1 = (V_1, A_1)$ , where  $V_1 = \{u_1, u_2, u_3, u_4\}$ ,  $A = \{e_1, e_2, e_3, e_4\}$  and  $D_2 = (V_2, A_2)$ , where  $V_2 = \{v_1, v_2, v_3\}$ ,  $A = \{e_5, e_6\}$ , see Figure 4.

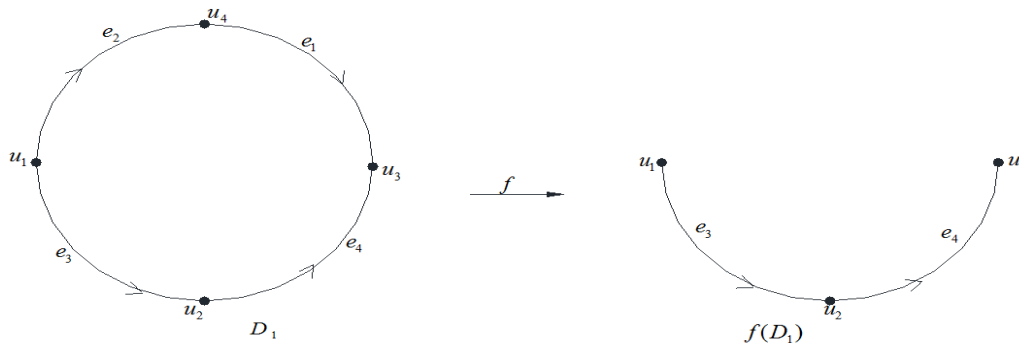


Figure 8:  $(f \times g)\{(u_4, v_1), (u_4, v_2), (u_4, v_3), (u_1, v_1), (u_2, v_1), (u_3, v_1)\} = \{(u_2, v_1), (u_2, v_2), (u_2, v_3), (u_1, v_3), (u_2, v_3), (u_3, v_3)\}$

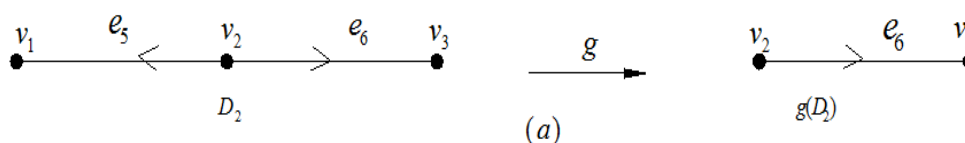


Figure 9:  $(f \times g)\{e_5\} = (f \times g)\{(v_1, v_2)\} = \{(v_2, v_2)\} = \{e_6\}$

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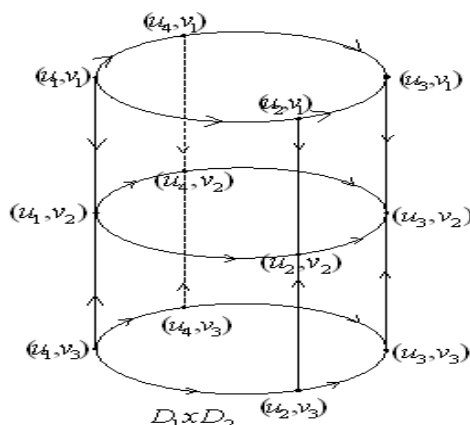


Figure 10:  $(fxg)\{(u_4,v_1),(u_3,v_1)\},\{(u_3,v_1),(u_3,v_2)\} = \{(u_2,v_3),(u_3,v_3)\},\{(u_3,v_3),(u_3,v_2)\}$

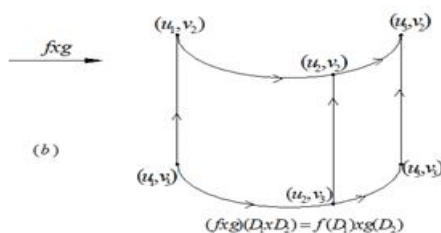


Figure 11:  $(fxg)\{(u_4,v_1),(u_3,v_1)\},\{(u_3,v_1),(u_3,v_2)\} = \{(u_2,v_3),(u_3,v_3)\},\{(u_3,v_3),(u_3,v_2)\}$

Let  $f \in D(D_1)$  defined by  $f\{u_4\} = \{u_2\}$  and  $f\{e_1, e_2\} = \{e_4, e_3\}$ . Also, let  $g \in D(D_2)$  defined by  $g\{v_1\} = \{v_3\}$  and  $g\{e_5\} = \{e_6\}$ . Then the Cartesian product  $\text{dimapfxg} : D_1 \times D_2 \rightarrow D_3 \times D_4$  is defined as follows :

$(fxg)\{(u_4,v_1),(u_4,v_2),(u_4,v_3),(u_1,v_1),(u_2,v_1),(u_3,v_1)\} = \{(u_2,v_1),(u_2,v_2),(u_2,v_3), (u_1,v_3),(u_2,v_3),(u_3,v_3)\}$ . Also,  $(fxg)\{(u_4,v_1),(u_3,v_1)\},\{(u_3,v_1),(u_3,v_2)\} = \{(u_2,v_3),(u_3,v_3)\},\{(u_3,v_3),(u_3,v_2)\}$  and so on, see Figure 4. The adjacency matrices of  $D_1$ ,  $D_2$  and  $D_1 \times D_2$  are as follows:

$$M(D_1) = \begin{matrix} u_1 \\ u_2 \\ u_3 \\ u_4 \end{matrix} \begin{bmatrix} u_1 & u_2 & u_3 & u_4 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}, M(D_2) = \begin{matrix} v_1 \\ v_2 \\ v_3 \end{matrix} \begin{bmatrix} v_1 & v_2 & v_3 \\ 0 & 1 & 0 \\ 0 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix} \text{ and}$$

$$M(D_1 \times D_2) = \begin{matrix} (u_1, v_1) \\ (u_2, v_1) \\ (u_3, v_1) \\ (u_4, v_1) \\ (u_1, v_2) \\ (u_2, v_2) \\ (u_3, v_2) \\ (u_4, v_2) \\ (u_1, v_3) \\ (u_2, v_3) \\ (u_3, v_3) \\ (u_4, v_3) \end{matrix} \begin{bmatrix} (u_1, v_1) & (u_2, v_1) & (u_3, v_1) & (u_4, v_1) & (u_1, v_2) & (u_2, v_2) & (u_3, v_2) & (u_4, v_2) & (u_1, v_3) & (u_2, v_3) & (u_3, v_3) & (u_4, v_3) \\ 0 & 1 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \end{bmatrix}$$



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Once again we can describe the digraph foldings by using  $M(D_1), M(D_2)$  and  $M(D_1XD_2)$ . For example, from  $M(D_1XD_2)$  we can see that the vertex  $(u_4, v_1)$  can be mapped to the vertex  $(u_2, v_1)$  since the second and fourth columns are the same. Also, the arc  $((u_1, v_1), (u_4, v_1))$  will be mapped to the arc  $((u_1, v_3), (u_2, v_3))$  since the vertex  $(u_4, v_1)$  is mapped to the vertex  $(u_2, v_3)$  and the vertex  $(u_1, v_1)$  is mapped to the vertex  $(u_1, v_3)$ , and so on, see Figure 4.

(b) Let  $D_1=(V_1, A_1)$ , where  $V_1=\{u_1, u_2, u_3, u_4\}$ ,  $A_1=\{e_1, e_2, e_3, e_4\}$  and  $D_2=(V_2, A_2)$ , where  $V_2 = \{v_1, v_2, v_3, v_4\}$ ,  $A_2 = \{e_5, e_6, e_7, e_8\}$ , see Figure 5.

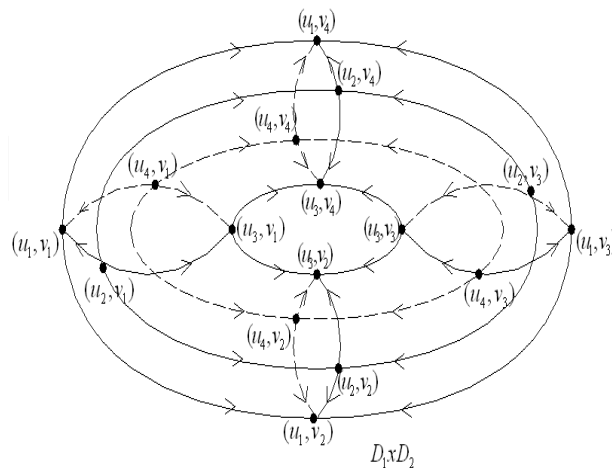
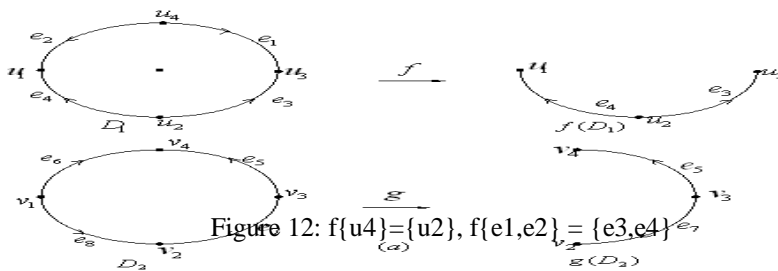


Figure 13:  $(fxg)\{((u_4, v_1), (u_3, v_1)), ((u_3, v_1), (u_3, v_2))\} = \{((u_2, v_3), (u_3, v_3)), ((u_3, v_3), (u_3, v_2))\}$

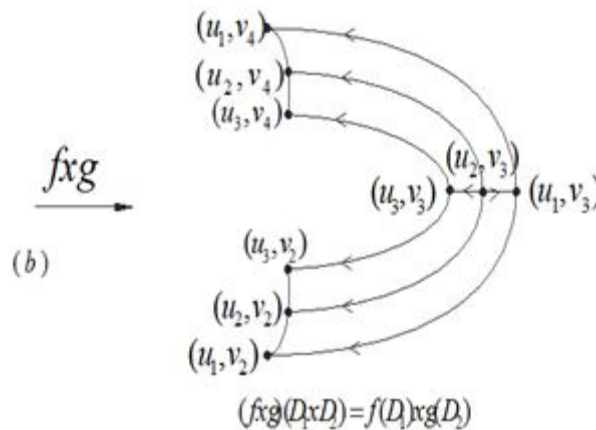


Figure 14:  $(fxg)\{((u_4, v_1), (u_3, v_1)), ((u_3, v_1), (u_3, v_2))\} = \{((u_2, v_3), (u_3, v_3)), ((u_3, v_3), (u_3, v_2))\}$

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Let  $f \in \mathcal{D}(D_1)$  be defined by  $f\{u_4\} = \{u_2\}$ ,  $f\{e_1, e_2\} = \{e_3, e_4\}$  and  $g \in \mathcal{D}(D_2)$  be defined by  $g\{v_1\} = \{v_3\}$ ,  $g\{e_6, e_8\} = \{e_5, e_7\}$ . Then the cartesian product  $\text{dimap } h = f \times g : D_1 \times D_2 \rightarrow D_1 \times D_2$  is defined as follows:  
 $h\{(u_4, v_2), (u_3, v_1)\} = \{(u_2, v_2), (u_3, v_3)\}$ , and so on. Also,  $h\{((u_4, v_2), (u_1, v_2)), ((u_3, v_1), (u_3, v_2))\} = \{((u_2, v_2), (u_1, v_2)), ((u_3, v_3), (u_3, v_2))\}$ , and so on, see Figure 6. The adjacency matrices of  $D_1, D_2$  and  $D_1 \times D_2$  are as follows:

$$M(D_1) = \begin{matrix} & u_1 & u_2 & u_3 & u_4 \\ \begin{matrix} u_1 \\ u_2 \\ u_3 \\ u_4 \end{matrix} & \begin{bmatrix} 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 \end{bmatrix} \end{matrix}, \quad M(D_2) = \begin{matrix} & v_1 & v_2 & v_3 & v_4 \\ \begin{matrix} v_1 \\ v_2 \\ v_3 \\ v_4 \end{matrix} & \begin{bmatrix} 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \end{matrix} \text{ and}$$

	$(u_1, v_1)$	$(u_2, v_1)$	$(u_3, v_1)$	$(u_4, v_1)$	$(u_1, v_2)$	$(u_2, v_2)$	$(u_3, v_2)$	$(u_4, v_2)$	$(u_1, v_3)$	$(u_2, v_3)$	$(u_3, v_3)$	$(u_4, v_3)$	$(u_1, v_4)$	$(u_2, v_4)$	$(u_3, v_4)$	$(u_4, v_4)$
$(u_1, v_1)$	0	0	0	0	1	0	0	0	0	0	0	0	1	0	0	0
$(u_2, v_1)$	1	0	1	0	0	1	0	0	0	0	0	0	0	1	0	0
$(u_3, v_1)$	0	0	0	0	0	0	1	0	0	0	0	0	0	0	1	0
$(u_4, v_1)$	1	0	1	0	0	0	0	1	0	0	0	0	0	0	0	1
$(u_1, v_2)$	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
$(u_2, v_2)$	0	0	0	0	1	0	1	0	0	0	0	0	0	0	0	0
$(u_3, v_2)$	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
$(u_4, v_2)$	0	0	0	0	1	0	1	0	0	0	0	0	0	0	0	0
$(u_1, v_3)$	0	0	0	0	0	1	0	0	0	0	0	0	1	0	0	0
$(u_2, v_3)$	0	0	0	0	0	0	1	0	0	1	0	1	0	0	1	0
$(u_3, v_3)$	0	0	0	0	0	0	1	0	0	0	0	0	0	0	1	0
$(u_4, v_3)$	0	0	0	0	0	0	0	1	1	0	1	0	0	0	0	1
$(u_1, v_4)$	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
$(u_2, v_4)$	0	0	0	0	0	0	0	0	0	0	0	0	1	0	1	0
$(u_3, v_4)$	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
$(u_4, v_4)$	0	0	0	0	0	0	0	0	0	0	0	0	1	0	1	0

Once again we can describe the digraph foldings by using  $M(D_1)$  and  $M(D_1 \times D_2)$ . For example, from  $M(D_1 \times D_2)$  we can see that the vertex  $(u_4, v_2)$  can be mapped to the vertex  $(u_2, v_2)$  since the 6<sup>th</sup> and 8<sup>th</sup> rows have the same entries. And the vertex  $(u_3, v_1)$  can be mapped to the vertex  $(u_3, v_3)$  since the 3<sup>rd</sup> and 11<sup>th</sup> rows are the same. Also, the arcs  $((u_4, v_2), (u_1, v_2))$  and  $((u_4, v_2), (u_3, v_2))$  can be mapped to the arcs  $((u_2, v_2), (u_1, v_2))$  and  $((u_2, v_2), (u_3, v_2))$ , respectively, since the 6<sup>th</sup> and 8<sup>th</sup> rows are the same. Finally the arcs  $((u_3, v_1), (u_3, v_2))$  and  $((u_3, v_1), (u_3, v_4))$  can be mapped to the arcs  $((u_3, v_3), (u_3, v_2))$  and  $((u_3, v_3), (u_3, v_4))$ , respectively, since the 3<sup>th</sup> and 11<sup>th</sup> rows are the same. And so on, see Figure 6.

## VI. THE COMPOSITION OF DIGRAPHS

### A. Definition

The composition  $D_1[D_2]$  of two simple digraphs is a simple digraph with  $V(D_1[D_2]) = V_1 \times V_2$ . The vertices  $u = (u_1, u_2)$  and  $v = (v_1, v_2)$  are adjacent if either  $u_1$  is adjacent to  $v_1$  and  $u_2 = v_2$  or  $u_1 = v_1$  and  $u_2$  is adjacent to  $v_2$ .

### B. Definition

Let  $D_1, D_2, D_3$  and  $D_4$  be simple digraphs. Let  $f: D_1 \rightarrow D_3$  and  $g: D_2 \rightarrow D_4$  be digraph maps. By the composition  $\text{dimap } f[g]: D_1[D_2] \rightarrow D_3[D_4]$  we mean a map defined as follows

(i) If  $v = (v_1, v_2) \in V(D_1[D_2]) = V_1 \times V_2$ , then  $f[g]\{(v_1, v_2)\} = \{(f(v_1), g(v_2))\} \in V(D_3[D_4])$

(ii) Let  $e = \{(\{v_1\}_i, \{v_2\}_j), (\{v_1\}_k, \{v_2\}_l)\}$ . If  $\{v_1\}_i = \{v_1\}_k$  and  $\{v_2\}_j$  is adjacent to  $\{v_2\}_l$ , then  $f[g]\{e\} = \{(\{v_1\}_i, g\{v_2\}_j), (\{v_1\}_i, g\{v_2\}_l)\}$ . Also, if  $\{v_2\}_j = \{v_2\}_l$  and  $\{v_1\}_i$  is adjacent to  $\{v_1\}_k$ , then  $f[g]\{e\} = \{(f\{v_1\}_i, \{v_2\}_j), (f\{v_1\}_k, \{v_2\}_j)\}$ .

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**THEOREM**

Let  $D_1, D_2, D_3$  and  $D_4$  be digraphs. Let  $f: D_1 \rightarrow D_2$  and  $g: D_2 \rightarrow D_4$  be digraph maps. Then the composition map  $f \circ g \in \mathcal{D}(D_1[D_2], D_3[D_4])$  is a digraph folding if  $f \in \mathcal{D}(D_1, D_3)$  and  $g \in \mathcal{D}(D_2, D_4)$  are digraph foldings.

Proof: Let  $f$  and  $g$  be digraph folding, then

(i) For each vertex  $v = (v_1, v_2) \in V(D_1[D_2]) = V_1 \times V_2$ ,  $f[g]\{(v_1, v_2)\} = \{(f(v_1), g(v_2))\}$ . But  $f(v_1) \in V(D_3)$  and  $g(v_2) \in V(D_4)$ , then  $\{(f(v_1), g(v_2))\} \in V(D_3[D_4])$ , i.e.,  $f[g]$  maps vertices to vertices.

(ii) Let  $e = \{(\{v_1\}_i, \{v_2\}_j), (\{v_1\}_k, \{v_2\}_l)\}$  and suppose  $\{v_1\}_i$  is adjacent to  $\{v_1\}_k$ , then there exists an arc  $\{(\{v_1\}_i, \{v_1\}_k)\} \in A_1$ , since  $f$  is a digraph folding and  $\{(\{v_1\}_i, \{v_1\}_k)\} \in A_1$ , then  $f[g]\{e\} \in A(D_3[D_4])$ . Now, if  $\{v_1\}_i = \{v_1\}_k$  and  $\{v_2\}_j$  is adjacent to  $\{v_2\}_l$ , then  $f[g]\{e\} = \{(\{v_1\}_i, g\{v_2\}_j), (\{v_1\}_i, g\{v_2\}_l)\}$ , since  $\{v_2\}_j$  is adjacent to  $\{v_2\}_l$ , then there exists an arc  $\{(\{v_2\}_j, \{v_2\}_l)\} \in A_2$  such that  $\{(g\{v_2\}_j, g\{v_2\}_l)\} \in A_3$ , i.e.,  $g\{v_2\}_j \neq g\{v_2\}_l$  and hence  $f[g]\{e\} \in A(D_3[D_4])$ , i.e.,  $f[g]$  maps arcs to arcs. The converse is not true since if  $f[g]$  is a digraph folding and for  $g$ , is not a digraph folding. In this case for  $g$ , maps an arc to a vertex, say  $f(u_1, v_1) = (u_3, u_3)$ ,  $u_3 \in V(D_3)$ .

Then  $f[g]\{(u_1, \{v_2\}_i), (v_1, \{v_2\}_j)\} = \{(f(u_1), g\{v_2\}_i), (f(v_1), g\{v_2\}_j)\} = \{(u_3, g\{v_2\}_i), (u_3, g\{v_2\}_j)\}$  which is an arc of  $D_3[D_4]$ .

*A. Example*

Let  $D_1, D_2, f$  and  $g$  be the digraphs and digraph foldings given in Example (A). The adjacency matrix of  $D_1[D_2]$  is as follows:

$$M(D_1[D_2]) = \begin{matrix} & \begin{matrix} (u_1, v_1) & (u_2, v_1) & (u_3, v_1) & (u_4, v_1) & (u_1, v_2) & (u_2, v_2) & (u_3, v_2) & (u_4, v_2) & (u_1, v_3) & (u_2, v_3) & (u_3, v_3) & (u_4, v_3) \end{matrix} \\ \begin{matrix} (u_1, v_1) \\ (u_2, v_1) \\ (u_3, v_1) \\ (u_4, v_1) \\ (u_1, v_2) \\ (u_2, v_2) \\ (u_3, v_2) \\ (u_4, v_2) \\ (u_1, v_3) \\ (u_2, v_3) \\ (u_3, v_3) \\ (u_4, v_3) \end{matrix} & \begin{bmatrix} 0 & 1 & 0 & 1 & 1 & 1 & 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 1 & 1 & 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 0 \end{bmatrix} \end{matrix}$$

Now a digraph folding  $f[g]: D_1[D_2] \rightarrow D_3[D_4]$  can be defined as follows:  $f[g]\{(u_4, v_1), (u_4, v_2), (u_4, v_3), (u_1, v_1), (u_2, v_1), (u_3, v_1)\} = \{(u_2, v_1), (u_2, v_2), (u_2, v_3), (u_1, v_3), (u_2, v_3), (u_3, v_3)\}$ . Also,  $f[g]\{((u_4, v_1), (u_3, v_1)), ((u_3, v_2), (u_3, v_1)), ((u_2, v_2), (u_3, v_1))\} = \{((u_2, v_1), (u_3, v_3)), ((u_3, v_2), (u_3, v_3)), ((u_2, v_2), (u_3, v_3))\}$ , and so on, see Figure 6.

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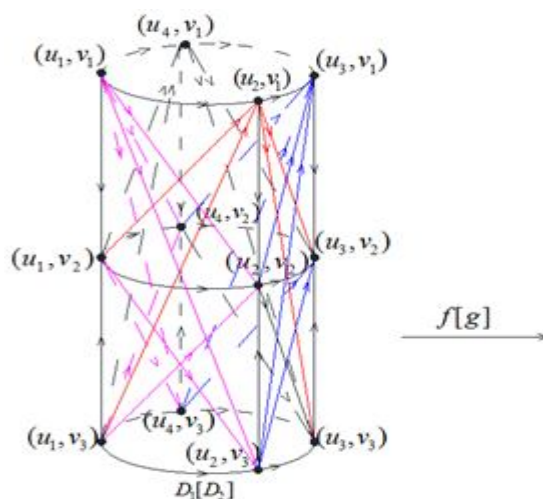


Figure 15:  $f[g]\{(u_4, v_1), (u_4, v_2), (u_4, v_3), (u_1, v_1), (u_2, v_1), (u_3, v_1)\} = \{(u_2, v_1), (u_2, v_2), (u_2, v_3), (u_1, v_3), (u_2, v_3), (u_3, v_3)\}$

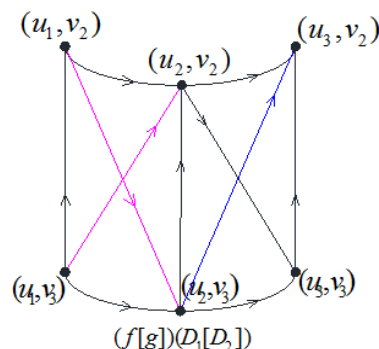


Figure 16:  $f[g]\{(u_4, v_1), (u_4, v_2), (u_4, v_3), (u_1, v_1), (u_2, v_1), (u_3, v_1)\} = \{(u_2, v_1), (u_2, v_2), (u_2, v_3), (u_1, v_3), (u_2, v_3), (u_3, v_3)\}$

We can describe the digraph foldings by using  $M(D_1), M(D_2)$  and  $M(D_1 [D_2])$ . For example, from  $M(D_1 [D_2])$  we can see that the vertex  $(u_4, v_1)$  can be mapped to the vertex  $(u_2, v_1)$  since the second and fourth rows have the same entries. Also, the arc  $((u_1, v_1), (u_4, v_1))$  can be mapped to the arc  $((u_1, v_1), (u_2, v_1))$  since the second and fourth rows are the same. Also the vertex  $(u_1, v_1)$  can be mapped to the vertex  $(u_1, v_3)$  since 1st and 9<sup>th</sup> rows have the same entries, and so on.

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