



# Securing from Untrusted Location Based Service Providers on Spatial Data Bases in K-Query Processing

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**ABSTRACT:** A well known distributed architecture for cooperative location-based data information and sharing which turn out to be progressively well known because of the high level development of Web skilled and location-based mobiles. The framework comprises of an information authority, information donors, location-based administration suppliers (LBSPs), and framework clients. The accumulated audits of information authority about points of interest (POIs) from information benefits, while LBSPs buy POI information sets from the information gatherer and permit clients to perform spatial top k inquiries which request the POIs in a sure locale and with the most noteworthy k appraisals for an intrigued POI trait. Practically, LBSPs are untrusted and may return fake question results for different awful thought processes, e.g., for POIs willing to pay. This paper presents three novel plans for clients to recognize fake spatial depiction and moving top-k question results as a push to cultivate the useful arrangement and utilization of the proposed framework. The effectiveness and productivity of these schemes are evaluated by various analytics.

**KEYWORDS:** Spatial top-k query, location-based service, security, POI, location based service providers

## I. INTRODUCTION

The explosive growth of Internet-capable location-aware cell phones and the surge in interpersonal organization use are Encouraging synergistic data era and sharing on an exceptional scale. All mobile phones have Wi-Fi Web get to and can simply get their exact locations through pre-introduced situating programming. Likewise attributable to the developing notoriety of informal communities, it is more advantageous and inspiring for versatile clients to impart to others their involvement with a wide range of purposes of intrigues. In the meantime, it gets to be regular spot for individuals to perform different spatial POI inquiries at online location-based administration suppliers (LBSPs). This paper concentrates on spatial top-k questions, and the expression "spatial" will be overlooked from now on for curtness.

## II. PROBLEM FORMULATION

This work is popularly called as information outsourcing, for which we can just review representative schemes because of space constraints. The way of information outsourcing was initially presented, in which an information proprietor outsources its information to an outsider administration supplier who is responsible of noting the information inquiries from either the information proprietor or different clients. As a rule, there are two security concerns in information outsourcing: information protection and question integrity. A bucketing methodology was proposed, to empower proficient reach inquiries over scrambled information, which was recently enhanced and the novel systems for multidimensional extent inquiries over encoded information.

### A. Existing System Model

The helpfulness and in the long run impede the more common utilization of spatial top-k question administrations. Proceed with the eatery case. The information sets at individual LBSPs may not cover all the Italian restaurants within a search. Also, the same restaurants may get assorted evaluations at distinctive LBSPs, so clients may get confounded by altogether different *First*, individual LBSPs regularly have little information sets including POI audits. This would



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to a great extent influence inquiry results from diverse LBSPs for the same question. A main purpose behind restricted information sets at individual LBSPs is that individuals tend to leave surveys for the same POI at one or at most just a couple LBSPs' sites which they frequently visit.

Second, LBSPs may alter their information sets by deleting a few surveys or including fake audits and return customized inquiry results for the eateries that are willing to pay or against those that decline to pay. Regardless of the possibility that LBSPs are not pernicious, they may return unfaithful question results affected by different assaults, for example, the Sybil assault whereby the same assailant can submit numerous fake surveys for the same POI.

Table 1: Survey on Secure Query Processing

Project Title	Algorithm/ Techniques	Remarks/ Problem Identification
Secure Top-k Query Processing via Untrusted Location-based Service Providers	Using of 2 scheme novel based approach	<ol style="list-style-type: none"> <li>1. It may falsely claim generating query results based on the review data from trusted data collectors.</li> <li>2. The non-malicious LBSPs may be compromised to return fake top-k query results.</li> </ol>
Sybil Guard: Defending Against Sybil Attacks via Social Networks	Sybil Guard, a novel protocol	<ol style="list-style-type: none"> <li>1. The Sybil Attack is a powerful threat faced by any decentralized distributed system that has no central, trusted authority to vouch for a one-to-one correspondence between users and identities.</li> <li>2. The important issues include how to bootstrap the social network and what applications can best benefit from Sybil Guard's fully decentralized approach.</li> </ol>
Sybil Limit: A Near Optimal Social Network Defense against Sybil Attacks	Sybil Limit's and Sybil Guard's approach	<ol style="list-style-type: none"> <li>1. Sybil Limit's guarantees on much smaller social networks with only 100 nodes. It cannot extract 100 node sub graphs from our social network data sets.</li> <li>2. It does not intend to implement Sybil Limit within the context of some real-world applications and demonstrate its utility.</li> </ol>
Query Integrity Assurance of Location based Services Accessing Outsourced Spatial Databases	Partially Materialized Digest scheme	<ol style="list-style-type: none"> <li>1. It does not extend algorithms to support more spatial query types such as spatial join, spatial path queries, etc.</li> </ol>



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### III. RELATED WORK

The solution for the two issues is to provide trusted data collectors as the focal points for gathering POI data. Specifically, data collectors can offer different motivators, for selling so as to invigorate survey entries and afterward benefit the audit data to individual LBSPs. Rather than submitting POI audits to individual LBSPs, individuals can now submit them to a couple data collectors to procure rewards. The data sets kept up by data collectors can in this manner be viewed as the union of the little data sets right now at individual LBSPs. Datagathering likewise makes it much less demanding and more attainable for data collectors to apply modern protections, for example, to sift through fake surveys from vindictive elements like Sybil assailants. Data collectors can be either new administration suppliers or all the more ideally existing ones with an expansive client base, for example, Google, Face book, Twitter, and MSN. A number of these administration suppliers (e.g., Google) have as of now been gathering surveys from their clients and offered open APIs for sending out chose data from their frameworks.

### IV. PROPOSED SYSTEM

In proposed system, three novel schemes are introduced to tackle the test for encouraging the handy sending and wide usage of the imagined framework. The key thought of our plans is that the information catcher pre register and verify some assistant data about its information set, which will be sold to LBSPs along with its information set. To answer a top-k inquiry reliably, a LBSP need to give back the right top-k POI information records and in addition appropriate proper authenticity and correctness proofs constructed from authenticated clues. The authenticity proof permits the query client to affirm that the inquiry come about just comprises of real information records from the trusted information auditor information set, and the rightness verification empowers the client to confirm that the returned top-k POIs are the one to fulfill the inquiry.

*The initial two schemes*, both target preview top-k questions yet vary in how authenticated hints are pre-processed and how authenticity and correctness proofs are developed and confirmed and also the related correspondence and calculation overhead.

*The third scheme*, based upon the first scheme, acknowledges productive and verifiable moving top-k questions. The adequacy and proficiency of our schemes are completely analyzed and evaluated.

#### A. Proposed Implementation Scheme

The proposed system implemented with three actors

- 1) *Data Collector*: Gathers the reviews about point of interest (POIs) from the data contributors.
- 2) *Data Contributors*: These are the people who submit POIs. Combined the data sets which gathered at individual LBSPs and provide centralized data sets.
- 3) *Location Based Service Providers (LBSP)*: purchases the POIs data sets from the data collector and set the stage for users to perform spatial top-k queries on POI in a certain region.

#### B. Secure Snapshot Top-K Query Processing

1) *Scheme 1*: Using of *M-hash tree* is for creating chaining ordered POIs according to zone. It allows efficient and secure Verification of the content of large data sets. Allow to verify any kind of data stored, handled, transferred in and between the computer.

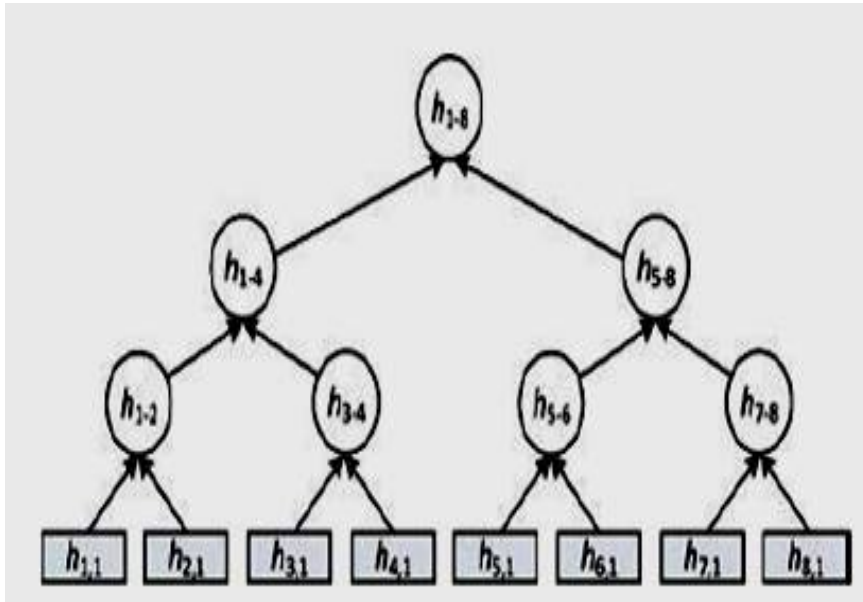


Figure 1: An example of constructing the M- hash tree

In **Scheme 1**, authenticated hints are created by chaining ordered POIs in every zone via cryptographic hash functions and then tie the POIs zone wise via an M- hash tree.

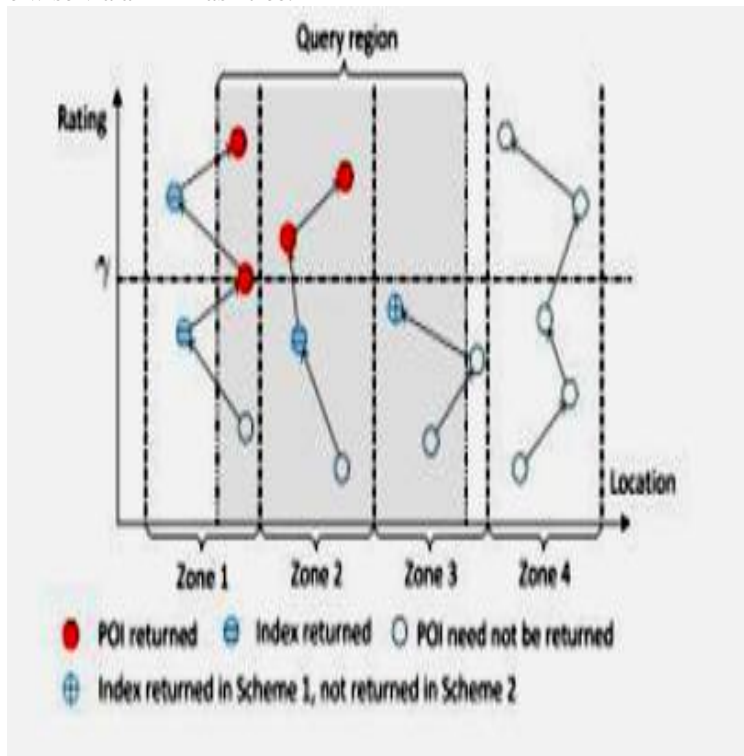


Figure 2: An example for scheme 1, the dots in zone i correspond to POI records from top to bottom

To perform true verification, the user first checks if zones I encloses the query region R. If I enclose R, then proceed with the following verifications in accordance with the mentioned true condition used in query processing:

So, the user there are exactly k data records in the query result with POI locations all in R, which correspond to the top-k POIs in R. If locates the lowest attribute-k rated with g.

2) *Scheme 2*: It points to the work by embedding or combining some information among nearby zones to reduce the amount of information return to the user. In scheme 2, LBSP return the information to the user, where no POIs are present. To implement the basic idea, the data collector binds to every POI data index some additional information about the POIs in adjacent zones. In particular, the data collector partitions the original M zones into non-overlapping macro zones, each consisting of m nearby zones, where m is a public system parameter. The LBSP purchases the original data set D, the signatures on M-tree root hashes, and all the intermediate results for constructing the M- hash tree of every interested POI category from the data collector.

### 3) *Scheme: Using Of Top-K Query*

This query is to return the k highest ranked answer or data sets quickly and efficiently. Reason for using Top-K query is to minimize the cost metrics that are associated with the retrieval data sets. To maximize the data set quality, this allows the users, not to overwhelm with irrelevant results. Updates in the top- k POIs may occur when a current top-k POI is no longer in the moving query region or when a new POI appears in the moving query region, which has an attribute-q rating higher than the lowest among the current top-k POIs. The user can directly tell when the first situation occurs based on the current top-k POIs he knows, in which case he can issue a new snapshot top-k query for the current query region.

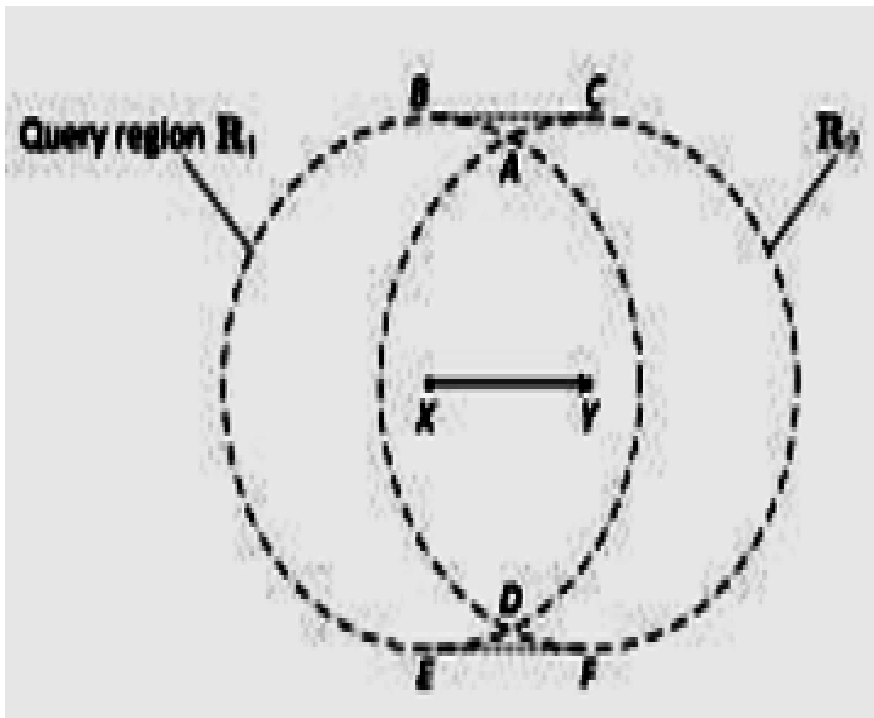


Figure 3: An example of two consecutive snapshot top-k query

The user, however, cannot tell when the second situation will occur. Without a sound defense in place, the LBSP can choose not to inform the user about updated top-k POIs in the second situation. Scheme 3 is done by using neither scheme 1 nor scheme 2. The data sets has been pre-processed by the data collector when it is selected. Using of scheme 1 is due to space constraints and for without loss of generality scheme 2 is used.



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**V. PERFORMANCE ANALYSIS**

**A. SCHEME 1 ANALYSIS**

Any incorrect and/or inauthentic query result from a misbehaving LBSP can be detected. The expected number of hash computations the user performs to verify the query result under Scheme 1 is formulated as:

$$E[N_{hash,1}] = k + |Z| \cdot \frac{(k + \delta)n + 1}{\delta n + 1} + \sum_{j=1}^{d-1} 2^{j-1} (1 - (1 - 2^{-(j-1)})^{|Z|}) \dots\dots\dots Eq (1)$$

The additional communication overhead between the LBSP and the user incurred by Scheme 1 is formulated.

**B. SCHEME 2 ANALYSIS**

Incorrect and/or inauthentic query result from a misbehaving LBSP is detected. The expected number of hash computations the user performs to verify the query result under Scheme 2 is formulated as:

$$E[N_{hash,2}] = |Z|\mu_1 + \sum_{j=1}^{d-1} 2^{j-1} (1 - (1 - 2^{-(j-1)})^{|Z|(1-\mu_2^j)}) \dots\dots Eq(2)$$

where  $\mu_1 = (n - n\mu_2 + 1 - \mu_2^n)$  and  $\mu_2 = \frac{\delta n - k + 1}{\delta n + 1}$ .

The expected additional communication overhead Scheme 2 incurs between the LBSP and user is bounded according to proposed theme.

Table 2: Default Simulation Settings

Para.	Val.	Para.	Val.	Para.	Val.	Para.	Val.
M	10000	M	100	n	100	Δ	10
k	5	D	14	D	20	L <sub>h</sub>	160
L <sub>loc</sub>	20	L <sub>sig</sub>	160	L <sub>r</sub>	10		

(Where **M** is no. of zones in merkle hash tree - **k** is records in query result - **L<sub>loc</sub>** is POI location- **m** is candidate of macro zone - **L<sub>sig</sub>** is data collectors signature- **n** is number of indexes - **d** is data sets - **L<sub>r</sub>** is location of radius for circle - **δ** is candidate zone - **L<sub>h</sub>** is bit lengths of hash value H. )

**C. SCHEME 3 ANALYSIS**

Misbehavior of the LBSP, including returning incorrect/inauthentic query result and omitting complete query results, will be eventually detected under this Scheme.

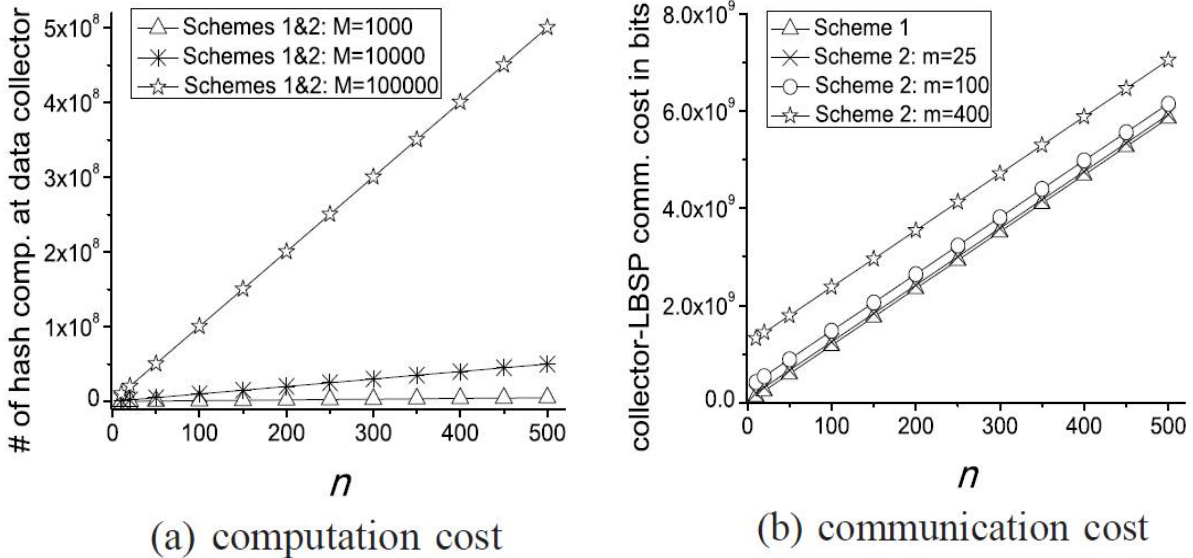


Figure 4: COMPARISON AMONG SCHEME 1 AND 2

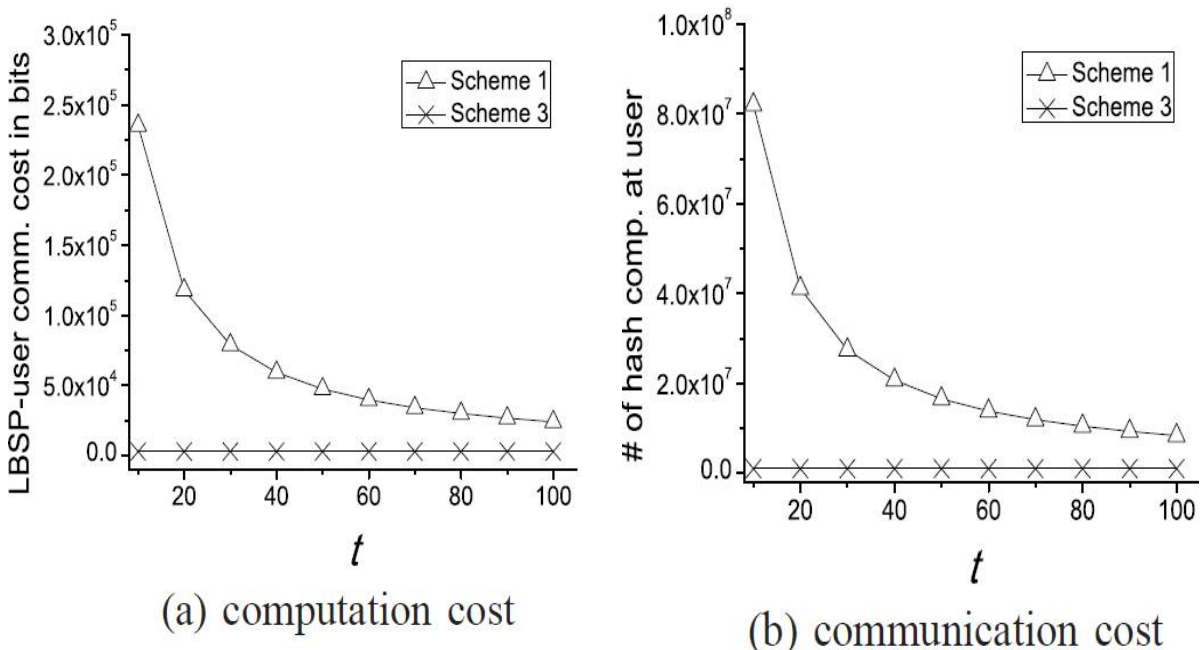


Figure 5: COMPARISON AMONG SCHEME 1 AND 3

the proposed one in this paper In figure 5 As the above graphs shows scheme 3 is proved as best in computation and communication cost factor, which is.

### V. CONCLUSION

This is the efficient distributed system for collaborative location-based information generating and sharing. The three good schemes implemented to enable secure top-k query processing via untrusted LBSPs for fostering the practical



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deployment and wide use of the system envisioning. The location based information generation and sharing for distributed system enables a secure processing, which enables the users to verify authenticity and correctness of the query result for untrusted location using novel schemes.

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### BIOGRAPHY

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